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QUIC: A UDP-Based Multiplexed and Secure Transport
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Abstract

This document defines the core of the QUIC transport protocol. Accompanying documents describe QUIC's loss detection and congestion control and the use of TLS for key negotiation.

Note to Readers

Discussion of this draft takes place on the QUIC working group mailing list (quic@ietf.org (<mailto:quic@ietf.org>)), which is archived at https://mailarchive.ietf.org/arch/search/?email_list=quic

Working Group information can be found at <https://github.com/quicwg>; source code and issues list for this draft can be found at <https://github.com/quicwg/base-drafts/labels/-transport>.

Status of This Memo

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1. Introduction

QUIC is a multiplexed and secure general-purpose transport protocol that provides:

- * Stream multiplexing
- * Stream and connection-level flow control
- * Low-latency connection establishment

- * Connection migration and resilience to NAT rebinding
- * Authenticated and encrypted header and payload

QUIC uses UDP as a substrate to avoid requiring changes to legacy client operating systems and middleboxes. QUIC authenticates all of its headers and encrypts most of the data it exchanges, including its signaling, to avoid incurring a dependency on middleboxes.

1.1. Document Structure

This document describes the core QUIC protocol and is structured as follows:

- * Streams are the basic service abstraction that QUIC provides.
 - [Section 2](#) describes core concepts related to streams,
 - [Section 3](#) provides a reference model for stream states, and
 - [Section 4](#) outlines the operation of flow control.
- * Connections are the context in which QUIC endpoints communicate.
 - [Section 5](#) describes core concepts related to connections,
 - [Section 6](#) describes version negotiation,
 - [Section 7](#) details the process for establishing connections,
 - [Section 8](#) specifies critical denial of service mitigation mechanisms,
 - [Section 9](#) describes how endpoints migrate a connection to a new network path,
 - [Section 10](#) lists the options for terminating an open connection, and
 - [Section 11](#) provides general guidance for error handling.
- * Packets and frames are the basic unit used by QUIC to communicate.
 - [Section 12](#) describes concepts related to packets and frames,
 - [Section 13](#) defines models for the transmission, retransmission, and acknowledgement of data, and

- [Section 14](#) specifies rules for managing the size of packets.
- * Finally, encoding details of QUIC protocol elements are described in:
- [Section 15](#) (Versions),
 - [Section 16](#) (Integer Encoding),
 - [Section 17](#) (Packet Headers),
 - [Section 18](#) (Transport Parameters),
 - [Section 19](#) (Frames), and
 - [Section 20](#) (Errors).

Accompanying documents describe QUIC's loss detection and congestion control [[QUIC-RECOVERY](#)], and the use of TLS for key negotiation [[QUIC-TLS](#)].

This document defines QUIC version 1, which conforms to the protocol invariants in [[QUIC-INVARIANTS](#)].

To refer to QUIC version 1, cite this document. References to the limited set of version-independent properties of QUIC can cite [[QUIC-INVARIANTS](#)].

1.2. Terms and Definitions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [BCP 14](#) [[RFC2119](#)] [[RFC8174](#)] when, and only when, they appear in all capitals, as shown here.

Commonly used terms in the document are described below.

QUIC: The transport protocol described by this document. QUIC is a name, not an acronym.

QUIC packet: A complete processable unit of QUIC that can be encapsulated in a UDP datagram. Multiple QUIC packets can be encapsulated in a single UDP datagram.

Ack-eliciting Packet: A QUIC packet that contains frames other than ACK, PADDING, and CONNECTION_CLOSE. These cause a recipient to send an acknowledgment; see [Section 13.2.1](#).

Out-of-order packet: A packet that does not increase the largest received packet number for its packet number space ([Section 12.3](#)) by exactly one. A packet can arrive out of order if it is delayed, if earlier packets are lost or delayed, or if the sender intentionally skips a packet number.

Endpoint: An entity that can participate in a QUIC connection by generating, receiving, and processing QUIC packets. There are only two types of endpoint in QUIC: client and server.

Client: The endpoint initiating a QUIC connection.

Server: The endpoint accepting incoming QUIC connections.

Address: When used without qualification, the tuple of IP version, IP address, UDP protocol, and UDP port number that represents one end of a network path.

Connection ID: An opaque identifier that is used to identify a QUIC connection at an endpoint. Each endpoint sets a value for its peer to include in packets sent towards the endpoint.

Stream: A unidirectional or bidirectional channel of ordered bytes within a QUIC connection. A QUIC connection can carry multiple simultaneous streams.

Application: An entity that uses QUIC to send and receive data.

1.3. Notational Conventions

Packet and frame diagrams in this document use a bespoke format. The purpose of this format is to summarize, not define, protocol elements. Prose defines the complete semantics and details of structures.

Complex fields are named and then followed by a list of fields surrounded by a pair of matching braces. Each field in this list is separated by commas.

Individual fields include length information, plus indications about fixed value, optionality, or repetitions. Individual fields use the following notational conventions, with all lengths in bits:

$x (A)$: Indicates that x is A bits long

$x (i)$: Indicates that x uses the variable-length encoding in [Section 16](#)

x (A..B): Indicates that x can be any length from A to B; A can be omitted to indicate a minimum of zero bits and B can be omitted to indicate no set upper limit; values in this format always end on an octet boundary

x (?) = C: Indicates that x has a fixed value of C

x (?) = C..D: Indicates that x has a value in the range from C to D, inclusive

[x (E)]: Indicates that x is optional (and has length of E)

x (E) ...: Indicates that x is repeated zero or more times (and that each instance is length E)

This document uses network byte order (that is, big endian) values. Fields are placed starting from the high-order bits of each byte.

By convention, individual fields reference a complex field by using the name of the complex field.

For example:

```
Example Structure {
  One-bit Field (1),
  7-bit Field with Fixed Value (7) = 61,
  Field with Variable-Length Integer (i),
  Arbitrary-Length Field (..),
  Variable-Length Field (8..24),
  Field With Minimum Length (16..),
  Field With Maximum Length (..128),
  [Optional Field (64)],
  Repeated Field (8) ...,
}
```

Figure 1: Example Format

2. Streams

Streams in QUIC provide a lightweight, ordered byte-stream abstraction to an application. Streams can be unidirectional or bidirectional. An alternative view of QUIC unidirectional streams is a "message" abstraction of practically unlimited length.

Streams can be created by sending data. Other processes associated with stream management - ending, cancelling, and managing flow control - are all designed to impose minimal overheads. For instance, a single STREAM frame ([Section 19.8](#)) can open, carry data for, and close a stream. Streams can also be long-lived and can last the entire duration of a connection.

Streams can be created by either endpoint, can concurrently send data interleaved with other streams, and can be cancelled. QUIC does not provide any means of ensuring ordering between bytes on different streams.

QUIC allows for an arbitrary number of streams to operate concurrently and for an arbitrary amount of data to be sent on any stream, subject to flow control constraints and stream limits; see [Section 4](#).

2.1. Stream Types and Identifiers

Streams can be unidirectional or bidirectional. Unidirectional streams carry data in one direction: from the initiator of the stream to its peer. Bidirectional streams allow for data to be sent in both directions.

Streams are identified within a connection by a numeric value, referred to as the stream ID. A stream ID is a 62-bit integer (0 to $2^{62}-1$) that is unique for all streams on a connection. Stream IDs are encoded as variable-length integers; see [Section 16](#). A QUIC endpoint MUST NOT reuse a stream ID within a connection.

The least significant bit (0x1) of the stream ID identifies the initiator of the stream. Client-initiated streams have even-numbered stream IDs (with the bit set to 0), and server-initiated streams have odd-numbered stream IDs (with the bit set to 1).

The second least significant bit (0x2) of the stream ID distinguishes between bidirectional streams (with the bit set to 0) and unidirectional streams (with the bit set to 1).

The least significant two bits from a stream ID therefore identify a stream as one of four types, as summarized in Table 1.

Bits	Stream Type
0x0	Client-Initiated, Bidirectional
0x1	Server-Initiated, Bidirectional
0x2	Client-Initiated, Unidirectional
0x3	Server-Initiated, Unidirectional

Table 1: Stream ID Types

Within each type, streams are created with numerically increasing stream IDs. A stream ID that is used out of order results in all streams of that type with lower-numbered stream IDs also being opened.

The first bidirectional stream opened by the client has a stream ID of 0.

2.2. Sending and Receiving Data

STREAM frames ([Section 19.8](#)) encapsulate data sent by an application. An endpoint uses the Stream ID and Offset fields in STREAM frames to place data in order.

Endpoints **MUST** be able to deliver stream data to an application as an ordered byte-stream. Delivering an ordered byte-stream requires that an endpoint buffer any data that is received out of order, up to the advertised flow control limit.

QUIC makes no specific allowances for delivery of stream data out of order. However, implementations **MAY** choose to offer the ability to deliver data out of order to a receiving application.

An endpoint could receive data for a stream at the same stream offset multiple times. Data that has already been received can be discarded. The data at a given offset **MUST NOT** change if it is sent multiple times; an endpoint **MAY** treat receipt of different data at the same offset within a stream as a connection error of type `PROTOCOL_VIOLATION`.

Streams are an ordered byte-stream abstraction with no other structure visible to QUIC. STREAM frame boundaries are not expected to be preserved when data is transmitted, retransmitted after packet loss, or delivered to the application at a receiver.

An endpoint **MUST NOT** send data on any stream without ensuring that it is within the flow control limits set by its peer. Flow control is described in detail in [Section 4](#).

2.3. Stream Prioritization

Stream multiplexing can have a significant effect on application performance if resources allocated to streams are correctly prioritized.

QUIC does not provide a mechanism for exchanging prioritization information. Instead, it relies on receiving priority information from the application that uses QUIC.

A QUIC implementation **SHOULD** provide ways in which an application can indicate the relative priority of streams. When deciding the streams to which resources are dedicated, the implementation **SHOULD** use the information provided by the application.

2.4. Required Operations on Streams

There are certain operations that an application **MUST** be able to perform when interacting with QUIC streams. This document does not specify an API, but any implementation of this version of QUIC **MUST** expose the ability to perform the operations described in this section on a QUIC stream.

On the sending part of a stream, application protocols need to be able to:

- * write data, understanding when stream flow control credit ([Section 4.1](#)) has successfully been reserved to send the written data;
- * end the stream (clean termination), resulting in a STREAM frame ([Section 19.8](#)) with the FIN bit set; and
- * reset the stream (abrupt termination), resulting in a RESET_STREAM frame ([Section 19.4](#)), if the stream was not already in a terminal state.

On the receiving part of a stream, application protocols need to be able to:

- * read data; and
- * abort reading of the stream and request closure, possibly resulting in a STOP_SENDING frame ([Section 19.5](#)).

Applications also need to be informed of state changes on streams, including when the peer has opened or reset a stream, when a peer aborts reading on a stream, when new data is available, and when data can or cannot be written to the stream due to flow control.

3. Stream States

This section describes streams in terms of their send or receive components. Two state machines are described: one for the streams on which an endpoint transmits data ([Section 3.1](#)), and another for streams on which an endpoint receives data ([Section 3.2](#)).

Unidirectional streams use the applicable state machine directly. Bidirectional streams use both state machines. For the most part, the use of these state machines is the same whether the stream is unidirectional or bidirectional. The conditions for opening a stream are slightly more complex for a bidirectional stream because the opening of either the send or receive side causes the stream to open in both directions.

An endpoint **MUST** open streams of the same type in increasing order of stream ID.

Note: These states are largely informative. This document uses stream states to describe rules for when and how different types of frames can be sent and the reactions that are expected when different types of frames are received. Though these state machines are intended to be useful in implementing QUIC, these states aren't intended to constrain implementations. An implementation can define a different state machine as long as its behavior is consistent with an implementation that implements these states.

3.1. Sending Stream States

Figure 2 shows the states for the part of a stream that sends data to a peer.

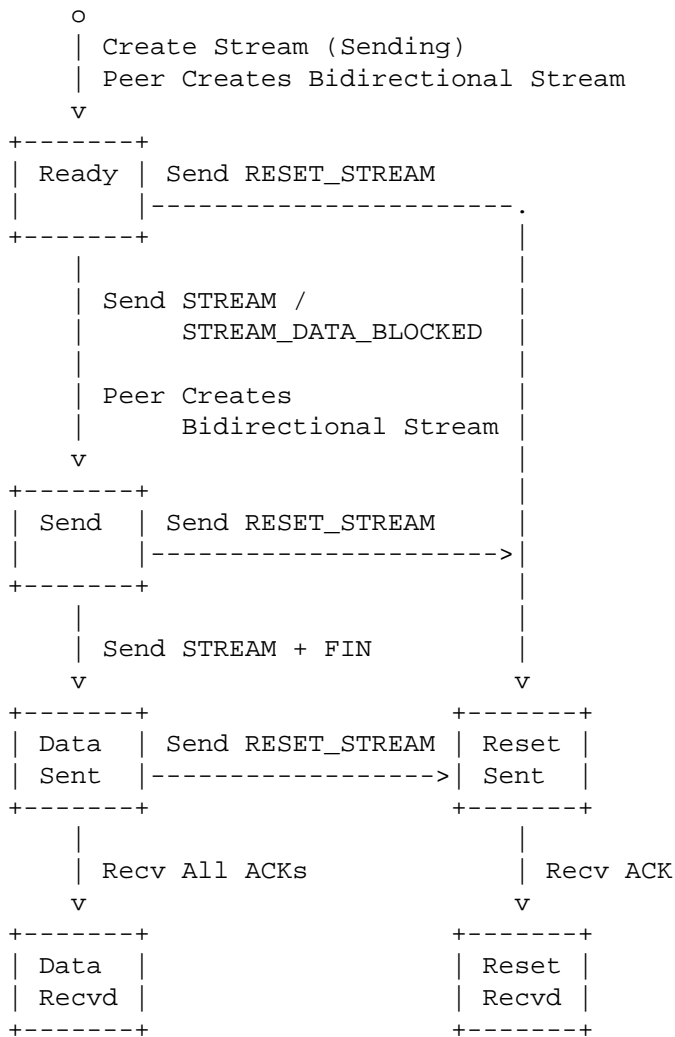


Figure 2: States for Sending Parts of Streams

The sending part of a stream that the endpoint initiates (types 0 and 2 for clients, 1 and 3 for servers) is opened by the application. The "Ready" state represents a newly created stream that is able to accept data from the application. Stream data might be buffered in this state in preparation for sending.

Sending the first STREAM or STREAM_DATA_BLOCKED frame causes a sending part of a stream to enter the "Send" state. An implementation might choose to defer allocating a stream ID to a stream until it sends the first STREAM frame and enters this state, which can allow for better stream prioritization.

The sending part of a bidirectional stream initiated by a peer (type 0 for a server, type 1 for a client) starts in the "Ready" state when the receiving part is created.

In the "Send" state, an endpoint transmits - and retransmits as necessary - stream data in STREAM frames. The endpoint respects the flow control limits set by its peer, and continues to accept and process MAX_STREAM_DATA frames. An endpoint in the "Send" state generates STREAM_DATA_BLOCKED frames if it is blocked from sending by stream or connection flow control limits [Section 4.1](#).

After the application indicates that all stream data has been sent and a STREAM frame containing the FIN bit is sent, the sending part of the stream enters the "Data Sent" state. From this state, the endpoint only retransmits stream data as necessary. The endpoint does not need to check flow control limits or send STREAM_DATA_BLOCKED frames for a stream in this state. MAX_STREAM_DATA frames might be received until the peer receives the final stream offset. The endpoint can safely ignore any MAX_STREAM_DATA frames it receives from its peer for a stream in this state.

Once all stream data has been successfully acknowledged, the sending part of the stream enters the "Data Recvd" state, which is a terminal state.

From any of the "Ready", "Send", or "Data Sent" states, an application can signal that it wishes to abandon transmission of stream data. Alternatively, an endpoint might receive a STOP_SENDING frame from its peer. In either case, the endpoint sends a RESET_STREAM frame, which causes the stream to enter the "Reset Sent" state.

An endpoint MAY send a RESET_STREAM as the first frame that mentions a stream; this causes the sending part of that stream to open and then immediately transition to the "Reset Sent" state.

Once a packet containing a RESET_STREAM has been acknowledged, the sending part of the stream enters the "Reset Recvd" state, which is a terminal state.

3.2. Receiving Stream States

Figure 3 shows the states for the part of a stream that receives data from a peer. The states for a receiving part of a stream mirror only some of the states of the sending part of the stream at the peer. The receiving part of a stream does not track states on the sending part that cannot be observed, such as the "Ready" state. Instead, the receiving part of a stream tracks the delivery of data to the application, some of which cannot be observed by the sender.

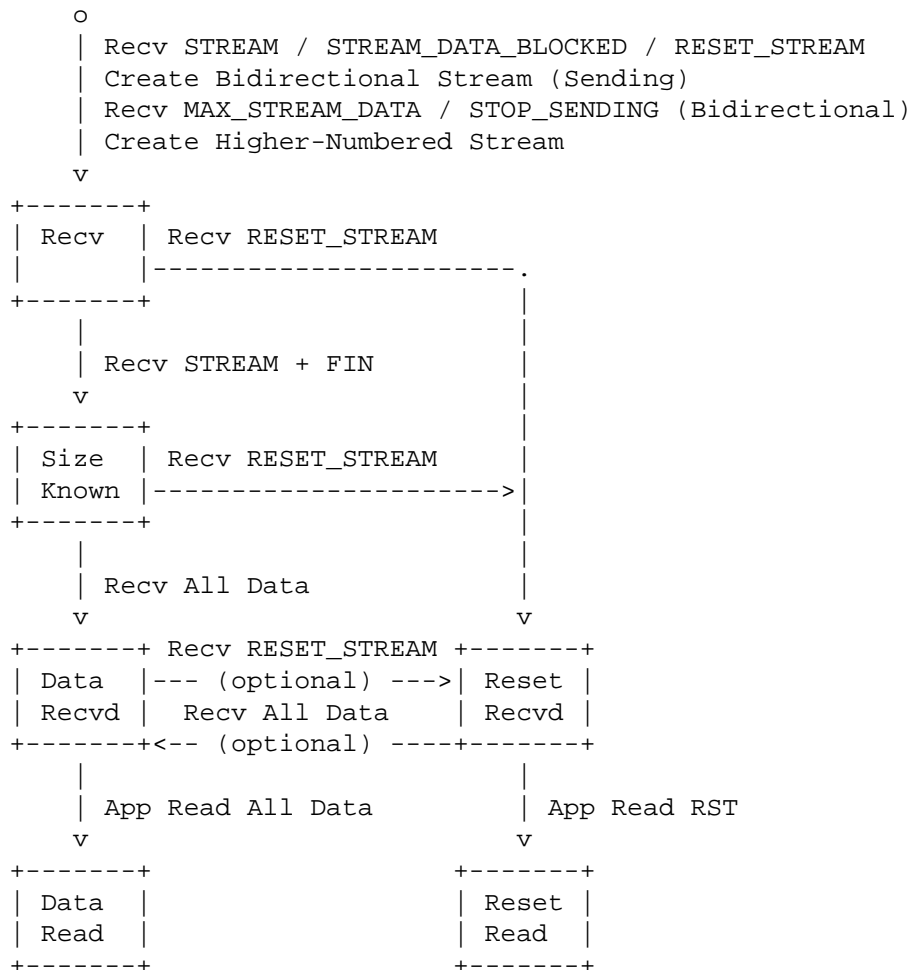


Figure 3: States for Receiving Parts of Streams

The receiving part of a stream initiated by a peer (types 1 and 3 for a client, or 0 and 2 for a server) is created when the first STREAM, STREAM_DATA_BLOCKED, or RESET_STREAM frame is received for that stream. For bidirectional streams initiated by a peer, receipt of a

MAX_STREAM_DATA or STOP_SENDING frame for the sending part of the stream also creates the receiving part. The initial state for the receiving part of stream is "Recv".

The receiving part of a stream enters the "Recv" state when the sending part of a bidirectional stream initiated by the endpoint (type 0 for a client, type 1 for a server) enters the "Ready" state.

An endpoint opens a bidirectional stream when a MAX_STREAM_DATA or STOP_SENDING frame is received from the peer for that stream. Receiving a MAX_STREAM_DATA frame for an unopened stream indicates that the remote peer has opened the stream and is providing flow control credit. Receiving a STOP_SENDING frame for an unopened stream indicates that the remote peer no longer wishes to receive data on this stream. Either frame might arrive before a STREAM or STREAM_DATA_BLOCKED frame if packets are lost or reordered.

Before a stream is created, all streams of the same type with lower-numbered stream IDs MUST be created. This ensures that the creation order for streams is consistent on both endpoints.

In the "Recv" state, the endpoint receives STREAM and STREAM_DATA_BLOCKED frames. Incoming data is buffered and can be reassembled into the correct order for delivery to the application. As data is consumed by the application and buffer space becomes available, the endpoint sends MAX_STREAM_DATA frames to allow the peer to send more data.

When a STREAM frame with a FIN bit is received, the final size of the stream is known; see [Section 4.4](#). The receiving part of the stream then enters the "Size Known" state. In this state, the endpoint no longer needs to send MAX_STREAM_DATA frames, it only receives any retransmissions of stream data.

Once all data for the stream has been received, the receiving part enters the "Data Recvd" state. This might happen as a result of receiving the same STREAM frame that causes the transition to "Size Known". After all data has been received, any STREAM or STREAM_DATA_BLOCKED frames for the stream can be discarded.

The "Data Recvd" state persists until stream data has been delivered to the application. Once stream data has been delivered, the stream enters the "Data Read" state, which is a terminal state.

Receiving a RESET_STREAM frame in the "Recv" or "Size Known" states causes the stream to enter the "Reset Recvd" state. This might cause the delivery of stream data to the application to be interrupted.

It is possible that all stream data is received when a RESET_STREAM is received (that is, from the "Data Recvd" state). Similarly, it is possible for remaining stream data to arrive after receiving a RESET_STREAM frame (the "Reset Recvd" state). An implementation is free to manage this situation as it chooses.

Sending RESET_STREAM means that an endpoint cannot guarantee delivery of stream data; however there is no requirement that stream data not be delivered if a RESET_STREAM is received. An implementation MAY interrupt delivery of stream data, discard any data that was not consumed, and signal the receipt of the RESET_STREAM. A RESET_STREAM signal might be suppressed or withheld if stream data is completely received and is buffered to be read by the application. If the RESET_STREAM is suppressed, the receiving part of the stream remains in "Data Recvd".

Once the application receives the signal indicating that the stream was reset, the receiving part of the stream transitions to the "Reset Read" state, which is a terminal state.

3.3. Permitted Frame Types

The sender of a stream sends just three frame types that affect the state of a stream at either sender or receiver: STREAM (Section 19.8), STREAM_DATA_BLOCKED (Section 19.13), and RESET_STREAM (Section 19.4).

A sender MUST NOT send any of these frames from a terminal state ("Data Recvd" or "Reset Recvd"). A sender MUST NOT send a STREAM or STREAM_DATA_BLOCKED frame for a stream in the "Reset Sent" state or any terminal state, that is, after sending a RESET_STREAM frame. A receiver could receive any of these three frames in any state, due to the possibility of delayed delivery of packets carrying them.

The receiver of a stream sends MAX_STREAM_DATA (Section 19.10) and STOP_SENDING frames (Section 19.5).

The receiver only sends MAX_STREAM_DATA in the "Recv" state. A receiver MAY send STOP_SENDING in any state where it has not received a RESET_STREAM frame; that is states other than "Reset Recvd" or "Reset Read". However there is little value in sending a STOP_SENDING frame in the "Data Recvd" state, since all stream data has been received. A sender could receive either of these two frames in any state as a result of delayed delivery of packets.

3.4. Bidirectional Stream States

A bidirectional stream is composed of sending and receiving parts. Implementations may represent states of the bidirectional stream as composites of sending and receiving stream states. The simplest model presents the stream as "open" when either sending or receiving parts are in a non-terminal state and "closed" when both sending and receiving streams are in terminal states.

Table 2 shows a more complex mapping of bidirectional stream states that loosely correspond to the stream states in HTTP/2 [HTTP2]. This shows that multiple states on sending or receiving parts of streams are mapped to the same composite state. Note that this is just one possibility for such a mapping; this mapping requires that data is acknowledged before the transition to a "closed" or "half-closed" state.

Sending Part	Receiving Part	Composite State
No Stream/Ready	No Stream/Recv *1	idle
Ready/Send/Data Sent	Recv/Size Known	open
Ready/Send/Data Sent	Data Recvd/Data Read	half-closed (remote)
Ready/Send/Data Sent	Reset Recvd/Reset Read	half-closed (remote)
Data Recvd	Recv/Size Known	half-closed (local)
Reset Sent/Reset Recv	Recv/Size Known	half-closed (local)
Reset Sent/Reset Recv	Data Recvd/Data Read	closed
Reset Sent/Reset Recv	Reset Recvd/Reset Read	closed
Data Recvd	Data Recvd/Data Read	closed
Data Recvd	Reset Recvd/Reset Read	closed

Table 2: Possible Mapping of Stream States to HTTP/2

Note (*1): A stream is considered "idle" if it has not yet been created, or if the receiving part of the stream is in the "Recv" state without yet having received any frames.

3.5. Solicited State Transitions

If an application is no longer interested in the data it is receiving on a stream, it can abort reading the stream and specify an application error code.

If the stream is in the "Recv" or "Size Known" states, the transport SHOULD signal this by sending a STOP_SENDING frame to prompt closure of the stream in the opposite direction. This typically indicates that the receiving application is no longer reading data it receives from the stream, but it is not a guarantee that incoming data will be ignored.

STREAM frames received after sending a STOP_SENDING frame are still counted toward connection and stream flow control, even though these frames can be discarded upon receipt.

A STOP_SENDING frame requests that the receiving endpoint send a RESET_STREAM frame. An endpoint that receives a STOP_SENDING frame MUST send a RESET_STREAM frame if the stream is in the Ready or Send state. If the stream is in the Data Sent state and any outstanding data is declared lost, an endpoint SHOULD send a RESET_STREAM frame in lieu of a retransmission.

An endpoint SHOULD copy the error code from the STOP_SENDING frame to the RESET_STREAM frame it sends, but MAY use any application error code. The endpoint that sends a STOP_SENDING frame MAY ignore the error code carried in any RESET_STREAM frame it receives.

If the STOP_SENDING frame is received on a stream that is already in the "Data Sent" state, an endpoint that wishes to cease retransmission of previously-sent STREAM frames on that stream MUST first send a RESET_STREAM frame.

STOP_SENDING SHOULD only be sent for a stream that has not been reset by the peer. STOP_SENDING is most useful for streams in the "Recv" or "Size Known" states.

An endpoint is expected to send another STOP_SENDING frame if a packet containing a previous STOP_SENDING is lost. However, once either all stream data or a RESET_STREAM frame has been received for the stream - that is, the stream is in any state other than "Recv" or "Size Known" - sending a STOP_SENDING frame is unnecessary.

An endpoint that wishes to terminate both directions of a bidirectional stream can terminate one direction by sending a RESET_STREAM frame, and it can encourage prompt termination in the opposite direction by sending a STOP_SENDING frame.

4. Flow Control

It is necessary to limit the amount of data that a receiver could buffer, to prevent a fast sender from overwhelming a slow receiver, or to prevent a malicious sender from consuming a large amount of memory at a receiver. To enable a receiver to limit memory commitment to a connection and to apply back pressure on the sender, streams are flow controlled both individually and as an aggregate. A QUIC receiver controls the maximum amount of data the sender can send on a stream at any time, as described in [Section 4.1](#) and [Section 4.2](#).

Similarly, to limit concurrency within a connection, a QUIC endpoint controls the maximum cumulative number of streams that its peer can initiate, as described in [Section 4.5](#).

Data sent in CRYPTO frames is not flow controlled in the same way as stream data. QUIC relies on the cryptographic protocol implementation to avoid excessive buffering of data; see [\[QUIC-TLS\]](#). The implementation SHOULD provide an interface to QUIC to tell it about its buffering limits so that there is not excessive buffering at multiple layers.

4.1. Data Flow Control

QUIC employs a credit-based flow-control scheme similar to that in HTTP/2 [\[HTTP2\]](#), where a receiver advertises the number of bytes it is prepared to receive on a given stream and for the entire connection. This leads to two levels of data flow control in QUIC:

- * Stream flow control, which prevents a single stream from consuming the entire receive buffer for a connection by limiting the amount of data that can be sent on any stream.
- * Connection flow control, which prevents senders from exceeding a receiver's buffer capacity for the connection, by limiting the total bytes of stream data sent in STREAM frames on all streams.

A receiver sets initial credits for all streams by sending transport parameters during the handshake ([Section 7.4](#)). A receiver sends MAX_STREAM_DATA ([Section 19.10](#)) or MAX_DATA ([Section 19.9](#)) frames to the sender to advertise additional credit.

A receiver advertises credit for a stream by sending a `MAX_STREAM_DATA` frame with the Stream ID field set appropriately. A `MAX_STREAM_DATA` frame indicates the maximum absolute byte offset of a stream. A receiver could use the current offset of data consumed to determine the flow control offset to be advertised. A receiver **MAY** send `MAX_STREAM_DATA` frames in multiple packets in order to make sure that the sender receives an update before running out of flow control credit, even if one of the packets is lost.

A receiver advertises credit for a connection by sending a `MAX_DATA` frame, which indicates the maximum of the sum of the absolute byte offsets of all streams. A receiver maintains a cumulative sum of bytes received on all streams, which is used to check for flow control violations. A receiver might use a sum of bytes consumed on all streams to determine the maximum data limit to be advertised.

A receiver can advertise a larger offset by sending `MAX_STREAM_DATA` or `MAX_DATA` frames. Once a receiver advertises an offset, it **MAY** advertise a smaller offset, but this has no effect.

A receiver **MUST** close the connection with a `FLOW_CONTROL_ERROR` error ([Section 11](#)) if the sender violates the advertised connection or stream data limits.

A sender **MUST** ignore any `MAX_STREAM_DATA` or `MAX_DATA` frames that do not increase flow control limits.

If a sender runs out of flow control credit, it will be unable to send new data and is considered blocked. A sender **SHOULD** send a `STREAM_DATA_BLOCKED` or `DATA_BLOCKED` frame to indicate it has data to write but is blocked by flow control limits. If a sender is blocked for a period longer than the idle timeout ([Section 10.2](#)), the connection might be closed even when data is available for transmission. To keep the connection from closing, a sender that is flow control limited **SHOULD** periodically send a `STREAM_DATA_BLOCKED` or `DATA_BLOCKED` frame when it has no ack-eliciting packets in flight.

4.2. Flow Credit Increments

Implementations decide when and how much credit to advertise in `MAX_STREAM_DATA` and `MAX_DATA` frames, but this section offers a few considerations.

To avoid blocking a sender, a receiver can send a `MAX_STREAM_DATA` or `MAX_DATA` frame multiple times within a round trip or send it early enough to allow for recovery from loss of the frame.

Control frames contribute to connection overhead. Therefore, frequently sending `MAX_STREAM_DATA` and `MAX_DATA` frames with small changes is undesirable. On the other hand, if updates are less frequent, larger increments to limits are necessary to avoid blocking a sender, requiring larger resource commitments at the receiver. There is a trade-off between resource commitment and overhead when determining how large a limit is advertised.

A receiver can use an autotuning mechanism to tune the frequency and amount of advertised additional credit based on a round-trip time estimate and the rate at which the receiving application consumes data, similar to common TCP implementations. As an optimization, an endpoint could send frames related to flow control only when there are other frames to send or when a peer is blocked, ensuring that flow control does not cause extra packets to be sent.

A blocked sender is not required to send `STREAM_DATA_BLOCKED` or `DATA_BLOCKED` frames. Therefore, a receiver **MUST NOT** wait for a `STREAM_DATA_BLOCKED` or `DATA_BLOCKED` frame before sending a `MAX_STREAM_DATA` or `MAX_DATA` frame; doing so could result in the sender being blocked for the rest of the connection. Even if the sender sends these frames, waiting for them will result in the sender being blocked for at least an entire round trip.

When a sender receives credit after being blocked, it might be able to send a large amount of data in response, resulting in short-term congestion; see Section 6.9 in [\[QUIC-RECOVERY\]](#) for a discussion of how a sender can avoid this congestion.

4.3. Handling Stream Cancellation

Endpoints need to eventually agree on the amount of flow control credit that has been consumed, to avoid either exceeding flow control limits or deadlocking.

On receipt of a `RESET_STREAM` frame, an endpoint will tear down state for the matching stream and ignore further data arriving on that stream. Without the offset included in `RESET_STREAM`, the two endpoints could disagree on the number of bytes that count towards connection flow control.

To remedy this issue, a `RESET_STREAM` frame ([Section 19.4](#)) includes the final size of data sent on the stream. On receiving a `RESET_STREAM` frame, a receiver definitively knows how many bytes were sent on that stream before the `RESET_STREAM` frame, and the receiver **MUST** use the final size of the stream to account for all bytes sent on the stream in its connection level flow controller.

RESET_STREAM terminates one direction of a stream abruptly. For a bidirectional stream, RESET_STREAM has no effect on data flow in the opposite direction. Both endpoints MUST maintain flow control state for the stream in the unterminated direction until that direction enters a terminal state, or until one of the endpoints sends CONNECTION_CLOSE.

4.4. Stream Final Size

The final size is the amount of flow control credit that is consumed by a stream. Assuming that every contiguous byte on the stream was sent once, the final size is the number of bytes sent. More generally, this is one higher than the offset of the byte with the largest offset sent on the stream, or zero if no bytes were sent.

The final size of a stream is always signaled to the recipient. The final size is the sum of the Offset and Length fields of a STREAM frame with a FIN flag, noting that these fields might be implicit. Alternatively, the Final Size field of a RESET_STREAM frame carries this value. This ensures that all ways that a stream can be closed result in the number of bytes on the stream being reliably transmitted. This guarantees that both endpoints agree on how much flow control credit was consumed by the stream.

An endpoint will know the final size for a stream when the receiving part of the stream enters the "Size Known" or "Reset Recvd" state ([Section 3](#)).

An endpoint MUST NOT send data on a stream at or beyond the final size.

Once a final size for a stream is known, it cannot change. If a RESET_STREAM or STREAM frame is received indicating a change in the final size for the stream, an endpoint SHOULD respond with a FINAL_SIZE_ERROR error; see [Section 11](#). A receiver SHOULD treat receipt of data at or beyond the final size as a FINAL_SIZE_ERROR error, even after a stream is closed. Generating these errors is not mandatory, but only because requiring that an endpoint generate these errors also means that the endpoint needs to maintain the final size state for closed streams, which could mean a significant state commitment.

4.5. Controlling Concurrency

An endpoint limits the cumulative number of incoming streams a peer can open. Only streams with a stream ID less than $(\text{max_stream} * 4 + \text{initial_stream_id_for_type})$ can be opened; see Table 5. Initial limits are set in the transport parameters (see [Section 18.2](#)) and subsequently limits are advertised using MAX_STREAMS frames ([Section 19.11](#)). Separate limits apply to unidirectional and bidirectional streams.

If a max_streams transport parameter or MAX_STREAMS frame is received with a value greater than 2^{60} , this would allow a maximum stream ID that cannot be expressed as a variable-length integer; see [Section 16](#). If either is received, the connection MUST be closed immediately with a connection error of type STREAM_LIMIT_ERROR; see [Section 10.3](#).

Endpoints MUST NOT exceed the limit set by their peer. An endpoint that receives a frame with a stream ID exceeding the limit it has sent MUST treat this as a connection error of type STREAM_LIMIT_ERROR ([Section 11](#)).

Once a receiver advertises a stream limit using the MAX_STREAMS frame, advertising a smaller limit has no effect. A receiver MUST ignore any MAX_STREAMS frame that does not increase the stream limit.

As with stream and connection flow control, this document leaves when and how many streams to advertise to a peer via MAX_STREAMS to implementations. Implementations might choose to increase limits as streams close to keep the number of streams available to peers roughly consistent.

An endpoint that is unable to open a new stream due to the peer's limits SHOULD send a STREAMS_BLOCKED frame ([Section 19.14](#)). This signal is considered useful for debugging. An endpoint MUST NOT wait to receive this signal before advertising additional credit, since doing so will mean that the peer will be blocked for at least an entire round trip, and potentially for longer if the peer chooses to not send STREAMS_BLOCKED frames.

5. Connections

QUIC's connection establishment combines version negotiation with the cryptographic and transport handshakes to reduce connection establishment latency, as described in [Section 7](#). Once established, a connection may migrate to a different IP or port at either endpoint as described in [Section 9](#). Finally, a connection may be terminated by either endpoint, as described in [Section 10](#).

5.1. Connection ID

Each connection possesses a set of connection identifiers, or connection IDs, each of which can identify the connection. Connection IDs are independently selected by endpoints; each endpoint selects the connection IDs that its peer uses.

The primary function of a connection ID is to ensure that changes in addressing at lower protocol layers (UDP, IP) don't cause packets for a QUIC connection to be delivered to the wrong endpoint. Each endpoint selects connection IDs using an implementation-specific (and perhaps deployment-specific) method which will allow packets with that connection ID to be routed back to the endpoint and to be identified by the endpoint upon receipt.

Connection IDs **MUST NOT** contain any information that can be used by an external observer (that is, one that does not cooperate with the issuer) to correlate them with other connection IDs for the same connection. As a trivial example, this means the same connection ID **MUST NOT** be issued more than once on the same connection.

Packets with long headers include Source Connection ID and Destination Connection ID fields. These fields are used to set the connection IDs for new connections; see [Section 7.2](#) for details.

Packets with short headers ([Section 17.3](#)) only include the Destination Connection ID and omit the explicit length. The length of the Destination Connection ID field is expected to be known to endpoints. Endpoints using a load balancer that routes based on connection ID could agree with the load balancer on a fixed length for connection IDs, or agree on an encoding scheme. A fixed portion could encode an explicit length, which allows the entire connection ID to vary in length and still be used by the load balancer.

A Version Negotiation ([Section 17.2.1](#)) packet echoes the connection IDs selected by the client, both to ensure correct routing toward the client and to allow the client to validate that the packet is in response to an Initial packet.

A zero-length connection ID can be used when a connection ID is not needed to route to the correct endpoint. However, multiplexing connections on the same local IP address and port while using zero-length connection IDs will cause failures in the presence of peer connection migration, NAT rebinding, and client port reuse; and therefore **MUST NOT** be done unless an endpoint is certain that those protocol features are not in use.

When an endpoint uses a non-zero-length connection ID, it needs to ensure that the peer has a supply of connection IDs from which to choose for packets sent to the endpoint. These connection IDs are supplied by the endpoint using the `NEW_CONNECTION_ID` frame ([Section 19.15](#)).

5.1.1. Issuing Connection IDs

Each Connection ID has an associated sequence number to assist in detecting when `NEW_CONNECTION_ID` or `RETIRE_CONNECTION_ID` frames refer to the same value. The initial connection ID issued by an endpoint is sent in the Source Connection ID field of the long packet header ([Section 17.2](#)) during the handshake. The sequence number of the initial connection ID is 0. If the `preferred_address` transport parameter is sent, the sequence number of the supplied connection ID is 1.

Additional connection IDs are communicated to the peer using `NEW_CONNECTION_ID` frames ([Section 19.15](#)). The sequence number on each newly-issued connection ID **MUST** increase by 1. The connection ID randomly selected by the client in the Initial packet and any connection ID provided by a Retry packet are not assigned sequence numbers unless a server opts to retain them as its initial connection ID.

When an endpoint issues a connection ID, it **MUST** accept packets that carry this connection ID for the duration of the connection or until its peer invalidates the connection ID via a `RETIRE_CONNECTION_ID` frame ([Section 19.16](#)). Connection IDs that are issued and not retired are considered active; any active connection ID is valid for use with the current connection at any time, in any packet type. This includes the connection ID issued by the server via the `preferred_address` transport parameter.

An endpoint **SHOULD** ensure that its peer has a sufficient number of available and unused connection IDs. Endpoints advertise the number of active connection IDs they are willing to maintain using the `active_connection_id_limit` transport parameter. An endpoint **MUST NOT** provide more connection IDs than the peer's limit. An endpoint **MAY** send connection IDs that temporarily exceed a peer's limit if the `NEW_CONNECTION_ID` frame also requires the retirement of any excess, by including a sufficiently large value in the Retire Prior To field.

A `NEW_CONNECTION_ID` frame might cause an endpoint to add some active connection IDs and retire others based on the value of the `Retire Prior To` field. After processing a `NEW_CONNECTION_ID` frame and adding and retiring active connection IDs, if the number of active connection IDs exceeds the value advertised in its `active_connection_id_limit` transport parameter, an endpoint **MUST** close the connection with an error of type `CONNECTION_ID_LIMIT_ERROR`.

An endpoint **SHOULD** supply a new connection ID when the peer retires a connection ID. If an endpoint provided fewer connection IDs than the peer's `active_connection_id_limit`, it **MAY** supply a new connection ID when it receives a packet with a previously unused connection ID. An endpoint **MAY** limit the frequency or the total number of connection IDs issued for each connection to avoid the risk of running out of connection IDs; see [Section 10.4.2](#). An endpoint **MAY** also limit the issuance of connection IDs to reduce the amount of per-path state it maintains, such as path validation status, as its peer might interact with it over as many paths as there are issued connection IDs.

An endpoint that initiates migration and requires non-zero-length connection IDs **SHOULD** ensure that the pool of connection IDs available to its peer allows the peer to use a new connection ID on migration, as the peer will close the connection if the pool is exhausted.

5.1.2. Consuming and Retiring Connection IDs

An endpoint can change the connection ID it uses for a peer to another available one at any time during the connection. An endpoint consumes connection IDs in response to a migrating peer; see [Section 9.5](#) for more.

An endpoint maintains a set of connection IDs received from its peer, any of which it can use when sending packets. When the endpoint wishes to remove a connection ID from use, it sends a `RETIRE_CONNECTION_ID` frame to its peer. Sending a `RETIRE_CONNECTION_ID` frame indicates that the connection ID will not be used again and requests that the peer replace it with a new connection ID using a `NEW_CONNECTION_ID` frame.

As discussed in [Section 9.5](#), endpoints limit the use of a connection ID to packets sent from a single local address to a single destination address. Endpoints **SHOULD** retire connection IDs when they are no longer actively using either the local or destination address for which the connection ID was used.

An endpoint might need to stop accepting previously issued connection IDs in certain circumstances. Such an endpoint can cause its peer to retire connection IDs by sending a `NEW_CONNECTION_ID` frame with an increased `Retire Prior To` field. The endpoint SHOULD continue to accept the previously issued connection IDs until they are retired by the peer. If the endpoint can no longer process the indicated connection IDs, it MAY close the connection.

Upon receipt of an increased `Retire Prior To` field, the peer MUST stop using the corresponding connection IDs and retire them with `RETIRE_CONNECTION_ID` frames before adding the newly provided connection ID to the set of active connection IDs. This ordering allows an endpoint to replace all active connection IDs without the possibility of a peer having no available connection IDs and without exceeding the limit the peer sets in the `active_connection_id_limit` transport parameter; see [Section 18.2](#). Failure to cease using the connection IDs when requested can result in connection failures, as the issuing endpoint might be unable to continue using the connection IDs with the active connection.

An endpoint SHOULD limit the number of connection IDs it has retired locally and have not yet been acknowledged. An endpoint SHOULD allow for sending and tracking a number of `RETIRE_CONNECTION_ID` frames of at least twice the `active_connection_id` limit. An endpoint MUST NOT forget a connection ID without retiring it, though it MAY choose to treat having connection IDs in need of retirement that exceed this limit as a connection error of type `CONNECTION_ID_LIMIT_ERROR`.

Endpoints SHOULD NOT issue updates of the `Retire Prior To` field before receiving `RETIRE_CONNECTION_ID` frames that retire all connection IDs indicated by the previous `Retire Prior To` value.

5.2. Matching Packets to Connections

Incoming packets are classified on receipt. Packets can either be associated with an existing connection, or - for servers - potentially create a new connection.

Endpoints try to associate a packet with an existing connection. If the packet has a non-zero-length `Destination Connection ID` corresponding to an existing connection, QUIC processes that packet accordingly. Note that more than one connection ID can be associated with a connection; see [Section 5.1](#).

If the `Destination Connection ID` is zero length and the addressing information in the packet matches the addressing information the endpoint uses to identify a connection with a zero-length connection ID, QUIC processes the packet as part of that connection. An

endpoint can use just destination IP and port or both source and destination addresses for identification, though this makes connections fragile as described in [Section 5.1](#).

Endpoints can send a Stateless Reset ([Section 10.4](#)) for any packets that cannot be attributed to an existing connection. A stateless reset allows a peer to more quickly identify when a connection becomes unusable.

Packets that are matched to an existing connection are discarded if the packets are inconsistent with the state of that connection. For example, packets are discarded if they indicate a different protocol version than that of the connection, or if the removal of packet protection is unsuccessful once the expected keys are available.

Invalid packets without packet protection, such as Initial, Retry, or Version Negotiation, MAY be discarded. An endpoint MUST generate a connection error if it commits changes to state before discovering an error.

5.2.1. Client Packet Handling

Valid packets sent to clients always include a Destination Connection ID that matches a value the client selects. Clients that choose to receive zero-length connection IDs can use the local address and port to identify a connection. Packets that don't match an existing connection are discarded.

Due to packet reordering or loss, a client might receive packets for a connection that are encrypted with a key it has not yet computed. The client MAY drop these packets, or MAY buffer them in anticipation of later packets that allow it to compute the key.

If a client receives a packet that has an unsupported version, it MUST discard that packet.

5.2.2. Server Packet Handling

If a server receives a packet that indicates an unsupported version but is large enough to initiate a new connection for any one supported version, the server SHOULD send a Version Negotiation packet as described in [Section 6.1](#). Servers MAY limit the number of packets that it responds to with a Version Negotiation packet. Servers MUST drop smaller packets that specify unsupported versions.

The first packet for an unsupported version can use different semantics and encodings for any version-specific field. In particular, different packet protection keys might be used for

different versions. Servers that do not support a particular version are unlikely to be able to decrypt the payload of the packet. Servers SHOULD NOT attempt to decode or decrypt a packet from an unknown version, but instead send a Version Negotiation packet, provided that the packet is sufficiently long.

Packets with a supported version, or no version field, are matched to a connection using the connection ID or - for packets with zero-length connection IDs - the local address and port. If the packet doesn't match an existing connection, the server continues below.

If the packet is an Initial packet fully conforming with the specification, the server proceeds with the handshake ([Section 7](#)). This commits the server to the version that the client selected.

If a server refuses to accept a new connection, it SHOULD send an Initial packet containing a CONNECTION_CLOSE frame with error code CONNECTION_REFUSED.

If the packet is a 0-RTT packet, the server MAY buffer a limited number of these packets in anticipation of a late-arriving Initial packet. Clients are not able to send Handshake packets prior to receiving a server response, so servers SHOULD ignore any such packets.

Servers MUST drop incoming packets under all other circumstances.

5.2.3. Considerations for Simple Load Balancers

A server deployment could load balance among servers using only source and destination IP addresses and ports. Changes to the client's IP address or port could result in packets being forwarded to the wrong server. Such a server deployment could use one of the following methods for connection continuity when a client's address changes.

- * Servers could use an out-of-band mechanism to forward packets to the correct server based on Connection ID.
- * If servers can use a dedicated server IP address or port, other than the one that the client initially connects to, they could use the preferred_address transport parameter to request that clients move connections to that dedicated address. Note that clients could choose not to use the preferred address.

A server in a deployment that does not implement a solution to maintain connection continuity when the client address changes SHOULD indicate migration is not supported using the

`disable_active_migration` transport parameter. The `disable_active_migration` transport parameter does not prohibit connection migration after a client has acted on a `preferred_address` transport parameter.

Server deployments that use this simple form of load balancing **MUST** avoid the creation of a stateless reset oracle; see [Section 21.9](#).

5.3. Life of a QUIC Connection

A QUIC connection is a stateful interaction between a client and server, the primary purpose of which is to support the exchange of data by an application protocol. Streams ([Section 2](#)) are the primary means by which an application protocol exchanges information.

Each connection starts with a handshake phase, during which client and server establish a shared secret using the cryptographic handshake protocol [[QUIC-TLS](#)] and negotiate the application protocol. The handshake ([Section 7](#)) confirms that both endpoints are willing to communicate ([Section 8.1](#)) and establishes parameters for the connection ([Section 7.4](#)).

An application protocol can also operate in a limited fashion during the handshake phase. 0-RTT allows application data to be sent by a client before receiving any response from the server. However, 0-RTT lacks certain key security guarantees. In particular, there is no protection against replay attacks in 0-RTT; see [[QUIC-TLS](#)]. Separately, a server can also send application data to a client before it receives the final cryptographic handshake messages that allow it to confirm the identity and liveness of the client. These capabilities allow an application protocol to offer the option to trade some security guarantees for reduced latency.

The use of connection IDs ([Section 5.1](#)) allows connections to migrate to a new network path, both as a direct choice of an endpoint and when forced by a change in a middlebox. [Section 9](#) describes mitigations for the security and privacy issues associated with migration.

For connections that are no longer needed or desired, there are several ways for a client and server to terminate a connection ([Section 10](#)).

5.4. Required Operations on Connections

There are certain operations that an application **MUST** be able to perform when interacting with the QUIC transport. This document does not specify an API, but any implementation of this version of QUIC **MUST** expose the ability to perform the operations described in this section on a QUIC connection.

When implementing the client role, applications need to be able to:

- * open a connection, which begins the exchange described in [Section 7](#);
- * enable 0-RTT when available; and
- * be informed when 0-RTT has been accepted or rejected by a server.

When implementing the server role, applications need to be able to:

- * listen for incoming connections, which prepares for the exchange described in [Section 7](#);
- * if Early Data is supported, embed application-controlled data in the TLS resumption ticket sent to the client; and
- * if Early Data is supported, retrieve application-controlled data from the client's resumption ticket and enable rejecting Early Data based on that information.

In either role, applications need to be able to:

- * configure minimum values for the initial number of permitted streams of each type, as communicated in the transport parameters ([Section 7.4](#));
- * control resource allocation of various types, including flow control and the number of permitted streams of each type;
- * identify whether the handshake has completed successfully or is still ongoing;
- * keep a connection from silently closing, either by generating PING frames ([Section 19.2](#)) or by requesting that the transport send additional frames before the idle timeout expires ([Section 10.2](#)); and
- * immediately close ([Section 10.3](#)) the connection.

6. Version Negotiation

Version negotiation ensures that client and server agree to a QUIC version that is mutually supported. A server sends a Version Negotiation packet in response to each packet that might initiate a new connection; see [Section 5.2](#) for details.

The size of the first packet sent by a client will determine whether a server sends a Version Negotiation packet. Clients that support multiple QUIC versions SHOULD pad the first packet they send to the largest of the minimum packet sizes across all versions they support. This ensures that the server responds if there is a mutually supported version.

6.1. Sending Version Negotiation Packets

If the version selected by the client is not acceptable to the server, the server responds with a Version Negotiation packet; see [Section 17.2.1](#). This includes a list of versions that the server will accept. An endpoint MUST NOT send a Version Negotiation packet in response to receiving a Version Negotiation packet.

This system allows a server to process packets with unsupported versions without retaining state. Though either the Initial packet or the Version Negotiation packet that is sent in response could be lost, the client will send new packets until it successfully receives a response or it abandons the connection attempt. As a result, the client discards all state for the connection and does not send any more packets on the connection.

A server MAY limit the number of Version Negotiation packets it sends. For instance, a server that is able to recognize packets as 0-RTT might choose not to send Version Negotiation packets in response to 0-RTT packets with the expectation that it will eventually receive an Initial packet.

6.2. Handling Version Negotiation Packets

Version Negotiation packets are designed to allow future versions of QUIC to negotiate the version in use between endpoints. Future versions of QUIC might change how implementations that support multiple versions of QUIC react to Version Negotiation packets when attempting to establish a connection using this version.

A client that supports only this version of QUIC MUST abandon the current connection attempt if it receives a Version Negotiation packet, with the following two exceptions. A client MUST discard any Version Negotiation packet if it has received and successfully

processed any other packet, including an earlier Version Negotiation packet. A client **MUST** discard a Version Negotiation packet that lists the QUIC version selected by the client.

How to perform version negotiation is left as future work defined by future versions of QUIC. In particular, that future work will ensure robustness against version downgrade attacks; see [Section 21.10](#).

6.2.1. Version Negotiation Between Draft Versions

[[RFC editor: please remove this section before publication.]]

When a draft implementation receives a Version Negotiation packet, it **MAY** use it to attempt a new connection with one of the versions listed in the packet, instead of abandoning the current connection attempt; see [Section 6.2](#).

The client **MUST** check that the Destination and Source Connection ID fields match the Source and Destination Connection ID fields in a packet that the client sent. If this check fails, the packet **MUST** be discarded.

Once the Version Negotiation packet is determined to be valid, the client then selects an acceptable protocol version from the list provided by the server. The client then attempts to create a new connection using that version. The new connection **MUST** use a new random Destination Connection ID different from the one it had previously sent.

Note that this mechanism does not protect against downgrade attacks and **MUST NOT** be used outside of draft implementations.

6.3. Using Reserved Versions

For a server to use a new version in the future, clients need to correctly handle unsupported versions. Some version numbers (0x?a?a?a as defined in [Section 15](#)) are reserved for inclusion in fields that contain version numbers.

Endpoints **MAY** add reserved versions to any field where unknown or unsupported versions are ignored to test that a peer correctly ignores the value. For instance, an endpoint could include a reserved version in a Version Negotiation packet; see [Section 17.2.1](#). Endpoints **MAY** send packets with a reserved version to test that a peer correctly discards the packet.

7. Cryptographic and Transport Handshake

QUIC relies on a combined cryptographic and transport handshake to minimize connection establishment latency. QUIC uses the CRYPTO frame [Section 19.6](#) to transmit the cryptographic handshake. Version 0x00000001 of QUIC uses TLS as described in [[QUIC-TLS](#)]; a different QUIC version number could indicate that a different cryptographic handshake protocol is in use.

QUIC provides reliable, ordered delivery of the cryptographic handshake data. QUIC packet protection is used to encrypt as much of the handshake protocol as possible. The cryptographic handshake **MUST** provide the following properties:

- * authenticated key exchange, where
 - a server is always authenticated,
 - a client is optionally authenticated,
 - every connection produces distinct and unrelated keys,
 - keying material is usable for packet protection for both 0-RTT and 1-RTT packets, and
 - 1-RTT keys have forward secrecy
- * authenticated values for transport parameters of both endpoints, and confidentiality protection for server transport parameters (see [Section 7.4](#))
- * authenticated negotiation of an application protocol (TLS uses ALPN [[RFC7301](#)] for this purpose)

An endpoint can verify support for Explicit Congestion Notification (ECN) in the first packets it sends, as described in [Section 13.4.2](#).

The CRYPTO frame can be sent in different packet number spaces ([Section 12.3](#)). The offsets used by CRYPTO frames to ensure ordered delivery of cryptographic handshake data start from zero in each packet number space.

Endpoints **MUST** explicitly negotiate an application protocol. This avoids situations where there is a disagreement about the protocol that is in use.

7.1. Example Handshake Flows

Details of how TLS is integrated with QUIC are provided in [QUIC-TLS], but some examples are provided here. An extension of this exchange to support client address validation is shown in Section 8.1.2.

Once any address validation exchanges are complete, the cryptographic handshake is used to agree on cryptographic keys. The cryptographic handshake is carried in Initial (Section 17.2.2) and Handshake (Section 17.2.4) packets.

Figure 4 provides an overview of the 1-RTT handshake. Each line shows a QUIC packet with the packet type and packet number shown first, followed by the frames that are typically contained in those packets. So, for instance the first packet is of type Initial, with packet number 0, and contains a CRYPTO frame carrying the ClientHello.

Multiple QUIC packets - even of different packet types - can be coalesced into a single UDP datagram; see Section 12.2. As a result, this handshake may consist of as few as 4 UDP datagrams, or any number more (subject to limits inherent to the protocol, such as congestion control or anti-amplification). For instance, the server's first flight contains Initial packets, Handshake packets, and "0.5-RTT data" in 1-RTT packets with a short header.

Client	Server
Initial[0]: CRYPTO[CH] ->	
	Initial[0]: CRYPTO[SH] ACK[0]
	Handshake[0]: CRYPTO[EE, CERT, CV, FIN]
	<- 1-RTT[0]: STREAM[1, "..."]
Initial[1]: ACK[0]	
Handshake[0]: CRYPTO[FIN], ACK[0]	
1-RTT[0]: STREAM[0, "..."], ACK[0] ->	
	Handshake[1]: ACK[0]
	<- 1-RTT[1]: STREAM[3, "..."], ACK[0]

Figure 4: Example 1-RTT Handshake

Figure 5 shows an example of a connection with a 0-RTT handshake and a single packet of 0-RTT data. Note that as described in Section 12.3, the server acknowledges 0-RTT data in 1-RTT packets, and the client sends 1-RTT packets in the same packet number space.

Client	Server
Initial[0]: CRYPTO[CH]	
0-RTT[0]: STREAM[0, "..."] ->	
	Initial[0]: CRYPTO[SH] ACK[0]
	Handshake[0] CRYPTO[EE, FIN]
	<- 1-RTT[0]: STREAM[1, "..."] ACK[0]
Initial[1]: ACK[0]	
Handshake[0]: CRYPTO[FIN], ACK[0]	
1-RTT[1]: STREAM[0, "..."] ACK[0] ->	
	Handshake[1]: ACK[0]
	<- 1-RTT[1]: STREAM[3, "..."], ACK[1]

Figure 5: Example 0-RTT Handshake

7.2. Negotiating Connection IDs

A connection ID is used to ensure consistent routing of packets, as described in [Section 5.1](#). The long header contains two connection IDs: the Destination Connection ID is chosen by the recipient of the packet and is used to provide consistent routing; the Source Connection ID is used to set the Destination Connection ID used by the peer.

During the handshake, packets with the long header ([Section 17.2](#)) are used to establish the connection IDs in each direction. Each endpoint uses the Source Connection ID field to specify the connection ID that is used in the Destination Connection ID field of packets being sent to them. Upon receiving a packet, each endpoint sets the Destination Connection ID it sends to match the value of the Source Connection ID that it receives.

When an Initial packet is sent by a client that has not previously received an Initial or Retry packet from the server, the client populates the Destination Connection ID field with an unpredictable value. This Destination Connection ID MUST be at least 8 bytes in length. Until a packet is received from the server, the client MUST use the same Destination Connection ID value on all packets in this connection. This Destination Connection ID is used to determine packet protection keys for Initial packets.

The client populates the Source Connection ID field with a value of its choosing and sets the Source Connection ID Length field to indicate the length.

The first flight of 0-RTT packets use the same Destination Connection ID and Source Connection ID values as the client's first Initial packet.

Upon first receiving an Initial or Retry packet from the server, the client uses the Source Connection ID supplied by the server as the Destination Connection ID for subsequent packets, including any 0-RTT packets. This means that a client might have to change the connection ID it sets in the Destination Connection ID field twice during connection establishment: once in response to a Retry, and once in response to an Initial packet from the server. Once a client has received a valid Initial packet from the server, it **MUST** discard any subsequent packet it receives with a different Source Connection ID.

A client **MUST** change the Destination Connection ID it uses for sending packets in response to only the first received Initial or Retry packet. A server **MUST** set the Destination Connection ID it uses for sending packets based on the first received Initial packet. Any further changes to the Destination Connection ID are only permitted if the values are taken from any received `NEW_CONNECTION_ID` frames; if subsequent Initial packets include a different Source Connection ID, they **MUST** be discarded. This avoids unpredictable outcomes that might otherwise result from stateless processing of multiple Initial packets with different Source Connection IDs.

The Destination Connection ID that an endpoint sends can change over the lifetime of a connection, especially in response to connection migration ([Section 9](#)); see [Section 5.1.1](#) for details.

7.3. Authenticating Connection IDs

The choice each endpoint makes about connection IDs during the handshake is authenticated by including all values in transport parameters; see [Section 7.4](#). This ensures that all connection IDs used for the handshake are also authenticated by the cryptographic handshake.

Each endpoint includes the value of the Source Connection ID field from the first Initial packet it sent in the `initial_source_connection_id` transport parameter; see [Section 18.2](#). A server includes the Destination Connection ID field from the first Initial packet it received from the client in the `original_destination_connection_id` transport parameter; if the server sent a Retry packet this refers to the first Initial packet received before sending the Retry packet. If it sends a Retry packet, a server also includes the Source Connection ID field from the Retry packet in the `retry_source_connection_id` transport parameter.

The values provided by a peer for these transport parameters MUST match the values that an endpoint used in the Destination and Source Connection ID fields of Initial packets that it sent. Including connection ID values in transport parameters and verifying them ensures that that an attacker cannot influence the choice of connection ID for a successful connection by injecting packets carrying attacker-chosen connection IDs during the handshake.

An endpoint MUST treat absence of the `initial_source_connection_id` transport parameter from either endpoint or absence of the `original_destination_connection_id` transport parameter from the server as a connection error of type `TRANSPORT_PARAMETER_ERROR`.

An endpoint MUST treat the following as a connection error of type `TRANSPORT_PARAMETER_ERROR` or `PROTOCOL_VIOLATION`:

- * absence of the `retry_source_connection_id` transport parameter from the server after receiving a Retry packet,
- * presence of the `retry_source_connection_id` transport parameter when no Retry packet was received, or
- * a mismatch between values received from a peer in these transport parameters and the value sent in the corresponding Destination or Source Connection ID fields of Initial packets.

If a zero-length connection ID is selected, the corresponding transport parameter is included with a zero-length value.

Figure 6 shows the connection IDs (with DCID=Destination Connection ID, SCID=Source Connection ID) that are used in a complete handshake. The exchange of Initial packets is shown, plus the later exchange of 1-RTT packets that includes the connection ID established during the handshake.

Client	Server
Initial: DCID=S1, SCID=C1 ->	<- Initial: DCID=C1, SCID=S3
	...
1-RTT: DCID=S3 ->	<- 1-RTT: DCID=C1

Figure 6: Use of Connection IDs in a Handshake

Figure 7 shows a similar handshake that includes a Retry packet.

Client	Server
Initial: DCID=S1, SCID=C1 ->	
	<- Retry: DCID=C1, SCID=S2
Initial: DCID=S2, SCID=C1 ->	<- Initial: DCID=C1, SCID=S3
	...
1-RTT: DCID=S3 ->	<- 1-RTT: DCID=C1

Figure 7: Use of Connection IDs in a Handshake with Retry

In both cases (Figure 6 and Figure 7), the client sets the value of the `initial_source_connection_id` transport parameter to "C1".

When the handshake does not include a Retry (Figure 6), the server sets `original_destination_connection_id` to "S1" and `initial_source_connection_id` to "S3". In this case, the server does not include a `retry_source_connection_id` transport parameter.

When the handshake includes a Retry (Figure 7), the server sets `original_destination_connection_id` to "S1", `retry_source_connection_id` to "S2", and `initial_source_connection_id` to "S3".

Each endpoint validates transport parameters set by the peer. The client confirms that the `retry_source_connection_id` transport parameter is absent if it did not process a Retry packet.

7.4. Transport Parameters

During connection establishment, both endpoints make authenticated declarations of their transport parameters. Endpoints are required to comply with the restrictions that each parameter defines; the description of each parameter includes rules for its handling.

Transport parameters are declarations that are made unilaterally by each endpoint. Each endpoint can choose values for transport parameters independent of the values chosen by its peer.

The encoding of the transport parameters is detailed in [Section 18](#).

QUIC includes the encoded transport parameters in the cryptographic handshake. Once the handshake completes, the transport parameters declared by the peer are available. Each endpoint validates the value provided by its peer.

Definitions for each of the defined transport parameters are included in [Section 18.2](#).

An endpoint MUST treat receipt of a transport parameter with an invalid value as a connection error of type `TRANSPORT_PARAMETER_ERROR`.

An endpoint MUST NOT send a parameter more than once in a given transport parameters extension. An endpoint SHOULD treat receipt of duplicate transport parameters as a connection error of type `TRANSPORT_PARAMETER_ERROR`.

Endpoints use transport parameters to authenticate the negotiation of connection IDs during the handshake; see [Section 7.3](#).

7.4.1. Values of Transport Parameters for 0-RTT

Both endpoints store the value of the server transport parameters from a connection and apply them to any 0-RTT packets that are sent in subsequent connections to that peer, except for transport parameters that are explicitly excluded. Remembered transport parameters apply to the new connection until the handshake completes and the client starts sending 1-RTT packets. Once the handshake completes, the client uses the transport parameters established in the handshake.

The definition of new transport parameters ([Section 7.4.2](#)) MUST specify whether they MUST, MAY, or MUST NOT be stored for 0-RTT. A client need not store a transport parameter it cannot process.

A client MUST NOT use remembered values for the following parameters: `ack_delay_exponent`, `max_ack_delay`, `initial_source_connection_id`, `original_destination_connection_id`, `preferred_address`, `retry_source_connection_id`, and `stateless_reset_token`. The client MUST use the server's new values in the handshake instead, and absent new values from the server, the default value.

A client that attempts to send 0-RTT data MUST remember all other transport parameters used by the server. The server can remember these transport parameters, or store an integrity-protected copy of the values in the ticket and recover the information when accepting 0-RTT data. A server uses the transport parameters in determining whether to accept 0-RTT data.

If 0-RTT data is accepted by the server, the server MUST NOT reduce any limits or alter any values that might be violated by the client with its 0-RTT data. In particular, a server that accepts 0-RTT data MUST NOT set values for the following parameters ([Section 18.2](#)) that are smaller than the remembered value of the parameters.

- * `active_connection_id_limit`
- * `initial_max_data`
- * `initial_max_stream_data_bidi_local`
- * `initial_max_stream_data_bidi_remote`
- * `initial_max_stream_data_uni`
- * `initial_max_streams_bidi`
- * `initial_max_streams_uni`

Omitting or setting a zero value for certain transport parameters can result in 0-RTT data being enabled, but not usable. The applicable subset of transport parameters that permit sending of application data SHOULD be set to non-zero values for 0-RTT. This includes `initial_max_data` and either `initial_max_streams_bidi` and `initial_max_stream_data_bidi_remote`, or `initial_max_streams_uni` and `initial_max_stream_data_uni`.

A server MUST either reject 0-RTT data or abort a handshake if the implied values for transport parameters cannot be supported.

When sending frames in 0-RTT packets, a client MUST only use remembered transport parameters; importantly, it MUST NOT use updated values that it learns from the server's updated transport parameters or from frames received in 1-RTT packets. Updated values of transport parameters from the handshake apply only to 1-RTT packets. For instance, flow control limits from remembered transport parameters apply to all 0-RTT packets even if those values are increased by the handshake or by frames sent in 1-RTT packets. A server MAY treat use of updated transport parameters in 0-RTT as a connection error of type `PROTOCOL_VIOLATION`.

7.4.2. New Transport Parameters

New transport parameters can be used to negotiate new protocol behavior. An endpoint **MUST** ignore transport parameters that it does not support. Absence of a transport parameter therefore disables any optional protocol feature that is negotiated using the parameter. As described in [Section 18.1](#), some identifiers are reserved in order to exercise this requirement.

New transport parameters can be registered according to the rules in [Section 22.2](#).

7.5. Cryptographic Message Buffering

Implementations need to maintain a buffer of CRYPTO data received out of order. Because there is no flow control of CRYPTO frames, an endpoint could potentially force its peer to buffer an unbounded amount of data.

Implementations **MUST** support buffering at least 4096 bytes of data received in out of order CRYPTO frames. Endpoints **MAY** choose to allow more data to be buffered during the handshake. A larger limit during the handshake could allow for larger keys or credentials to be exchanged. An endpoint's buffer size does not need to remain constant during the life of the connection.

Being unable to buffer CRYPTO frames during the handshake can lead to a connection failure. If an endpoint's buffer is exceeded during the handshake, it can expand its buffer temporarily to complete the handshake. If an endpoint does not expand its buffer, it **MUST** close the connection with a CRYPTO_BUFFER_EXCEEDED error code.

Once the handshake completes, if an endpoint is unable to buffer all data in a CRYPTO frame, it **MAY** discard that CRYPTO frame and all CRYPTO frames received in the future, or it **MAY** close the connection with a CRYPTO_BUFFER_EXCEEDED error code. Packets containing discarded CRYPTO frames **MUST** be acknowledged because the packet has been received and processed by the transport even though the CRYPTO frame was discarded.

8. Address Validation

Address validation ensures that an endpoint cannot be used for a traffic amplification attack. In such an attack, a packet is sent to a server with spoofed source address information that identifies a victim. If a server generates more or larger packets in response to that packet, the attacker can use the server to send more data toward the victim than it would be able to send on its own.

The primary defense against amplification attack is verifying that an endpoint is able to receive packets at the transport address that it claims. Address validation is performed both during connection establishment (see [Section 8.1](#)) and during connection migration (see [Section 8.2](#)).

8.1. Address Validation During Connection Establishment

Connection establishment implicitly provides address validation for both endpoints. In particular, receipt of a packet protected with Handshake keys confirms that the client received the Initial packet from the server. Once the server has successfully processed a Handshake packet from the client, it can consider the client address to have been validated.

Prior to validating the client address, servers **MUST NOT** send more than three times as many bytes as the number of bytes they have received. This limits the magnitude of any amplification attack that can be mounted using spoofed source addresses. For the purposes of avoiding amplification prior to address validation, servers **MUST** count all of the payload bytes received in datagrams that are uniquely attributed to a single connection. This includes datagrams that contain packets that are successfully processed and datagrams that contain packets that are all discarded.

Clients **MUST** ensure that UDP datagrams containing Initial packets have UDP payloads of at least 1200 bytes, adding padding to packets in the datagram as necessary. A client that sends padded datagrams allows the server to send more data prior to completing address validation.

Loss of an Initial or Handshake packet from the server can cause a deadlock if the client does not send additional Initial or Handshake packets. A deadlock could occur when the server reaches its anti-amplification limit and the client has received acknowledgements for all the data it has sent. In this case, when the client has no reason to send additional packets, the server will be unable to send more data because it has not validated the client's address. To prevent this deadlock, clients **MUST** send a packet on a probe timeout (PTO, see Section 6.2 of [\[QUIC-RECOVERY\]](#)). Specifically, the client **MUST** send an Initial packet in a UDP datagram that contains at least 1200 bytes if it does not have Handshake keys, and otherwise send a Handshake packet.

A server might wish to validate the client address before starting the cryptographic handshake. QUIC uses a token in the Initial packet to provide address validation prior to completing the handshake. This token is delivered to the client during connection establishment with a Retry packet (see [Section 8.1.2](#)) or in a previous connection using the NEW_TOKEN frame (see [Section 8.1.3](#)).

In addition to sending limits imposed prior to address validation, servers are also constrained in what they can send by the limits set by the congestion controller. Clients are only constrained by the congestion controller.

8.1.1. Token Construction

A token sent in a NEW_TOKEN frames or a Retry packet MUST be constructed in a way that allows the server to identify how it was provided to a client. These tokens are carried in the same field, but require different handling from servers.

8.1.2. Address Validation using Retry Packets

Upon receiving the client's Initial packet, the server can request address validation by sending a Retry packet ([Section 17.2.5](#)) containing a token. This token MUST be repeated by the client in all Initial packets it sends for that connection after it receives the Retry packet. In response to processing an Initial containing a token, a server can either abort the connection or permit it to proceed.

As long as it is not possible for an attacker to generate a valid token for its own address (see [Section 8.1.4](#)) and the client is able to return that token, it proves to the server that it received the token.

A server can also use a Retry packet to defer the state and processing costs of connection establishment. Requiring the server to provide a different connection ID, along with the `original_destination_connection_id` transport parameter defined in [Section 18.2](#), forces the server to demonstrate that it, or an entity it cooperates with, received the original Initial packet from the client. Providing a different connection ID also grants a server some control over how subsequent packets are routed. This can be used to direct connections to a different server instance.

If a server receives a client Initial that can be unprotected but contains an invalid Retry token, it knows the client will not accept another Retry token. The server can discard such a packet and allow the client to time out to detect handshake failure, but that could

impose a significant latency penalty on the client. Instead, the server SHOULD immediately close ([Section 10.3](#)) the connection with an `INVALID_TOKEN` error. Note that a server has not established any state for the connection at this point and so does not enter the closing period.

A flow showing the use of a Retry packet is shown in Figure 8.

```

Client                                                    Server

Initial[0]: CRYPTO[CH] ->

                                                    <- Retry+Token

Initial+Token[1]: CRYPTO[CH] ->

                                                    Initial[0]: CRYPTO[SH] ACK[1]
                                                    Handshake[0]: CRYPTO[EE, CERT, CV, FIN]
                                                    <- 1-RTT[0]: STREAM[1, "..."]

```

Figure 8: Example Handshake with Retry

8.1.3. Address Validation for Future Connections

A server MAY provide clients with an address validation token during one connection that can be used on a subsequent connection. Address validation is especially important with 0-RTT because a server potentially sends a significant amount of data to a client in response to 0-RTT data.

The server uses the `NEW_TOKEN` frame [Section 19.7](#) to provide the client with an address validation token that can be used to validate future connections. The client includes this token in Initial packets to provide address validation in a future connection. The client MUST include the token in all Initial packets it sends, unless a Retry replaces the token with a newer one. The client MUST NOT use the token provided in a Retry for future connections. Servers MAY discard any Initial packet that does not carry the expected token.

Unlike the token that is created for a Retry packet, which is used immediately, the token sent in the `NEW_TOKEN` frame might be used after some period of time has passed. Thus, a token SHOULD have an expiration time, which could be either an explicit expiration time or an issued timestamp that can be used to dynamically calculate the expiration time. A server can store the expiration time or include it in an encrypted form in the token.

A token issued with `NEW_TOKEN` MUST NOT include information that would allow values to be linked by an observer to the connection on which it was issued, unless the values are encrypted. For example, it cannot include the previous connection ID or addressing information. A server MUST ensure that every `NEW_TOKEN` frame it sends is unique across all clients, with the exception of those sent to repair losses of previously sent `NEW_TOKEN` frames. Information that allows the server to distinguish between tokens from `Retry` and `NEW_TOKEN` MAY be accessible to entities other than the server.

It is unlikely that the client port number is the same on two different connections; validating the port is therefore unlikely to be successful.

A token received in a `NEW_TOKEN` frame is applicable to any server that the connection is considered authoritative for (e.g., server names included in the certificate). When connecting to a server for which the client retains an applicable and unused token, it SHOULD include that token in the `Token` field of its Initial packet. Including a token might allow the server to validate the client address without an additional round trip. A client MUST NOT include a token that is not applicable to the server that it is connecting to, unless the client has the knowledge that the server that issued the token and the server the client is connecting to are jointly managing the tokens. A client MAY use a token from any previous connection to that server.

A token allows a server to correlate activity between the connection where the token was issued and any connection where it is used. Clients that want to break continuity of identity with a server MAY discard tokens provided using the `NEW_TOKEN` frame. In comparison, a token obtained in a `Retry` packet MUST be used immediately during the connection attempt and cannot be used in subsequent connection attempts.

A client SHOULD NOT reuse a `NEW_TOKEN` token for different connection attempts. Reusing a token allows connections to be linked by entities on the network path; see [Section 9.5](#).

Clients might receive multiple tokens on a single connection. Aside from preventing linkability, any token can be used in any connection attempt. Servers can send additional tokens to either enable address validation for multiple connection attempts or to replace older tokens that might become invalid. For a client, this ambiguity means that sending the most recent unused token is most likely to be effective. Though saving and using older tokens has no negative consequences, clients can regard older tokens as being less likely to be useful to the server for address validation.

When a server receives an Initial packet with an address validation token, it **MUST** attempt to validate the token, unless it has already completed address validation. If the token is invalid then the server **SHOULD** proceed as if the client did not have a validated address, including potentially sending a Retry. A server **SHOULD** encode tokens provided with `NEW_TOKEN` frames and Retry packets differently, and validate the latter more strictly. If the validation succeeds, the server **SHOULD** then allow the handshake to proceed.

Note: The rationale for treating the client as unvalidated rather than discarding the packet is that the client might have received the token in a previous connection using the `NEW_TOKEN` frame, and if the server has lost state, it might be unable to validate the token at all, leading to connection failure if the packet is discarded.

In a stateless design, a server can use encrypted and authenticated tokens to pass information to clients that the server can later recover and use to validate a client address. Tokens are not integrated into the cryptographic handshake and so they are not authenticated. For instance, a client might be able to reuse a token. To avoid attacks that exploit this property, a server can limit its use of tokens to only the information needed to validate client addresses.

Clients **MAY** use tokens obtained on one connection for any connection attempt using the same version. When selecting a token to use, clients do not need to consider other properties of the connection that is being attempted, including the choice of possible application protocols, session tickets, or other connection properties.

Attackers could replay tokens to use servers as amplifiers in DDoS attacks. To protect against such attacks, servers **SHOULD** ensure that tokens sent in Retry packets are only accepted for a short time. Tokens that are provided in `NEW_TOKEN` frames ([Section 19.7](#)) need to be valid for longer, but **SHOULD NOT** be accepted multiple times in a short period. Servers are encouraged to allow tokens to be used only once, if possible.

8.1.4. Address Validation Token Integrity

An address validation token **MUST** be difficult to guess. Including a large enough random value in the token would be sufficient, but this depends on the server remembering the value it sends to clients.

A token-based scheme allows the server to offload any state associated with validation to the client. For this design to work, the token **MUST** be covered by integrity protection against modification or falsification by clients. Without integrity protection, malicious clients could generate or guess values for tokens that would be accepted by the server. Only the server requires access to the integrity protection key for tokens.

There is no need for a single well-defined format for the token because the server that generates the token also consumes it. Tokens sent in Retry packets **SHOULD** include information that allows the server to verify that the source IP address and port in client packets remain constant.

Tokens sent in `NEW_TOKEN` frames **MUST** include information that allows the server to verify that the client IP address has not changed from when the token was issued. Servers can use tokens from `NEW_TOKEN` in deciding not to send a Retry packet, even if the client address has changed. If the client IP address has changed, the server **MUST** adhere to the anti-amplification limits found in [Section 8.1](#). Note that in the presence of NAT, this requirement might be insufficient to protect other hosts that share the NAT from amplification attack.

Servers **MUST** ensure that replay of tokens is prevented or limited. For instance, servers might limit the time over which a token is accepted. Tokens provided in `NEW_TOKEN` frames might need to allow longer validity periods. Tokens **MAY** include additional information about clients to further narrow applicability or reuse.

8.2. Path Validation

Path validation is used during connection migration (see [Section 9](#) and [Section 9.6](#)) by the migrating endpoint to verify reachability of a peer from a new local address. In path validation, endpoints test reachability between a specific local address and a specific peer address, where an address is the two-tuple of IP address and port.

Path validation tests that packets (`PATH_CHALLENGE`) can be both sent to and received (`PATH_RESPONSE`) from a peer on the path. Importantly, it validates that the packets received from the migrating endpoint do not carry a spoofed source address.

Path validation can be used at any time by either endpoint. For instance, an endpoint might check that a peer is still in possession of its address after a period of quiescence.

Path validation is not designed as a NAT traversal mechanism. Though the mechanism described here might be effective for the creation of NAT bindings that support NAT traversal, the expectation is that one or other peer is able to receive packets without first having sent a packet on that path. Effective NAT traversal needs additional synchronization mechanisms that are not provided here.

An endpoint MAY bundle `PATH_CHALLENGE` and `PATH_RESPONSE` frames that are used for path validation with other frames. In particular, an endpoint may pad a packet carrying a `PATH_CHALLENGE` for PMTU discovery, or an endpoint may bundle a `PATH_RESPONSE` with its own `PATH_CHALLENGE`.

When probing a new path, an endpoint might want to ensure that its peer has an unused connection ID available for responses. The endpoint can send `NEW_CONNECTION_ID` and `PATH_CHALLENGE` frames in the same packet. This ensures that an unused connection ID will be available to the peer when sending a response.

8.3. Initiating Path Validation

To initiate path validation, an endpoint sends a `PATH_CHALLENGE` frame containing a random payload on the path to be validated.

An endpoint MAY send multiple `PATH_CHALLENGE` frames to guard against packet loss. However, an endpoint SHOULD NOT send multiple `PATH_CHALLENGE` frames in a single packet. An endpoint SHOULD NOT send a `PATH_CHALLENGE` more frequently than it would an Initial packet, ensuring that connection migration is no more load on a new path than establishing a new connection.

The endpoint MUST use unpredictable data in every `PATH_CHALLENGE` frame so that it can associate the peer's response with the corresponding `PATH_CHALLENGE`.

8.4. Path Validation Responses

On receiving a `PATH_CHALLENGE` frame, an endpoint MUST respond immediately by echoing the data contained in the `PATH_CHALLENGE` frame in a `PATH_RESPONSE` frame.

An endpoint MUST NOT send more than one `PATH_RESPONSE` frame in response to one `PATH_CHALLENGE` frame; see [Section 13.3](#). The peer is expected to send more `PATH_CHALLENGE` frames as necessary to evoke additional `PATH_RESPONSE` frames.

8.5. Successful Path Validation

A new address is considered valid when a `PATH_RESPONSE` frame is received that contains the data that was sent in a previous `PATH_CHALLENGE`. Receipt of an acknowledgment for a packet containing a `PATH_CHALLENGE` frame is not adequate validation, since the acknowledgment can be spoofed by a malicious peer.

Note that receipt on a different local address does not result in path validation failure, as it might be a result of a forwarded packet (see [Section 9.3.3](#)) or misrouting. It is possible that a valid `PATH_RESPONSE` might be received in the future.

8.6. Failed Path Validation

Path validation only fails when the endpoint attempting to validate the path abandons its attempt to validate the path.

Endpoints SHOULD abandon path validation based on a timer. When setting this timer, implementations are cautioned that the new path could have a longer round-trip time than the original. A value of three times the larger of the current Probe Timeout (PTO) or the initial timeout (that is, $2 * kInitialRtt$) as defined in [\[QUIC-RECOVERY\]](#) is RECOMMENDED. That is:

$$\text{validation_timeout} = \max(3 * \text{PTO}, 6 * kInitialRtt)$$

Note that the endpoint might receive packets containing other frames on the new path, but a `PATH_RESPONSE` frame with appropriate data is required for path validation to succeed.

When an endpoint abandons path validation, it determines that the path is unusable. This does not necessarily imply a failure of the connection - endpoints can continue sending packets over other paths as appropriate. If no paths are available, an endpoint can wait for a new path to become available or close the connection.

A path validation might be abandoned for other reasons besides failure. Primarily, this happens if a connection migration to a new path is initiated while a path validation on the old path is in progress.

9. Connection Migration

The use of a connection ID allows connections to survive changes to endpoint addresses (IP address and port), such as those caused by an endpoint migrating to a new network. This section describes the process by which an endpoint migrates to a new address.

The design of QUIC relies on endpoints retaining a stable address for the duration of the handshake. An endpoint **MUST NOT** initiate connection migration before the handshake is confirmed, as defined in section 4.1.2 of [QUIC-TLS].

If the peer sent the `disable_active_migration` transport parameter, an endpoint also **MUST NOT** send packets (including probing packets; see [Section 9.1](#)) from a different local address to the address the peer used during the handshake. An endpoint that has sent this transport parameter, but detects that a peer has nonetheless migrated to a different remote address **MUST** either drop the incoming packets on that path without generating a stateless reset or proceed with path validation and allow the peer to migrate. Generating a stateless reset or closing the connection would allow third parties in the network to cause connections to close by spoofing or otherwise manipulating observed traffic.

Not all changes of peer address are intentional, or active, migrations. The peer could experience NAT rebinding: a change of address due to a middlebox, usually a NAT, allocating a new outgoing port or even a new outgoing IP address for a flow. An endpoint **MUST** perform path validation ([Section 8.2](#)) if it detects any change to a peer's address, unless it has previously validated that address.

When an endpoint has no validated path on which to send packets, it **MAY** discard connection state. An endpoint capable of connection migration **MAY** wait for a new path to become available before discarding connection state.

This document limits migration of connections to new client addresses, except as described in [Section 9.6](#). Clients are responsible for initiating all migrations. Servers do not send non-probing packets (see [Section 9.1](#)) toward a client address until they see a non-probing packet from that address. If a client receives packets from an unknown server address, the client **MUST** discard these packets.

9.1. Probing a New Path

An endpoint **MAY** probe for peer reachability from a new local address using path validation [Section 8.2](#) prior to migrating the connection to the new local address. Failure of path validation simply means that the new path is not usable for this connection. Failure to validate a path does not cause the connection to end unless there are no valid alternative paths available.

An endpoint uses a new connection ID for probes sent from a new local address; see [Section 9.5](#) for further discussion. An endpoint that uses a new local address needs to ensure that at least one new connection ID is available at the peer. That can be achieved by including a `NEW_CONNECTION_ID` frame in the probe.

Receiving a `PATH_CHALLENGE` frame from a peer indicates that the peer is probing for reachability on a path. An endpoint sends a `PATH_RESPONSE` in response as per [Section 8.2](#).

`PATH_CHALLENGE`, `PATH_RESPONSE`, `NEW_CONNECTION_ID`, and `PADDING` frames are "probing frames", and all other frames are "non-probing frames". A packet containing only probing frames is a "probing packet", and a packet containing any other frame is a "non-probing packet".

9.2. Initiating Connection Migration

An endpoint can migrate a connection to a new local address by sending packets containing non-probing frames from that address.

Each endpoint validates its peer's address during connection establishment. Therefore, a migrating endpoint can send to its peer knowing that the peer is willing to receive at the peer's current address. Thus an endpoint can migrate to a new local address without first validating the peer's address.

When migrating, the new path might not support the endpoint's current sending rate. Therefore, the endpoint resets its congestion controller, as described in [Section 9.4](#).

The new path might not have the same ECN capability. Therefore, the endpoint verifies ECN capability as described in [Section 13.4](#).

Receiving acknowledgments for data sent on the new path serves as proof of the peer's reachability from the new address. Note that since acknowledgments may be received on any path, return reachability on the new path is not established. To establish return reachability on the new path, an endpoint MAY concurrently initiate path validation [Section 8.2](#) on the new path or it MAY choose to wait for the peer to send the next non-probing frame to its new address.

9.3. Responding to Connection Migration

Receiving a packet from a new peer address containing a non-probing frame indicates that the peer has migrated to that address.

In response to such a packet, an endpoint MUST start sending subsequent packets to the new peer address and MUST initiate path validation ([Section 8.2](#)) to verify the peer's ownership of the unvalidated address.

An endpoint MAY send data to an unvalidated peer address, but it MUST protect against potential attacks as described in [Section 9.3.1](#) and [Section 9.3.2](#). An endpoint MAY skip validation of a peer address if that address has been seen recently. In particular, if an endpoint returns to a previously-validated path after detecting some form of spurious migration, skipping address validation and restoring loss detection and congestion state can reduce the performance impact of the attack.

An endpoint only changes the address that it sends packets to in response to the highest-numbered non-probing packet. This ensures that an endpoint does not send packets to an old peer address in the case that it receives reordered packets.

After changing the address to which it sends non-probing packets, an endpoint could abandon any path validation for other addresses.

Receiving a packet from a new peer address might be the result of a NAT rebinding at the peer.

After verifying a new client address, the server SHOULD send new address validation tokens ([Section 8](#)) to the client.

9.3.1. Peer Address Spoofing

It is possible that a peer is spoofing its source address to cause an endpoint to send excessive amounts of data to an unwilling host. If the endpoint sends significantly more data than the spoofing peer, connection migration might be used to amplify the volume of data that an attacker can generate toward a victim.

As described in [Section 9.3](#), an endpoint is required to validate a peer's new address to confirm the peer's possession of the new address. Until a peer's address is deemed valid, an endpoint MUST limit the rate at which it sends data to this address. The endpoint MUST NOT send more than a minimum congestion window's worth of data per estimated round-trip time (`kMinimumWindow`, as defined in [\[QUIC-RECOVERY\]](#)). In the absence of this limit, an endpoint risks being used for a denial of service attack against an unsuspecting victim. Note that since the endpoint will not have any round-trip time measurements to this address, the estimate SHOULD be the default initial value; see [\[QUIC-RECOVERY\]](#).

If an endpoint skips validation of a peer address as described in [Section 9.3](#), it does not need to limit its sending rate.

9.3.2. On-Path Address Spoofing

An on-path attacker could cause a spurious connection migration by copying and forwarding a packet with a spoofed address such that it arrives before the original packet. The packet with the spoofed address will be seen to come from a migrating connection, and the original packet will be seen as a duplicate and dropped. After a spurious migration, validation of the source address will fail because the entity at the source address does not have the necessary cryptographic keys to read or respond to the `PATH_CHALLENGE` frame that is sent to it even if it wanted to.

To protect the connection from failing due to such a spurious migration, an endpoint **MUST** revert to using the last validated peer address when validation of a new peer address fails.

If an endpoint has no state about the last validated peer address, it **MUST** close the connection silently by discarding all connection state. This results in new packets on the connection being handled generically. For instance, an endpoint **MAY** send a stateless reset in response to any further incoming packets.

Note that receipt of packets with higher packet numbers from the legitimate peer address will trigger another connection migration. This will cause the validation of the address of the spurious migration to be abandoned.

9.3.3. Off-Path Packet Forwarding

An off-path attacker that can observe packets might forward copies of genuine packets to endpoints. If the copied packet arrives before the genuine packet, this will appear as a NAT rebinding. Any genuine packet will be discarded as a duplicate. If the attacker is able to continue forwarding packets, it might be able to cause migration to a path via the attacker. This places the attacker on path, giving it the ability to observe or drop all subsequent packets.

Unlike the attack described in [Section 9.3.2](#), the attacker can ensure that the new path is successfully validated.

This style of attack relies on the attacker using a path that is approximately as fast as the direct path between endpoints. The attack is more reliable if relatively few packets are sent or if packet loss coincides with the attempted attack.

A non-probing packet received on the original path that increases the maximum received packet number will cause the endpoint to move back to that path. Eliciting packets on this path increases the likelihood that the attack is unsuccessful. Therefore, mitigation of this attack relies on triggering the exchange of packets.

In response to an apparent migration, endpoints MUST validate the previously active path using a PATH_CHALLENGE frame. This induces the sending of new packets on that path. If the path is no longer viable, the validation attempt will time out and fail; if the path is viable, but no longer desired, the validation will succeed, but only results in probing packets being sent on the path.

An endpoint that receives a PATH_CHALLENGE on an active path SHOULD send a non-probing packet in response. If the non-probing packet arrives before any copy made by an attacker, this results in the connection being migrated back to the original path. Any subsequent migration to another path restarts this entire process.

This defense is imperfect, but this is not considered a serious problem. If the path via the attack is reliably faster than the original path despite multiple attempts to use that original path, it is not possible to distinguish between attack and an improvement in routing.

An endpoint could also use heuristics to improve detection of this style of attack. For instance, NAT rebinding is improbable if packets were recently received on the old path, similarly rebinding is rare on IPv6 paths. Endpoints can also look for duplicated packets. Conversely, a change in connection ID is more likely to indicate an intentional migration rather than an attack.

9.4. Loss Detection and Congestion Control

The capacity available on the new path might not be the same as the old path. Packets sent on the old path MUST NOT contribute to congestion control or RTT estimation for the new path.

On confirming a peer's ownership of its new address, an endpoint **MUST** immediately reset the congestion controller and round-trip time estimator for the new path to initial values (see Sections A.3 and B.3 in [QUIC-RECOVERY]) unless it has knowledge that a previous send rate or round-trip time estimate is valid for the new path. For instance, an endpoint might infer that a change in only the client's port number is indicative of a NAT rebinding, meaning that the new path is likely to have similar bandwidth and round-trip time. However, this determination will be imperfect. If the determination is incorrect, the congestion controller and the RTT estimator are expected to adapt to the new path. Generally, implementations are advised to be cautious when using previous values on a new path.

There may be apparent reordering at the receiver when an endpoint sends data and probes from/to multiple addresses during the migration period, since the two resulting paths may have different round-trip times. A receiver of packets on multiple paths will still send ACK frames covering all received packets.

While multiple paths might be used during connection migration, a single congestion control context and a single loss recovery context (as described in [QUIC-RECOVERY]) may be adequate. For instance, an endpoint might delay switching to a new congestion control context until it is confirmed that an old path is no longer needed (such as the case in Section 9.3.3).

A sender can make exceptions for probe packets so that their loss detection is independent and does not unduly cause the congestion controller to reduce its sending rate. An endpoint might set a separate timer when a `PATH_CHALLENGE` is sent, which is cancelled if the corresponding `PATH_RESPONSE` is received. If the timer fires before the `PATH_RESPONSE` is received, the endpoint might send a new `PATH_CHALLENGE`, and restart the timer for a longer period of time. This timer **SHOULD** be set as described in Section 6.2.1 of [QUIC-RECOVERY] and **MUST NOT** be more aggressive.

9.5. Privacy Implications of Connection Migration

Using a stable connection ID on multiple network paths allows a passive observer to correlate activity between those paths. An endpoint that moves between networks might not wish to have their activity correlated by any entity other than their peer, so different connection IDs are used when sending from different local addresses, as discussed in Section 5.1. For this to be effective endpoints need to ensure that connection IDs they provide cannot be linked by any other entity.

At any time, endpoints MAY change the Destination Connection ID they send to a value that has not been used on another path.

An endpoint MUST NOT reuse a connection ID when sending from more than one local address, for example when initiating connection migration as described in [Section 9.2](#) or when probing a new network path as described in [Section 9.1](#).

Similarly, an endpoint MUST NOT reuse a connection ID when sending to more than one destination address. Due to network changes outside the control of its peer, an endpoint might receive packets from a new source address with the same destination connection ID, in which case it MAY continue to use the current connection ID with the new remote address while still sending from the same local address.

These requirements regarding connection ID reuse apply only to the sending of packets, as unintentional changes in path without a change in connection ID are possible. For example, after a period of network inactivity, NAT rebinding might cause packets to be sent on a new path when the client resumes sending. An endpoint responds to such an event as described in [Section 9.3](#).

Using different connection IDs for packets sent in both directions on each new network path eliminates the use of the connection ID for linking packets from the same connection across different network paths. Header protection ensures that packet numbers cannot be used to correlate activity. This does not prevent other properties of packets, such as timing and size, from being used to correlate activity.

An endpoint SHOULD NOT initiate migration with a peer that has requested a zero-length connection ID, because traffic over the new path might be trivially linkable to traffic over the old one. If the server is able to route packets with a zero-length connection ID to the right connection, it means that the server is using other information to demultiplex packets. For example, a server might provide a unique address to every client, for instance using HTTP alternative services [[ALTSVC](#)]. Information that might allow correct routing of packets across multiple network paths will also allow activity on those paths to be linked by entities other than the peer.

A client might wish to reduce linkability by employing a new connection ID and source UDP port when sending traffic after a period of inactivity. Changing the UDP port from which it sends packets at the same time might cause the packet to appear as a connection migration. This ensures that the mechanisms that support migration are exercised even for clients that don't experience NAT rebindings or genuine migrations. Changing port number can cause a peer to reset its congestion state (see [Section 9.4](#)), so the port SHOULD only be changed infrequently.

An endpoint that exhausts available connection IDs cannot probe new paths or initiate migration, nor can it respond to probes or attempts by its peer to migrate. To ensure that migration is possible and packets sent on different paths cannot be correlated, endpoints SHOULD provide new connection IDs before peers migrate; see [Section 5.1.1](#). If a peer might have exhausted available connection IDs, a migrating endpoint could include a NEW_CONNECTION_ID frame in all packets sent on a new network path.

9.6. Server's Preferred Address

QUIC allows servers to accept connections on one IP address and attempt to transfer these connections to a more preferred address shortly after the handshake. This is particularly useful when clients initially connect to an address shared by multiple servers but would prefer to use a unicast address to ensure connection stability. This section describes the protocol for migrating a connection to a preferred server address.

Migrating a connection to a new server address mid-connection is left for future work. If a client receives packets from a new server address not indicated by the preferred_address transport parameter, the client SHOULD discard these packets.

9.6.1. Communicating a Preferred Address

A server conveys a preferred address by including the preferred_address transport parameter in the TLS handshake.

Servers MAY communicate a preferred address of each address family (IPv4 and IPv6) to allow clients to pick the one most suited to their network attachment.

Once the handshake is confirmed, the client SHOULD select one of the two server's preferred addresses and initiate path validation (see [Section 8.2](#)) of that address using any previously unused active connection ID, taken from either the preferred_address transport parameter or a NEW_CONNECTION_ID frame.

If path validation succeeds, the client SHOULD immediately begin sending all future packets to the new server address using the new connection ID and discontinue use of the old server address. If path validation fails, the client MUST continue sending all future packets to the server's original IP address.

9.6.2. Responding to Connection Migration

A server might receive a packet addressed to its preferred IP address at any time after it accepts a connection. If this packet contains a PATH_CHALLENGE frame, the server sends a PATH_RESPONSE frame as per [Section 8.2](#). The server MUST send other non-probing frames from its original address until it receives a non-probing packet from the client at its preferred address and until the server has validated the new path.

The server MUST probe on the path toward the client from its preferred address. This helps to guard against spurious migration initiated by an attacker.

Once the server has completed its path validation and has received a non-probing packet with a new largest packet number on its preferred address, the server begins sending non-probing packets to the client exclusively from its preferred IP address. It SHOULD drop packets for this connection received on the old IP address, but MAY continue to process delayed packets.

The addresses that a server provides in the preferred_address transport parameter are only valid for the connection in which they are provided. A client MUST NOT use these for other connections, including connections that are resumed from the current connection.

9.6.3. Interaction of Client Migration and Preferred Address

A client might need to perform a connection migration before it has migrated to the server's preferred address. In this case, the client SHOULD perform path validation to both the original and preferred server address from the client's new address concurrently.

If path validation of the server's preferred address succeeds, the client MUST abandon validation of the original address and migrate to using the server's preferred address. If path validation of the server's preferred address fails but validation of the server's original address succeeds, the client MAY migrate to its new address and continue sending to the server's original address.

If the connection to the server's preferred address is not from the same client address, the server MUST protect against potential attacks as described in [Section 9.3.1](#) and [Section 9.3.2](#). In addition to intentional simultaneous migration, this might also occur because the client's access network used a different NAT binding for the server's preferred address.

Servers SHOULD initiate path validation to the client's new address upon receiving a probe packet from a different address. Servers MUST NOT send more than a minimum congestion window's worth of non-probing packets to the new address before path validation is complete.

A client that migrates to a new address SHOULD use a preferred address from the same address family for the server.

The connection ID provided in the `preferred_address` transport parameter is not specific to the addresses that are provided. This connection ID is provided to ensure that the client has a connection ID available for migration, but the client MAY use this connection ID on any path.

9.7. Use of IPv6 Flow-Label and Migration

Endpoints that send data using IPv6 SHOULD apply an IPv6 flow label in compliance with [\[RFC6437\]](#), unless the local API does not allow setting IPv6 flow labels.

The IPv6 flow label SHOULD be a pseudo-random function of the source and destination addresses, source and destination UDP ports, and the destination CID. The flow label generation MUST be designed to minimize the chances of linkability with a previously used flow label, as this would enable correlating activity on multiple paths; see [Section 9.5](#).

A possible implementation is to compute the flow label as a cryptographic hash function of the source and destination addresses, source and destination UDP ports, destination CID, and a local secret.

10. Connection Termination

An established QUIC connection can be terminated in one of three ways:

- * idle timeout ([Section 10.2](#))
- * immediate close ([Section 10.3](#))

- * stateless reset ([Section 10.4](#))

An endpoint MAY discard connection state if it does not have a validated path on which it can send packets; see [Section 8.2](#).

10.1. Closing and Draining Connection States

The closing and draining connection states exist to ensure that connections close cleanly and that delayed or reordered packets are properly discarded. These states SHOULD persist for at least three times the current Probe Timeout (PTO) interval as defined in [\[QUIC-RECOVERY\]](#).

An endpoint enters a closing period after initiating an immediate close; [Section 10.3](#). While closing, an endpoint MUST NOT send packets unless they contain a CONNECTION_CLOSE frame; see [Section 10.3](#) for details. An endpoint retains only enough information to generate a packet containing a CONNECTION_CLOSE frame and to identify packets as belonging to the connection. The endpoint's selected connection ID and the QUIC version are sufficient information to identify packets for a closing connection; an endpoint can discard all other connection state. An endpoint MAY retain packet protection keys for incoming packets to allow it to read and process a CONNECTION_CLOSE frame.

The draining state is entered once an endpoint receives a signal that its peer is closing or draining. While otherwise identical to the closing state, an endpoint in the draining state MUST NOT send any packets. Retaining packet protection keys is unnecessary once a connection is in the draining state.

An endpoint MAY transition from the closing period to the draining period if it receives a CONNECTION_CLOSE frame or stateless reset, both of which indicate that the peer is also closing or draining. The draining period SHOULD end when the closing period would have ended. In other words, the endpoint can use the same end time, but cease retransmission of the closing packet.

Disposing of connection state prior to the end of the closing or draining period could cause delayed or reordered packets to generate an unnecessary stateless reset. Endpoints that have some alternative means to ensure that late-arriving packets on the connection do not induce a response, such as those that are able to close the UDP socket, MAY use an abbreviated draining period which can allow for faster resource recovery. Servers that retain an open socket for accepting new connections SHOULD NOT exit the closing or draining period early.

Once the closing or draining period has ended, an endpoint SHOULD discard all connection state. This results in new packets on the connection being handled generically. For instance, an endpoint MAY send a stateless reset in response to any further incoming packets.

The draining and closing periods do not apply when a stateless reset ([Section 10.4](#)) is sent.

An endpoint is not expected to handle key updates when it is closing or draining. A key update might prevent the endpoint from moving from the closing state to draining, but it otherwise has no impact.

While in the closing period, an endpoint could receive packets from a new source address, indicating a connection migration; [Section 9](#). An endpoint in the closing state MUST strictly limit the number of packets it sends to this new address until the address is validated; see [Section 8.2](#). A server in the closing state MAY instead choose to discard packets received from a new source address.

10.2. Idle Timeout

If a `max_idle_timeout` is specified by either peer in its transport parameters ([Section 18.2](#)), the connection is silently closed and its state is discarded when it remains idle for longer than the minimum of both peers `max_idle_timeout` values and three times the current Probe Timeout (PTO).

Each endpoint advertises a `max_idle_timeout`, but the effective value at an endpoint is computed as the minimum of the two advertised values. By announcing a `max_idle_timeout`, an endpoint commits to initiating an immediate close ([Section 10.3](#)) if it abandons the connection prior to the effective value.

An endpoint restarts its idle timer when a packet from its peer is received and processed successfully. An endpoint also restarts its idle timer when sending an ack-eliciting packet if no other ack-eliciting packets have been sent since last receiving and processing a packet. Restarting this timer when sending a packet ensures that connections are not closed after new activity is initiated.

10.2.1. Liveness Testing

An endpoint that sends packets close to the effective timeout risks having them be discarded at the peer, since the peer might enter its draining state before these packets arrive.

An endpoint can send a PING or another ack-eliciting frame to test the connection for liveness if the peer could time out soon, such as within a PTO; see Section 6.2 of [QUIC-RECOVERY]. This is especially useful if any available application data cannot be safely retried. Note that the application determines what data is safe to retry.

10.2.2. Deferring Idle Timeout

An endpoint might need to send ack-eliciting packets to avoid an idle timeout if it is expecting response data, but does not have or is unable to send application data.

An implementation of QUIC might provide applications with an option to defer an idle timeout. This facility could be used when the application wishes to avoid losing state that has been associated with an open connection, but does not expect to exchange application data for some time. With this option, an endpoint could send a PING frame periodically to defer an idle timeout; see Section 19.2.

Application protocols that use QUIC SHOULD provide guidance on when deferring an idle timeout is appropriate. Unnecessary sending of PING frames could have a detrimental effect on performance.

A connection will time out if no packets are sent or received for a period longer than the time negotiated using the `max_idle_timeout` transport parameter; see Section 10. However, state in middleboxes might time out earlier than that. Though REQ-5 in [RFC4787] recommends a 2 minute timeout interval, experience shows that sending packets every 15 to 30 seconds is necessary to prevent the majority of middleboxes from losing state for UDP flows.

10.3. Immediate Close

An endpoint sends a CONNECTION_CLOSE frame (Section 19.19) to terminate the connection immediately. A CONNECTION_CLOSE frame causes all streams to immediately become closed; open streams can be assumed to be implicitly reset.

After sending a CONNECTION_CLOSE frame, an endpoint immediately enters the closing state.

During the closing period, an endpoint that sends a CONNECTION_CLOSE frame SHOULD respond to any incoming packet that can be decrypted with another packet containing a CONNECTION_CLOSE frame. Such an endpoint SHOULD limit the number of packets it generates containing a CONNECTION_CLOSE frame. For instance, an endpoint could wait for a progressively increasing number of received packets or amount of time before responding to a received packet.

An endpoint is allowed to drop the packet protection keys when entering the closing period ([Section 10.1](#)) and send a packet containing a CONNECTION_CLOSE in response to any UDP datagram that is received. However, an endpoint without the packet protection keys cannot identify and discard invalid packets. To avoid creating an unwitting amplification attack, such endpoints MUST reduce the frequency with which it sends packets containing a CONNECTION_CLOSE frame. To minimize the state that an endpoint maintains for a closing connection, endpoints MAY send the exact same packet.

Note: Allowing retransmission of a closing packet contradicts other advice in this document that recommends the creation of new packet numbers for every packet. Sending new packet numbers is primarily of advantage to loss recovery and congestion control, which are not expected to be relevant for a closed connection. Retransmitting the final packet requires less state.

New packets from unverified addresses could be used to create an amplification attack; see [Section 8](#). To avoid this, endpoints MUST either limit transmission of CONNECTION_CLOSE frames to validated addresses or drop packets without response if the response would be more than three times larger than the received packet.

After receiving a CONNECTION_CLOSE frame, endpoints enter the draining state. An endpoint that receives a CONNECTION_CLOSE frame MAY send a single packet containing a CONNECTION_CLOSE frame before entering the draining state, using a CONNECTION_CLOSE frame and a NO_ERROR code if appropriate. An endpoint MUST NOT send further packets, which could result in a constant exchange of CONNECTION_CLOSE frames until the closing period on either peer ended.

An immediate close can be used after an application protocol has arranged to close a connection. This might be after the application protocols negotiates a graceful shutdown. The application protocol exchanges whatever messages that are needed to cause both endpoints to agree to close the connection, after which the application requests that the connection be closed. The application protocol can use a CONNECTION_CLOSE frame with an appropriate error code to signal closure.

10.3.1. Immediate Close During the Handshake

When sending `CONNECTION_CLOSE`, the goal is to ensure that the peer will process the frame. Generally, this means sending the frame in a packet with the highest level of packet protection to avoid the packet being discarded. After the handshake is confirmed (see Section 4.1.2 of [QUIC-TLS]), an endpoint **MUST** send any `CONNECTION_CLOSE` frames in a 1-RTT packet. However, prior to confirming the handshake, it is possible that more advanced packet protection keys are not available to the peer, so another `CONNECTION_CLOSE` frame **MAY** be sent in a packet that uses a lower packet protection level. More specifically:

- * A client will always know whether the server has Handshake keys (see Section 17.2.2.1), but it is possible that a server does not know whether the client has Handshake keys. Under these circumstances, a server **SHOULD** send a `CONNECTION_CLOSE` frame in both Handshake and Initial packets to ensure that at least one of them is processable by the client.
- * A client that sends `CONNECTION_CLOSE` in a 0-RTT packet cannot be assured of the server has accepted 0-RTT and so sending a `CONNECTION_CLOSE` frame in an Initial packet makes it more likely that the server can receive the close signal, even if the application error code might not be received.
- * Prior to confirming the handshake, a peer might be unable to process 1-RTT packets, so an endpoint **SHOULD** send `CONNECTION_CLOSE` in both Handshake and 1-RTT packets. A server **SHOULD** also send `CONNECTION_CLOSE` in an Initial packet.

Sending a `CONNECTION_CLOSE` of type `0x1d` in an Initial or Handshake packet could expose application state or be used to alter application state. A `CONNECTION_CLOSE` of type `0x1d` **MUST** be replaced by a `CONNECTION_CLOSE` of type `0x1c` when sending the frame in Initial or Handshake packets. Otherwise, information about the application state might be revealed. Endpoints **MUST** clear the value of the Reason Phrase field and **SHOULD** use the `APPLICATION_ERROR` code when converting to a `CONNECTION_CLOSE` of type `0x1c`.

`CONNECTION_CLOSE` frames sent in multiple packet types can be coalesced into a single UDP datagram; see Section 12.2.

An endpoint might send a `CONNECTION_CLOSE` frame in an Initial packet or in response to unauthenticated information received in Initial or Handshake packets. Such an immediate close might expose legitimate connections to a denial of service. QUIC does not include defensive measures for on-path attacks during the handshake; see Section 21.1.

However, at the cost of reducing feedback about errors for legitimate peers, some forms of denial of service can be made more difficult for an attacker if endpoints discard illegal packets rather than terminating a connection with `CONNECTION_CLOSE`. For this reason, endpoints *MAY* discard packets rather than immediately close if errors are detected in packets that lack authentication.

An endpoint that has not established state, such as a server that detects an error in an Initial packet, does not enter the closing state. An endpoint that has no state for the connection does not enter a closing or draining period on sending a `CONNECTION_CLOSE` frame.

10.4. Stateless Reset

A stateless reset is provided as an option of last resort for an endpoint that does not have access to the state of a connection. A crash or outage might result in peers continuing to send data to an endpoint that is unable to properly continue the connection. An endpoint *MAY* send a stateless reset in response to receiving a packet that it cannot associate with an active connection.

A stateless reset is not appropriate for signaling error conditions. An endpoint that wishes to communicate a fatal connection error *MUST* use a `CONNECTION_CLOSE` frame if it has sufficient state to do so.

To support this process, a token is sent by endpoints. The token is carried in the Stateless Reset Token field of a `NEW_CONNECTION_ID` frame. Servers can also specify a `stateless_reset_token` transport parameter during the handshake that applies to the connection ID that it selected during the handshake; clients cannot use this transport parameter because their transport parameters don't have confidentiality protection. These tokens are protected by encryption, so only client and server know their value. Tokens are invalidated when their associated connection ID is retired via a `RETIRE_CONNECTION_ID` frame ([Section 19.16](#)).

An endpoint that receives packets that it cannot process sends a packet in the following layout:

```
Stateless Reset {  
  Fixed Bits (2) = 1,  
  Unpredictable Bits (38..),  
  Stateless Reset Token (128),  
}
```

Figure 9: Stateless Reset Packet

This design ensures that a stateless reset packet is - to the extent possible - indistinguishable from a regular packet with a short header.

A stateless reset uses an entire UDP datagram, starting with the first two bits of the packet header. The remainder of the first byte and an arbitrary number of bytes following it that are set to unpredictable values. The last 16 bytes of the datagram contain a Stateless Reset Token.

To entities other than its intended recipient, a stateless reset will appear to be a packet with a short header. For the stateless reset to appear as a valid QUIC packet, the Unpredictable Bits field needs to include at least 38 bits of data (or 5 bytes, less the two fixed bits).

A minimum size of 21 bytes does not guarantee that a stateless reset is difficult to distinguish from other packets if the recipient requires the use of a connection ID. To prevent a resulting stateless reset from being trivially distinguishable from a valid packet, all packets sent by an endpoint SHOULD be padded to at least 22 bytes longer than the minimum connection ID that the endpoint might use. An endpoint that sends a stateless reset in response to a packet that is 43 bytes or less in length SHOULD send a stateless reset that is one byte shorter than the packet it responds to.

These values assume that the Stateless Reset Token is the same length as the minimum expansion of the packet protection AEAD. Additional unpredictable bytes are necessary if the endpoint could have negotiated a packet protection scheme with a larger minimum expansion.

An endpoint MUST NOT send a stateless reset that is three times or more larger than the packet it receives to avoid being used for amplification. [Section 10.4.3](#) describes additional limits on stateless reset size.

Endpoints MUST discard packets that are too small to be valid QUIC packets. With the set of AEAD functions defined in [\[QUIC-TLS\]](#), packets that are smaller than 21 bytes are never valid.

Endpoints MUST send stateless reset packets formatted as a packet with a short header. However, endpoints MUST treat any packet ending in a valid stateless reset token as a stateless reset, as other QUIC versions might allow the use of a long header.

An endpoint MAY send a stateless reset in response to a packet with a long header. Sending a stateless reset is not effective prior to the stateless reset token being available to a peer. In this QUIC version, packets with a long header are only used during connection establishment. Because the stateless reset token is not available until connection establishment is complete or near completion, ignoring an unknown packet with a long header might be as effective as sending a stateless reset.

An endpoint cannot determine the Source Connection ID from a packet with a short header, therefore it cannot set the Destination Connection ID in the stateless reset packet. The Destination Connection ID will therefore differ from the value used in previous packets. A random Destination Connection ID makes the connection ID appear to be the result of moving to a new connection ID that was provided using a NEW_CONNECTION_ID frame ([Section 19.15](#)).

Using a randomized connection ID results in two problems:

- * The packet might not reach the peer. If the Destination Connection ID is critical for routing toward the peer, then this packet could be incorrectly routed. This might also trigger another Stateless Reset in response; see [Section 10.4.3](#). A Stateless Reset that is not correctly routed is an ineffective error detection and recovery mechanism. In this case, endpoints will need to rely on other methods - such as timers - to detect that the connection has failed.
- * The randomly generated connection ID can be used by entities other than the peer to identify this as a potential stateless reset. An endpoint that occasionally uses different connection IDs might introduce some uncertainty about this.

This stateless reset design is specific to QUIC version 1. An endpoint that supports multiple versions of QUIC needs to generate a stateless reset that will be accepted by peers that support any version that the endpoint might support (or might have supported prior to losing state). Designers of new versions of QUIC need to be aware of this and either reuse this design, or use a portion of the packet other than the last 16 bytes for carrying data.

10.4.1. Detecting a Stateless Reset

An endpoint detects a potential stateless reset using the trailing 16 bytes of the UDP datagram. An endpoint remembers all Stateless Reset Tokens associated with the connection IDs and remote addresses for datagrams it has recently sent. This includes Stateless Reset Tokens from NEW_CONNECTION_ID frames and the server's transport parameters but excludes Stateless Reset Tokens associated with connection IDs that are either unused or retired. The endpoint identifies a received datagram as a stateless reset by comparing the last 16 bytes of the datagram with all Stateless Reset Tokens associated with the remote address on which the datagram was received.

This comparison can be performed for every inbound datagram. Endpoints MAY skip this check if any packet from a datagram is successfully processed. However, the comparison MUST be performed when the first packet in an incoming datagram either cannot be associated with a connection, or cannot be decrypted.

An endpoint MUST NOT check for any Stateless Reset Tokens associated with connection IDs it has not used or for connection IDs that have been retired.

When comparing a datagram to Stateless Reset Token values, endpoints MUST perform the comparison without leaking information about the value of the token. For example, performing this comparison in constant time protects the value of individual Stateless Reset Tokens from information leakage through timing side channels. Another approach would be to store and compare the transformed values of Stateless Reset Tokens instead of the raw token values, where the transformation is defined as a cryptographically-secure pseudo-random function using a secret key (e.g., block cipher, HMAC [RFC2104]). An endpoint is not expected to protect information about whether a packet was successfully decrypted, or the number of valid Stateless Reset Tokens.

If the last 16 bytes of the datagram are identical in value to a Stateless Reset Token, the endpoint MUST enter the draining period and not send any further packets on this connection.

10.4.2. Calculating a Stateless Reset Token

The stateless reset token MUST be difficult to guess. In order to create a Stateless Reset Token, an endpoint could randomly generate [RFC4086] a secret for every connection that it creates. However, this presents a coordination problem when there are multiple instances in a cluster or a storage problem for an endpoint that might lose state. Stateless reset specifically exists to handle the case where state is lost, so this approach is suboptimal.

A single static key can be used across all connections to the same endpoint by generating the proof using a second iteration of a preimage-resistant function that takes a static key and the connection ID chosen by the endpoint (see Section 5.1) as input. An endpoint could use HMAC [RFC2104] (for example, `HMAC(static_key, connection_id)`) or HKDF [RFC5869] (for example, using the static key as input keying material, with the connection ID as salt). The output of this function is truncated to 16 bytes to produce the Stateless Reset Token for that connection.

An endpoint that loses state can use the same method to generate a valid Stateless Reset Token. The connection ID comes from the packet that the endpoint receives.

This design relies on the peer always sending a connection ID in its packets so that the endpoint can use the connection ID from a packet to reset the connection. An endpoint that uses this design MUST either use the same connection ID length for all connections or encode the length of the connection ID such that it can be recovered without state. In addition, it cannot provide a zero-length connection ID.

Revealing the Stateless Reset Token allows any entity to terminate the connection, so a value can only be used once. This method for choosing the Stateless Reset Token means that the combination of connection ID and static key MUST NOT be used for another connection. A denial of service attack is possible if the same connection ID is used by instances that share a static key, or if an attacker can cause a packet to be routed to an instance that has no state but the same static key; see Section 21.9. A connection ID from a connection that is reset by revealing the Stateless Reset Token MUST NOT be reused for new connections at nodes that share a static key.

The same Stateless Reset Token MUST NOT be used for multiple connection IDs. Endpoints are not required to compare new values against all previous values, but a duplicate value MAY be treated as a connection error of type `PROTOCOL_VIOLATION`.

Note that Stateless Reset packets do not have any cryptographic protection.

10.4.3. Looping

The design of a Stateless Reset is such that without knowing the stateless reset token it is indistinguishable from a valid packet. For instance, if a server sends a Stateless Reset to another server it might receive another Stateless Reset in response, which could lead to an infinite exchange.

An endpoint **MUST** ensure that every Stateless Reset that it sends is smaller than the packet which triggered it, unless it maintains state sufficient to prevent looping. In the event of a loop, this results in packets eventually being too small to trigger a response.

An endpoint can remember the number of Stateless Reset packets that it has sent and stop generating new Stateless Reset packets once a limit is reached. Using separate limits for different remote addresses will ensure that Stateless Reset packets can be used to close connections when other peers or connections have exhausted limits.

Reducing the size of a Stateless Reset below 41 bytes means that the packet could reveal to an observer that it is a Stateless Reset, depending upon the length of the peer's connection IDs. Conversely, refusing to send a Stateless Reset in response to a small packet might result in Stateless Reset not being useful in detecting cases of broken connections where only very small packets are sent; such failures might only be detected by other means, such as timers.

11. Error Handling

An endpoint that detects an error **SHOULD** signal the existence of that error to its peer. Both transport-level and application-level errors can affect an entire connection; see [Section 11.1](#). Only application-level errors can be isolated to a single stream; see [Section 11.2](#).

The most appropriate error code ([Section 20](#)) **SHOULD** be included in the frame that signals the error. Where this specification identifies error conditions, it also identifies the error code that is used; though these are worded as requirements, different implementation strategies might lead to different errors being reported. In particular, an endpoint **MAY** use any applicable error code when it detects an error condition; a generic error code (such as `PROTOCOL_VIOLATION` or `INTERNAL_ERROR`) can always be used in place of specific error codes.

A stateless reset ([Section 10.4](#)) is not suitable for any error that can be signaled with a CONNECTION_CLOSE or RESET_STREAM frame. A stateless reset MUST NOT be used by an endpoint that has the state necessary to send a frame on the connection.

11.1. Connection Errors

Errors that result in the connection being unusable, such as an obvious violation of protocol semantics or corruption of state that affects an entire connection, MUST be signaled using a CONNECTION_CLOSE frame ([Section 19.19](#)). An endpoint MAY close the connection in this manner even if the error only affects a single stream.

Application protocols can signal application-specific protocol errors using the application-specific variant of the CONNECTION_CLOSE frame. Errors that are specific to the transport, including all those described in this document, are carried in the QUIC-specific variant of the CONNECTION_CLOSE frame.

A CONNECTION_CLOSE frame could be sent in a packet that is lost. An endpoint SHOULD be prepared to retransmit a packet containing a CONNECTION_CLOSE frame if it receives more packets on a terminated connection. Limiting the number of retransmissions and the time over which this final packet is sent limits the effort expended on terminated connections.

An endpoint that chooses not to retransmit packets containing a CONNECTION_CLOSE frame risks a peer missing the first such packet. The only mechanism available to an endpoint that continues to receive data for a terminated connection is to use the stateless reset process ([Section 10.4](#)).

11.2. Stream Errors

If an application-level error affects a single stream, but otherwise leaves the connection in a recoverable state, the endpoint can send a RESET_STREAM frame ([Section 19.4](#)) with an appropriate error code to terminate just the affected stream.

Resetting a stream without the involvement of the application protocol could cause the application protocol to enter an unrecoverable state. RESET_STREAM MUST only be instigated by the application protocol that uses QUIC.

The semantics of the application error code carried in RESET_STREAM are defined by the application protocol. Only the application protocol is able to cause a stream to be terminated. A local

instance of the application protocol uses a direct API call and a remote instance uses the `STOP_SENDING` frame, which triggers an automatic `RESET_STREAM`.

Application protocols SHOULD define rules for handling streams that are prematurely cancelled by either endpoint.

12. Packets and Frames

QUIC endpoints communicate by exchanging packets. Packets have confidentiality and integrity protection; see [Section 12.1](#). Packets are carried in UDP datagrams; see [Section 12.2](#).

This version of QUIC uses the long packet header during connection establishment; see [Section 17.2](#). Packets with the long header are Initial ([Section 17.2.2](#)), 0-RTT ([Section 17.2.3](#)), Handshake ([Section 17.2.4](#)), and Retry ([Section 17.2.5](#)). Version negotiation uses a version-independent packet with a long header; see [Section 17.2.1](#).

Packets with the short header are designed for minimal overhead and are used after a connection is established and 1-RTT keys are available; see [Section 17.3](#).

12.1. Protected Packets

All QUIC packets except Version Negotiation packets use authenticated encryption with additional data (AEAD) [[RFC5116](#)] to provide confidentiality and integrity protection. Retry packets use AEAD to provide integrity protection. Details of packet protection are found in [[QUIC-TLS](#)]; this section includes an overview of the process.

Initial packets are protected using keys that are statically derived. This packet protection is not effective confidentiality protection. Initial protection only exists to ensure that the sender of the packet is on the network path. Any entity that receives the Initial packet from a client can recover the keys necessary to remove packet protection or to generate packets that will be successfully authenticated.

All other packets are protected with keys derived from the cryptographic handshake. The type of the packet from the long header or key phase from the short header are used to identify which encryption keys are used. Packets protected with 0-RTT and 1-RTT keys are expected to have confidentiality and data origin authentication; the cryptographic handshake ensures that only the communicating endpoints receive the corresponding keys.

The packet number field contains a packet number, which has additional confidentiality protection that is applied after packet protection is applied; see [QUIC-TLS] for details. The underlying packet number increases with each packet sent in a given packet number space; see Section 12.3 for details.

12.2. Coalescing Packets

Initial (Section 17.2.2), 0-RTT (Section 17.2.3), and Handshake (Section 17.2.4) packets contain a Length field, which determines the end of the packet. The length includes both the Packet Number and Payload fields, both of which are confidentiality protected and initially of unknown length. The length of the Payload field is learned once header protection is removed.

Using the Length field, a sender can coalesce multiple QUIC packets into one UDP datagram. This can reduce the number of UDP datagrams needed to complete the cryptographic handshake and start sending data. This can also be used to construct PMTU probes; see Section 14.4.1. Receivers MUST be able to process coalesced packets.

Coalescing packets in order of increasing encryption levels (Initial, 0-RTT, Handshake, 1-RTT; see Section 4.1.4 of [QUIC-TLS]) makes it more likely the receiver will be able to process all the packets in a single pass. A packet with a short header does not include a length, so it can only be the last packet included in a UDP datagram. An endpoint SHOULD NOT coalesce multiple packets at the same encryption level.

Senders MUST NOT coalesce QUIC packets for different connections into a single UDP datagram. Receivers SHOULD ignore any subsequent packets with a different Destination Connection ID than the first packet in the datagram.

Every QUIC packet that is coalesced into a single UDP datagram is separate and complete. The receiver of coalesced QUIC packets MUST individually process each QUIC packet and separately acknowledge them, as if they were received as the payload of different UDP datagrams. For example, if decryption fails (because the keys are not available or any other reason), the receiver MAY either discard or buffer the packet for later processing and MUST attempt to process the remaining packets.

Retry packets ([Section 17.2.5](#)), Version Negotiation packets ([Section 17.2.1](#)), and packets with a short header ([Section 17.3](#)) do not contain a Length field and so cannot be followed by other packets in the same UDP datagram. Note also that there is no situation where a Retry or Version Negotiation packet is coalesced with another packet.

12.3. Packet Numbers

The packet number is an integer in the range 0 to $2^{62}-1$. This number is used in determining the cryptographic nonce for packet protection. Each endpoint maintains a separate packet number for sending and receiving.

Packet numbers are limited to this range because they need to be representable in whole in the Largest Acknowledged field of an ACK frame ([Section 19.3](#)). When present in a long or short header however, packet numbers are reduced and encoded in 1 to 4 bytes; see [Section 17.1](#).

Version Negotiation ([Section 17.2.1](#)) and Retry ([Section 17.2.5](#)) packets do not include a packet number.

Packet numbers are divided into 3 spaces in QUIC:

- * Initial space: All Initial packets ([Section 17.2.2](#)) are in this space.
- * Handshake space: All Handshake packets ([Section 17.2.4](#)) are in this space.
- * Application data space: All 0-RTT and 1-RTT encrypted packets ([Section 12.1](#)) are in this space.

As described in [[QUIC-TLS](#)], each packet type uses different protection keys.

Conceptually, a packet number space is the context in which a packet can be processed and acknowledged. Initial packets can only be sent with Initial packet protection keys and acknowledged in packets which are also Initial packets. Similarly, Handshake packets are sent at the Handshake encryption level and can only be acknowledged in Handshake packets.

This enforces cryptographic separation between the data sent in the different packet number spaces. Packet numbers in each space start at packet number 0. Subsequent packets sent in the same packet number space MUST increase the packet number by at least one.

0-RTT and 1-RTT data exist in the same packet number space to make loss recovery algorithms easier to implement between the two packet types.

A QUIC endpoint **MUST NOT** reuse a packet number within the same packet number space in one connection. If the packet number for sending reaches $2^{62} - 1$, the sender **MUST** close the connection without sending a `CONNECTION_CLOSE` frame or any further packets; an endpoint **MAY** send a Stateless Reset ([Section 10.4](#)) in response to further packets that it receives.

A receiver **MUST** discard a newly unprotected packet unless it is certain that it has not processed another packet with the same packet number from the same packet number space. Duplicate suppression **MUST** happen after removing packet protection for the reasons described in [Section 9.3](#) of [[QUIC-TLS](#)].

Endpoints that track all individual packets for the purposes of detecting duplicates are at risk of accumulating excessive state. The data required for detecting duplicates can be limited by maintaining a minimum packet number below which all packets are immediately dropped. Any minimum needs to account for large variations in round trip time, which includes the possibility that a peer might probe network paths with a much larger round trip times; see [Section 9](#).

Packet number encoding at a sender and decoding at a receiver are described in [Section 17.1](#).

12.4. Frames and Frame Types

The payload of QUIC packets, after removing packet protection, consists of a sequence of complete frames, as shown in [Figure 10](#). Version Negotiation, Stateless Reset, and Retry packets do not contain frames.

```
Packet Payload {  
    Frame (...) ...,  
}
```

Figure 10: QUIC Payload

The payload of a packet that contains frames **MUST** contain at least one frame, and **MAY** contain multiple frames and multiple frame types. Frames always fit within a single QUIC packet and cannot span multiple packets.

Each frame begins with a Frame Type, indicating its type, followed by additional type-dependent fields:

```
Frame {  
  Frame Type (i),  
  Type-Dependent Fields (...),  
}
```

Figure 11: Generic Frame Layout

The frame types defined in this specification are listed in Table 3. The Frame Type in ACK, STREAM, MAX_STREAMS, STREAMS_BLOCKED, and CONNECTION_CLOSE frames is used to carry other frame-specific flags. For all other frames, the Frame Type field simply identifies the frame. These frames are explained in more detail in [Section 19](#).

Type Value	Frame Type Name	Definition	Packets
0x00	PADDING	Section 19.1	IH01
0x01	PING	Section 19.2	IH01
0x02 - 0x03	ACK	Section 19.3	IH_1
0x04	RESET_STREAM	Section 19.4	__01
0x05	STOP_SENDING	Section 19.5	__01
0x06	CRYPTO	Section 19.6	IH_1
0x07	NEW_TOKEN	Section 19.7	___1
0x08 - 0x0f	STREAM	Section 19.8	__01
0x10	MAX_DATA	Section 19.9	__01
0x11	MAX_STREAM_DATA	Section 19.10	__01
0x12 - 0x13	MAX_STREAMS	Section 19.11	__01
0x14	DATA_BLOCKED	Section 19.12	__01
0x15	STREAM_DATA_BLOCKED	Section 19.13	__01
0x16 - 0x17	STREAMS_BLOCKED	Section 19.14	__01
0x18	NEW_CONNECTION_ID	Section 19.15	__01
0x19	RETIRE_CONNECTION_ID	Section 19.16	__01
0x1a	PATH_CHALLENGE	Section 19.17	__01
0x1b	PATH_RESPONSE	Section 19.18	__01
0x1c - 0x1d	CONNECTION_CLOSE	Section 19.19	ih01
0x1e	HANDSHAKE_DONE	Section 19.20	___1

Table 3: Frame Types

The "Packets" column in Table 3 does not form part of the IANA registry; see [Section 22.3](#). This column lists the types of packets that each frame type could appear in, indicated by the following characters:

I: Initial ([Section 17.2.2](#))

H: Handshake ([Section 17.2.4](#))

0: 0-RTT ([Section 17.2.3](#))

1: 1-RTT ([Section 17.3](#))

ih: A CONNECTION_CLOSE frame of type 0x1d cannot appear in Initial or Handshake packets.

Section 4 of [[QUIC-TLS](#)] provides more detail about these restrictions. Note that all frames can appear in 1-RTT packets.

An endpoint MUST treat the receipt of a frame of unknown type as a connection error of type FRAME_ENCODING_ERROR.

All QUIC frames are idempotent in this version of QUIC. That is, a valid frame does not cause undesirable side effects or errors when received more than once.

The Frame Type field uses a variable length integer encoding (see [Section 16](#)) with one exception. To ensure simple and efficient implementations of frame parsing, a frame type MUST use the shortest possible encoding. For frame types defined in this document, this means a single-byte encoding, even though it is possible to encode these values as a two-, four- or eight-byte variable length integer. For instance, though 0x4001 is a legitimate two-byte encoding for a variable-length integer with a value of 1, PING frames are always encoded as a single byte with the value 0x01. This rule applies to all current and future QUIC frame types. An endpoint MAY treat the receipt of a frame type that uses a longer encoding than necessary as a connection error of type PROTOCOL_VIOLATION.

13. Packetization and Reliability

A sender bundles one or more frames in a QUIC packet; see [Section 12.4](#).

A sender can minimize per-packet bandwidth and computational costs by including as many frames as possible in each QUIC packet. A sender MAY wait for a short period of time to collect multiple frames before sending a packet that is not maximally packed, to avoid sending out

large numbers of small packets. An implementation MAY use knowledge about application sending behavior or heuristics to determine whether and for how long to wait. This waiting period is an implementation decision, and an implementation should be careful to delay conservatively, since any delay is likely to increase application-visible latency.

Stream multiplexing is achieved by interleaving STREAM frames from multiple streams into one or more QUIC packets. A single QUIC packet can include multiple STREAM frames from one or more streams.

One of the benefits of QUIC is avoidance of head-of-line blocking across multiple streams. When a packet loss occurs, only streams with data in that packet are blocked waiting for a retransmission to be received, while other streams can continue making progress. Note that when data from multiple streams is included in a single QUIC packet, loss of that packet blocks all those streams from making progress. Implementations are advised to include as few streams as necessary in outgoing packets without losing transmission efficiency to underfilled packets.

13.1. Packet Processing

A packet MUST NOT be acknowledged until packet protection has been successfully removed and all frames contained in the packet have been processed. For STREAM frames, this means the data has been enqueued in preparation to be received by the application protocol, but it does not require that data is delivered and consumed.

Once the packet has been fully processed, a receiver acknowledges receipt by sending one or more ACK frames containing the packet number of the received packet.

An endpoint SHOULD treat receipt of an acknowledgment for a packet it did not send as a connection error of type `PROTOCOL_VIOLATION`, if it is able to detect the condition.

13.2. Generating Acknowledgements

Endpoints acknowledge all packets they receive and process. However, only ack-eliciting packets cause an ACK frame to be sent within the maximum ack delay. Packets that are not ack-eliciting are only acknowledged when an ACK frame is sent for other reasons.

When sending a packet for any reason, an endpoint SHOULD attempt to bundle an ACK frame if one has not been sent recently. Doing so helps with timely loss detection at the peer.

In general, frequent feedback from a receiver improves loss and congestion response, but this has to be balanced against excessive load generated by a receiver that sends an ACK frame in response to every ack-eliciting packet. The guidance offered below seeks to strike this balance.

13.2.1. Sending ACK Frames

Every packet SHOULD be acknowledged at least once, and ack-eliciting packets MUST be acknowledged at least once within the maximum ack delay. An endpoint communicates its maximum delay using the `max_ack_delay` transport parameter; see [Section 18.2](#). `max_ack_delay` declares an explicit contract: an endpoint promises to never intentionally delay acknowledgments of an ack-eliciting packet by more than the indicated value. If it does, any excess accrues to the RTT estimate and could result in spurious or delayed retransmissions from the peer. For Initial and Handshake packets, a `max_ack_delay` of 0 is used. The sender uses the receiver's `max_ack_delay` value in determining timeouts for timer-based retransmission, as detailed in [Section 6.2](#) of [\[QUIC-RECOVERY\]](#).

Since packets containing only ACK frames are not congestion controlled, an endpoint MUST NOT send more than one such packet in response to receiving an ack-eliciting packet.

An endpoint MUST NOT send a non-ack-eliciting packet in response to a non-ack-eliciting packet, even if there are packet gaps which precede the received packet. This avoids an infinite feedback loop of acknowledgments, which could prevent the connection from ever becoming idle. Non-ack-eliciting packets are eventually acknowledged when the endpoint sends an ACK frame in response to other events.

In order to assist loss detection at the sender, an endpoint SHOULD send an ACK frame immediately on receiving an ack-eliciting packet that is out of order. The endpoint SHOULD NOT continue sending ACK frames immediately unless more ack-eliciting packets are received out of order. If every subsequent ack-eliciting packet arrives out of order, then an ACK frame SHOULD be sent immediately for every received ack-eliciting packet.

Similarly, packets marked with the ECN Congestion Experienced (CE) codepoint in the IP header SHOULD be acknowledged immediately, to reduce the peer's response time to congestion events.

The algorithms in [\[QUIC-RECOVERY\]](#) are expected to be resilient to receivers that do not follow guidance offered above. However, an implementer should only deviate from these requirements after careful consideration of the performance implications of doing so.

An endpoint that is only sending ACK frames will not receive acknowledgments from its peer unless those acknowledgements are included in packets with ack-eliciting frames. An endpoint SHOULD bundle ACK frames with other frames when there are new ack-eliciting packets to acknowledge. When only non-ack-eliciting packets need to be acknowledged, an endpoint MAY wait until an ack-eliciting packet has been received to bundle an ACK frame with outgoing frames.

13.2.2. Acknowledgement Frequency

A receiver determines how frequently to send acknowledgements in response to ack-eliciting packets. This determination involves a tradeoff.

Endpoints rely on timely acknowledgment to detect loss; see Section 6 of [QUIC-RECOVERY]. Window-based congestion controllers, such as the one in Section 7 of [QUIC-RECOVERY], rely on acknowledgments to manage their congestion window. In both cases, delaying acknowledgments can adversely affect performance.

On the other hand, reducing the frequency of packets that carry only acknowledgements reduces packet transmission and processing cost at both endpoints. It can also improve connection throughput on severely asymmetric links; see Section 3 of [RFC3449].

A receiver SHOULD send an ACK frame after receiving at least two ack-eliciting packets. This recommendation is general in nature and consistent with recommendations for TCP endpoint behavior [RFC5681]. Knowledge of network conditions, knowledge of the peer's congestion controller, or further research and experimentation might suggest alternative acknowledgment strategies with better performance characteristics.

A receiver MAY process multiple available packets before determining whether to send an ACK frame in response.

13.2.3. Managing ACK Ranges

When an ACK frame is sent, one or more ranges of acknowledged packets are included. Including older packets reduces the chance of spurious retransmits caused by losing previously sent ACK frames, at the cost of larger ACK frames.

ACK frames SHOULD always acknowledge the most recently received packets, and the more out-of-order the packets are, the more important it is to send an updated ACK frame quickly, to prevent the peer from declaring a packet as lost and spuriously retransmitting the frames it contains. An ACK frame is expected to fit within a single QUIC packet. If it does not, then older ranges (those with the smallest packet numbers) are omitted.

[Section 13.2.4](#) and [Section 13.2.5](#) describe an exemplary approach for determining what packets to acknowledge in each ACK frame. Though the goal of these algorithms is to generate an acknowledgment for every packet that is processed, it is still possible for acknowledgments to be lost. A sender cannot expect to receive an acknowledgment for every packet that the receiver processes.

13.2.4. Receiver Tracking of ACK Frames

When a packet containing an ACK frame is sent, the largest acknowledged in that frame may be saved. When a packet containing an ACK frame is acknowledged, the receiver can stop acknowledging packets less than or equal to the largest acknowledged in the sent ACK frame.

In cases without ACK frame loss, this algorithm allows for a minimum of 1 RTT of reordering. In cases with ACK frame loss and reordering, this approach does not guarantee that every acknowledgement is seen by the sender before it is no longer included in the ACK frame. Packets could be received out of order and all subsequent ACK frames containing them could be lost. In this case, the loss recovery algorithm could cause spurious retransmits, but the sender will continue making forward progress.

13.2.5. Limiting ACK Ranges

A receiver limits the number of ACK Ranges ([Section 19.3.1](#)) it remembers and sends in ACK frames, both to limit the size of ACK frames and to avoid resource exhaustion. After receiving acknowledgments for an ACK frame, the receiver SHOULD stop tracking those acknowledged ACK Ranges.

It is possible that retaining many ACK Ranges could cause an ACK frame to become too large. A receiver can discard unacknowledged ACK Ranges to limit ACK frame size, at the cost of increased retransmissions from the sender. This is necessary if an ACK frame would be too large to fit in a packet, however receivers MAY also limit ACK frame size further to preserve space for other frames.

A receiver MUST retain an ACK Range unless it can ensure that it will not subsequently accept packets with numbers in that range. Maintaining a minimum packet number that increases as ranges are discarded is one way to achieve this with minimal state.

Receivers can discard all ACK Ranges, but they MUST retain the largest packet number that has been successfully processed as that is used to recover packet numbers from subsequent packets; see [Section 17.1](#).

A receiver SHOULD include an ACK Range containing the largest received packet number in every ACK frame. The Largest Acknowledged field is used in ECN validation at a sender and including a lower value than what was included in a previous ACK frame could cause ECN to be unnecessarily disabled; see [Section 13.4.2](#).

A receiver that sends only non-ack-eliciting packets, such as ACK frames, might not receive an acknowledgement for a long period of time. This could cause the receiver to maintain state for a large number of ACK frames for a long period of time, and ACK frames it sends could be unnecessarily large. In such a case, a receiver could bundle a PING or other small ack-eliciting frame occasionally, such as once per round trip, to elicit an ACK from the peer.

A receiver MUST NOT bundle an ack-eliciting frame with all packets that would otherwise be non-ack-eliciting, to avoid an infinite feedback loop of acknowledgements.

13.2.6. Measuring and Reporting Host Delay

An endpoint measures the delays intentionally introduced between the time the packet with the largest packet number is received and the time an acknowledgment is sent. The endpoint encodes this delay in the Ack Delay field of an ACK frame; see [Section 19.3](#). This allows the receiver of the ACK to adjust for any intentional delays, which is important for getting a better estimate of the path RTT when acknowledgments are delayed. A packet might be held in the OS kernel or elsewhere on the host before being processed. An endpoint MUST NOT include delays that it does not control when populating the Ack Delay field in an ACK frame.

13.2.7. ACK Frames and Packet Protection

ACK frames MUST only be carried in a packet that has the same packet number space as the packet being ACKed; see [Section 12.1](#). For instance, packets that are protected with 1-RTT keys MUST be acknowledged in packets that are also protected with 1-RTT keys.

Packets that a client sends with 0-RTT packet protection **MUST** be acknowledged by the server in packets protected by 1-RTT keys. This can mean that the client is unable to use these acknowledgments if the server cryptographic handshake messages are delayed or lost. Note that the same limitation applies to other data sent by the server protected by the 1-RTT keys.

13.2.8. PADDING Frames Consume Congestion Window

Packets containing PADDING frames are considered to be in flight for congestion control purposes [[QUIC-RECOVERY](#)]. Packets containing only PADDING frames therefore consume congestion window but do not generate acknowledgments that will open the congestion window. To avoid a deadlock, a sender **SHOULD** ensure that other frames are sent periodically in addition to PADDING frames to elicit acknowledgments from the receiver.

13.3. Retransmission of Information

QUIC packets that are determined to be lost are not retransmitted whole. The same applies to the frames that are contained within lost packets. Instead, the information that might be carried in frames is sent again in new frames as needed.

New frames and packets are used to carry information that is determined to have been lost. In general, information is sent again when a packet containing that information is determined to be lost and sending ceases when a packet containing that information is acknowledged.

- * Data sent in CRYPTO frames is retransmitted according to the rules in [[QUIC-RECOVERY](#)], until all data has been acknowledged. Data in CRYPTO frames for Initial and Handshake packets is discarded when keys for the corresponding packet number space are discarded.
- * Application data sent in STREAM frames is retransmitted in new STREAM frames unless the endpoint has sent a RESET_STREAM for that stream. Once an endpoint sends a RESET_STREAM frame, no further STREAM frames are needed.
- * ACK frames carry the most recent set of acknowledgements and the Ack Delay from the largest acknowledged packet, as described in [Section 13.2.1](#). Delaying the transmission of packets containing ACK frames or sending old ACK frames can cause the peer to generate an inflated RTT sample or unnecessarily disable ECN.

- * Cancellation of stream transmission, as carried in a RESET_STREAM frame, is sent until acknowledged or until all stream data is acknowledged by the peer (that is, either the "Reset Recvd" or "Data Recvd" state is reached on the sending part of the stream). The content of a RESET_STREAM frame MUST NOT change when it is sent again.
- * Similarly, a request to cancel stream transmission, as encoded in a STOP_SENDING frame, is sent until the receiving part of the stream enters either a "Data Recvd" or "Reset Recvd" state; see [Section 3.5](#).
- * Connection close signals, including packets that contain CONNECTION_CLOSE frames, are not sent again when packet loss is detected, but as described in [Section 10](#).
- * The current connection maximum data is sent in MAX_DATA frames. An updated value is sent in a MAX_DATA frame if the packet containing the most recently sent MAX_DATA frame is declared lost, or when the endpoint decides to update the limit. Care is necessary to avoid sending this frame too often as the limit can increase frequently and cause an unnecessarily large number of MAX_DATA frames to be sent.
- * The current maximum stream data offset is sent in MAX_STREAM_DATA frames. Like MAX_DATA, an updated value is sent when the packet containing the most recent MAX_STREAM_DATA frame for a stream is lost or when the limit is updated, with care taken to prevent the frame from being sent too often. An endpoint SHOULD stop sending MAX_STREAM_DATA frames when the receiving part of the stream enters a "Size Known" state.
- * The limit on streams of a given type is sent in MAX_STREAMS frames. Like MAX_DATA, an updated value is sent when a packet containing the most recent MAX_STREAMS for a stream type frame is declared lost or when the limit is updated, with care taken to prevent the frame from being sent too often.
- * Blocked signals are carried in DATA_BLOCKED, STREAM_DATA_BLOCKED, and STREAMS_BLOCKED frames. DATA_BLOCKED frames have connection scope, STREAM_DATA_BLOCKED frames have stream scope, and STREAMS_BLOCKED frames are scoped to a specific stream type. New frames are sent if packets containing the most recent frame for a scope is lost, but only while the endpoint is blocked on the corresponding limit. These frames always include the limit that is causing blocking at the time that they are transmitted.

- * A liveness or path validation check using `PATH_CHALLENGE` frames is sent periodically until a matching `PATH_RESPONSE` frame is received or until there is no remaining need for liveness or path validation checking. `PATH_CHALLENGE` frames include a different payload each time they are sent.
- * Responses to path validation using `PATH_RESPONSE` frames are sent just once. The peer is expected to send more `PATH_CHALLENGE` frames as necessary to evoke additional `PATH_RESPONSE` frames.
- * New connection IDs are sent in `NEW_CONNECTION_ID` frames and retransmitted if the packet containing them is lost. Retransmissions of this frame carry the same sequence number value. Likewise, retired connection IDs are sent in `RETIRE_CONNECTION_ID` frames and retransmitted if the packet containing them is lost.
- * `NEW_TOKEN` frames are retransmitted if the packet containing them is lost. No special support is made for detecting reordered and duplicated `NEW_TOKEN` frames other than a direct comparison of the frame contents.
- * `PING` and `PADDING` frames contain no information, so lost `PING` or `PADDING` frames do not require repair.
- * The `HANDSHAKE_DONE` frame **MUST** be retransmitted until it is acknowledged.

Endpoints **SHOULD** prioritize retransmission of data over sending new data, unless priorities specified by the application indicate otherwise; see [Section 2.3](#).

Even though a sender is encouraged to assemble frames containing up-to-date information every time it sends a packet, it is not forbidden to retransmit copies of frames from lost packets. A sender that retransmits copies of frames needs to handle decreases in available payload size due to change in packet number length, connection ID length, and path MTU. A receiver **MUST** accept packets containing an outdated frame, such as a `MAX_DATA` frame carrying a smaller maximum data than one found in an older packet.

Upon detecting losses, a sender **MUST** take appropriate congestion control action. The details of loss detection and congestion control are described in [\[QUIC-RECOVERY\]](#).

13.4. Explicit Congestion Notification

QUIC endpoints can use Explicit Congestion Notification (ECN) [RFC3168] to detect and respond to network congestion. ECN allows a network node to indicate congestion in the network by setting a codepoint in the IP header of a packet instead of dropping it. Endpoints react to congestion by reducing their sending rate in response, as described in [QUIC-RECOVERY].

To use ECN, QUIC endpoints first determine whether a path supports ECN marking and the peer is able to access the ECN codepoint in the IP header. A network path does not support ECN if ECN marked packets get dropped or ECN markings are rewritten on the path. An endpoint validates the use of ECN on the path, both during connection establishment and when migrating to a new path (Section 9).

13.4.1. ECN Counts

On receiving a QUIC packet with an ECT or CE codepoint, an ECN-enabled endpoint that can access the ECN codepoints from the enclosing IP packet increases the corresponding ECT(0), ECT(1), or CE count, and includes these counts in subsequent ACK frames; see Section 13.2 and Section 19.3. Note that this requires being able to read the ECN codepoints from the enclosing IP packet, which is not possible on all platforms.

An IP packet that results in no QUIC packets being processed does not increase ECN counts. A QUIC packet detected by a receiver as a duplicate does not affect the receiver's local ECN codepoint counts; see Section 21.8 for relevant security concerns.

If an endpoint receives a QUIC packet without an ECT or CE codepoint in the IP packet header, it responds per Section 13.2 with an ACK frame without increasing any ECN counts. If an endpoint does not implement ECN support or does not have access to received ECN codepoints, it does not increase ECN counts.

Coalesced packets (see Section 12.2) mean that several packets can share the same IP header. The ECN counts for the ECN codepoint received in the associated IP header are incremented once for each QUIC packet, not per enclosing IP packet or UDP datagram.

Each packet number space maintains separate acknowledgement state and separate ECN counts. For example, if one each of an Initial, 0-RTT, Handshake, and 1-RTT QUIC packet are coalesced, the corresponding counts for the Initial and Handshake packet number space will be incremented by one and the counts for the 1-RTT packet number space will be increased by two.

13.4.2. ECN Validation

It is possible for faulty network devices to corrupt or erroneously drop packets with ECN markings. To provide robust connectivity in the presence of such devices, each endpoint independently validates ECN counts and disables ECN if errors are detected.

Endpoints validate ECN for packets sent on each network path independently. An endpoint thus validates ECN on new connection establishment, when switching to a new server preferred address, and on active connection migration to a new path. [Appendix B](#) describes one possible algorithm for testing paths for ECN support.

Even if an endpoint does not use ECN markings on packets it transmits, the endpoint **MUST** provide feedback about ECN markings received from the peer if they are accessible. Failing to report ECN counts will cause the peer to disable ECN marking.

13.4.2.1. Sending ECN Markings

To start ECN validation, an endpoint **SHOULD** do the following when sending packets on a new path to a peer:

- * Set the ECT(0) codepoint in the IP header of early outgoing packets sent on a new path to the peer [[RFC8311](#)].
- * If all packets that were sent with the ECT(0) codepoint are eventually deemed lost [[QUIC-RECOVERY](#)], validation is deemed to have failed.

To reduce the chances of misinterpreting congestive loss as packets dropped by a faulty network element, an endpoint could set the ECT(0) codepoint for only the first ten outgoing packets on a path, or for a period of three RTTs, whichever occurs first.

Other methods of probing paths for ECN support are possible, as are different marking strategies. Implementations **MAY** use other methods defined in RFCs; see [[RFC8311](#)]. Implementations that use the ECT(1) codepoint need to perform ECN validation using ECT(1) counts.

13.4.2.2. Receiving ACK Frames

An endpoint that sets ECT(0) or ECT(1) codepoints on packets it transmits **MUST** use the following steps on receiving an ACK frame to validate ECN.

- * If this ACK frame newly acknowledges a packet that the endpoint sent with either ECT(0) or ECT(1) codepoints set, and if no ECN feedback is present in the ACK frame, validation fails. This step protects against both a network element that zeroes out ECN bits and a peer that is unable to access ECN markings, since the peer could respond without ECN feedback in these cases.
- * For validation to succeed, the total increase in ECT(0), ECT(1), and CE counts MUST be no smaller than the total number of QUIC packets sent with an ECT codepoint that are newly acknowledged in this ACK frame. This step detects any network remarking from ECT(0), ECT(1), or CE codepoints to Not-ECT.
- * Any increase in either ECT(0) or ECT(1) counts, plus any increase in the CE count, MUST be no smaller than the number of packets sent with the corresponding ECT codepoint that are newly acknowledged in this ACK frame. This step detects any erroneous network remarking from ECT(0) to ECT(1) (or vice versa).

Processing ECN counts out of order can result in validation failure. An endpoint SHOULD NOT perform this validation if this ACK frame does not advance the largest packet number acknowledged in this connection.

An endpoint could miss acknowledgements for a packet when ACK frames are lost. It is therefore possible for the total increase in ECT(0), ECT(1), and CE counts to be greater than the number of packets acknowledged in an ACK frame. When this happens, and if validation succeeds, the local reference counts MUST be increased to match the counts in the ACK frame.

13.4.2.3. Validation Outcomes

If validation fails, then the endpoint stops sending ECN markings in subsequent IP packets with the expectation that either the network path or the peer does not support ECN.

Upon successful validation, an endpoint can continue to set ECT codepoints in subsequent packets with the expectation that the path is ECN-capable. Network routing and path elements can change mid-connection however; an endpoint MUST disable ECN if validation fails at any point in the connection.

Even if validation fails, an endpoint MAY revalidate ECN on the same path at any later time in the connection.

14. Packet Size

The QUIC packet size includes the QUIC header and protected payload, but not the UDP or IP headers.

QUIC depends upon a minimum IP packet size of at least 1280 bytes. This is the IPv6 minimum size [RFC8200] and is also supported by most modern IPv4 networks. Assuming the minimum IP header size, this results in a QUIC maximum packet size of 1232 bytes for IPv6 and 1252 bytes for IPv4.

The QUIC maximum packet size is the largest size of QUIC packet that can be sent across a network path using a single packet. Any maximum packet size larger than 1200 bytes can be discovered using Path Maximum Transmission Unit Discovery (PMTUD; see [Section 14.2.1](#)) or Datagram Packetization Layer PMTU Discovery (DPLPMTUD; see [Section 14.3](#)).

Enforcement of the `max_udp_payload_size` transport parameter ([Section 18.2](#)) might act as an additional limit on the maximum packet size. A sender can avoid exceeding this limit, once the value is known. However, prior to learning the value of the transport parameter, endpoints risk datagrams being lost if they send packets larger than the smallest allowed maximum packet size of 1200 bytes.

UDP datagrams MUST NOT be fragmented at the IP layer. In IPv4 [IPv4], the DF bit MUST be set to prevent fragmentation on the path.

14.1. Initial Packet Size

A client MUST expand the payload of all UDP datagrams carrying Initial packets to at least the smallest allowed maximum packet size (1200 bytes) by adding PADDING frames to the Initial packet or by coalescing the Initial packet; see [Section 12.2](#). Sending a UDP datagram of this size ensures that the network path from the client to the server supports a reasonable Path Maximum Transmission Unit (PMTU). This also helps reduce the amplitude of amplification attacks caused by server responses toward an unverified client address; see [Section 8](#).

Datagrams containing Initial packets MAY exceed 1200 bytes if the client believes that the network path and peer both support the size that it chooses.

A server MUST discard an Initial packet that is carried in a UDP datagram with a payload that is less than the smallest allowed maximum packet size of 1200 bytes. A server MAY also immediately close the connection by sending a CONNECTION_CLOSE frame with an error code of PROTOCOL_VIOLATION; see [Section 10.3.1](#).

The server MUST also limit the number of bytes it sends before validating the address of the client; see [Section 8](#).

14.2. Path Maximum Transmission Unit (PMTU)

The Path Maximum Transmission Unit (PMTU) is the maximum size of the entire IP packet including the IP header, UDP header, and UDP payload. The UDP payload includes the QUIC packet header, protected payload, and any authentication fields. The PMTU can depend on path characteristics, and can therefore change over time. The largest UDP payload an endpoint sends at any given time is referred to as the endpoint's maximum packet size.

An endpoint SHOULD use DPLPMTUD ([Section 14.3](#)) or PMTUD ([Section 14.2.1](#)) to determine whether the path to a destination will support a desired maximum packet size without fragmentation. In the absence of these mechanisms, QUIC endpoints SHOULD NOT send IP packets larger than the smallest allowed maximum packet size.

Both DPLPMTUD and PMTUD send IP packets that are larger than the current maximum packet size. We refer to these as PMTU probes. All QUIC packets that are not sent in a PMTU probe SHOULD be sized to fit within the maximum packet size to avoid the packet being fragmented or dropped [[RFC8085](#)].

If a QUIC endpoint determines that the PMTU between any pair of local and remote IP addresses has fallen below the smallest allowed maximum packet size of 1200 bytes, it MUST immediately cease sending QUIC packets, except for those in PMTU probes or those containing CONNECTION_CLOSE frames, on the affected path. An endpoint MAY terminate the connection if an alternative path cannot be found.

Each pair of local and remote addresses could have a different PMTU. QUIC implementations that implement any kind of PMTU discovery therefore SHOULD maintain a maximum packet size for each combination of local and remote IP addresses.

A QUIC implementation MAY be more conservative in computing the maximum packet size to allow for unknown tunnel overheads or IP header options/extensions.

14.2.1. Handling of ICMP Messages by PMTUD

Path Maximum Transmission Unit Discovery (PMTUD; [RFC1191], [RFC8201]) relies on reception of ICMP messages (e.g., IPv6 Packet Too Big messages) that indicate when a packet is dropped because it is larger than the local router MTU. DPLPMTUD can also optionally use these messages. This use of ICMP messages is potentially vulnerable to off-path attacks that successfully guess the addresses used on the path and reduce the PMTU to a bandwidth-inefficient value.

An endpoint MUST ignore an ICMP message that claims the PMTU has decreased below the minimum QUIC packet size.

The requirements for generating ICMP ([RFC1812], [RFC4443]) state that the quoted packet should contain as much of the original packet as possible without exceeding the minimum MTU for the IP version. The size of the quoted packet can actually be smaller, or the information unintelligible, as described in Section 1.1 of [DPLPMTUD].

QUIC endpoints using PMTUD SHOULD validate ICMP messages to protect from off-path injection as specified in [RFC8201] and Section 5.2 of [RFC8085]. This validation SHOULD use the quoted packet supplied in the payload of an ICMP message to associate the message with a corresponding transport connection (see Section 4.6.1 of [DPLPMTUD]). ICMP message validation MUST include matching IP addresses and UDP ports [RFC8085] and, when possible, connection IDs to an active QUIC session. The endpoint SHOULD ignore all ICMP messages that fail validation.

An endpoint MUST NOT increase PMTU based on ICMP messages; see Section 3, clause 6 of [DPLPMTUD]. Any reduction in the QUIC maximum packet size in response to ICMP messages MAY be provisional until QUIC's loss detection algorithm determines that the quoted packet has actually been lost.

14.3. Datagram Packetization Layer PMTU Discovery

Datagram Packetization Layer PMTU Discovery (DPLPMTUD; [DPLPMTUD]) relies on tracking loss or acknowledgment of QUIC packets that are carried in PMTU probes. PMTU probes for DPLPMTUD that use the PADDING frame implement "Probing using padding data", as defined in Section 4.1 of [DPLPMTUD].

Endpoints SHOULD set the initial value of BASE_PMTU (see Section 5.1 of [DPLPMTUD]) to be consistent with the minimum QUIC packet size. The MIN_PLPMTU is the same as the BASE_PMTU.

QUIC endpoints implementing DPLPMTUD maintain a maximum packet size (DPLPMTUD MPS) for each combination of local and remote IP addresses.

14.3.1. DPLPMTUD and Initial Connectivity

From the perspective of DPLPMTUD, QUIC transport is an acknowledged packetization layer (PL). A sender can therefore enter the DPLPMTUD BASE state when the QUIC connection handshake has been completed.

14.3.2. Validating the QUIC Path with DPLPMTUD

QUIC provides an acknowledged PL, therefore a sender does not implement the DPLPMTUD CONFIRMATION_TIMER while in the SEARCH_COMPLETE state; see Section 5.2 of [DPLPMTUD].

14.3.3. Handling of ICMP Messages by DPLPMTUD

An endpoint using DPLPMTUD requires the validation of any received ICMP Packet Too Big (PTB) message before using the PTB information, as defined in Section 4.6 of [DPLPMTUD]. In addition to UDP port validation, QUIC validates an ICMP message by using other PL information (e.g., validation of connection IDs in the quoted packet of any received ICMP message).

The considerations for processing ICMP messages described in Section 14.2.1 also apply if these messages are used by DPLPMTUD.

14.4. Sending QUIC PMTU Probes

PMTU probes are ack-eliciting packets.

Endpoints could limit the content of PMTU probes to PING and PADDING frames as packets that are larger than the current maximum packet size are more likely to be dropped by the network. Loss of a QUIC packet that is carried in a PMTU probe is therefore not a reliable indication of congestion and SHOULD NOT trigger a congestion control reaction; see Section 3, Bullet 7 of [DPLPMTUD]. However, PMTU probes consume congestion window, which could delay subsequent transmission by an application.

14.4.1. PMTU Probes Containing Source Connection ID

Endpoints that rely on the destination connection ID for routing incoming QUIC packets are likely to require that the connection ID be included in PMTU probes to route any resulting ICMP messages ([Section 14.2.1](#)) back to the correct endpoint. However, only long header packets ([Section 17.2](#)) contain the Source Connection ID field, and long header packets are not decrypted or acknowledged by the peer once the handshake is complete.

One way to construct a PMTU probe is to coalesce (see [Section 12.2](#)) a packet with a long header, such as a Handshake or 0-RTT packet ([Section 17.2](#)), with a short header packet in a single UDP datagram. If the resulting PMTU probe reaches the endpoint, the packet with the long header will be ignored, but the short header packet will be acknowledged. If the PMTU probe causes an ICMP message to be sent, the first part of the probe will be quoted in that message. If the Source Connection ID field is within the quoted portion of the probe, that could be used for routing or validation of the ICMP message.

Note: The purpose of using a packet with a long header is only to ensure that the quoted packet contained in the ICMP message contains a Source Connection ID field. This packet does not need to be a valid packet and it can be sent even if there is no current use for packets of that type.

15. Versions

QUIC versions are identified using a 32-bit unsigned number.

The version 0x00000000 is reserved to represent version negotiation. This version of the specification is identified by the number 0x00000001.

Other versions of QUIC might have different properties to this version. The properties of QUIC that are guaranteed to be consistent across all versions of the protocol are described in [\[QUIC-INVARIANTS\]](#).

Version 0x00000001 of QUIC uses TLS as a cryptographic handshake protocol, as described in [\[QUIC-TLS\]](#).

Versions with the most significant 16 bits of the version number cleared are reserved for use in future IETF consensus documents.

Versions that follow the pattern `0x?a?a?a` are reserved for use in forcing version negotiation to be exercised. That is, any version number where the low four bits of all bytes is 1010 (in binary). A client or server MAY advertise support for any of these reserved versions.

Reserved version numbers will never represent a real protocol; a client MAY use one of these version numbers with the expectation that the server will initiate version negotiation; a server MAY advertise support for one of these versions and can expect that clients ignore the value.

[[RFC editor: please remove the remainder of this section before publication.]]

The version number for the final version of this specification (`0x00000001`), is reserved for the version of the protocol that is published as an RFC.

Version numbers used to identify IETF drafts are created by adding the draft number to `0xff000000`. For example, [draft-ietf-quick-transport-13](#) would be identified as `0xff00000D`.

Implementors are encouraged to register version numbers of QUIC that they are using for private experimentation on the GitHub wiki at <https://github.com/quickwg/base-drafts/wiki/QUIC-Versions>.

16. Variable-Length Integer Encoding

QUIC packets and frames commonly use a variable-length encoding for non-negative integer values. This encoding ensures that smaller integer values need fewer bytes to encode.

The QUIC variable-length integer encoding reserves the two most significant bits of the first byte to encode the base 2 logarithm of the integer encoding length in bytes. The integer value is encoded on the remaining bits, in network byte order.

This means that integers are encoded on 1, 2, 4, or 8 bytes and can encode 6, 14, 30, or 62 bit values respectively. Table 4 summarizes the encoding properties.

2Bit	Length	Usable Bits	Range
00	1	6	0-63
01	2	14	0-16383
10	4	30	0-1073741823
11	8	62	0-4611686018427387903

Table 4: Summary of Integer Encodings

For example, the eight byte sequence c2 19 7c 5e ff 14 e8 8c (in hexadecimal) decodes to the decimal value 151288809941952652; the four byte sequence 9d 7f 3e 7d decodes to 494878333; the two byte sequence 7b bd decodes to 15293; and the single byte 25 decodes to 37 (as does the two byte sequence 40 25).

Error codes ([Section 20](#)) and versions ([Section 15](#)) are described using integers, but do not use this encoding.

17. Packet Formats

All numeric values are encoded in network byte order (that is, big-endian) and all field sizes are in bits. Hexadecimal notation is used for describing the value of fields.

17.1. Packet Number Encoding and Decoding

Packet numbers are integers in the range 0 to $2^{62}-1$ ([Section 12.3](#)). When present in long or short packet headers, they are encoded in 1 to 4 bytes. The number of bits required to represent the packet number is reduced by including the least significant bits of the packet number.

The encoded packet number is protected as described in [Section 5.4](#) of [\[QUIC-TLS\]](#).

The sender **MUST** use a packet number size able to represent more than twice as large a range than the difference between the largest acknowledged packet and packet number being sent. A peer receiving the packet will then correctly decode the packet number, unless the packet is delayed in transit such that it arrives after many higher-numbered packets have been received. An endpoint **SHOULD** use a large enough packet number encoding to allow the packet number to be recovered even if the packet arrives after packets that are sent afterwards.

As a result, the size of the packet number encoding is at least one bit more than the base-2 logarithm of the number of contiguous unacknowledged packet numbers, including the new packet.

For example, if an endpoint has received an acknowledgment for packet 0xab8bc, sending a packet with a number of 0xac5c02 requires a packet number encoding with 16 bits or more; whereas the 24-bit packet number encoding is needed to send a packet with a number of 0xace8fe.

At a receiver, protection of the packet number is removed prior to recovering the full packet number. The full packet number is then reconstructed based on the number of significant bits present, the value of those bits, and the largest packet number received on a successfully authenticated packet. Recovering the full packet number is necessary to successfully remove packet protection.

Once header protection is removed, the packet number is decoded by finding the packet number value that is closest to the next expected packet. The next expected packet is the highest received packet number plus one. For example, if the highest successfully authenticated packet had a packet number of 0xa82f30ea, then a packet containing a 16-bit value of 0x9b32 will be decoded as 0xa82f9b32. Example pseudo-code for packet number decoding can be found in [Appendix A](#).

17.2. Long Header Packets

```
Long Header Packet {
  Header Form (1) = 1,
  Fixed Bit (1) = 1,
  Long Packet Type (2),
  Type-Specific Bits (4),
  Version (32),
  Destination Connection ID Length (8),
  Destination Connection ID (0..160),
  Source Connection ID Length (8),
  Source Connection ID (0..160),
}
```

Figure 12: Long Header Packet Format

Long headers are used for packets that are sent prior to the establishment of 1-RTT keys. Once 1-RTT keys are available, a sender switches to sending packets using the short header ([Section 17.3](#)). The long form allows for special packets - such as the Version Negotiation packet - to be represented in this uniform fixed-length packet format. Packets that use the long header contain the following fields:

Header Form: The most significant bit (0x80) of byte 0 (the first byte) is set to 1 for long headers.

Fixed Bit: The next bit (0x40) of byte 0 is set to 1. Packets containing a zero value for this bit are not valid packets in this version and **MUST** be discarded.

Long Packet Type: The next two bits (those with a mask of 0x30) of byte 0 contain a packet type. Packet types are listed in Table 5.

Type-Specific Bits: The lower four bits (those with a mask of 0x0f) of byte 0 are type-specific.

Version: The QUIC Version is a 32-bit field that follows the first byte. This field indicates which version of QUIC is in use and determines how the rest of the protocol fields are interpreted.

Destination Connection ID Length: The byte following the version contains the length in bytes of the Destination Connection ID field that follows it. This length is encoded as an 8-bit unsigned integer. In QUIC version 1, this value **MUST NOT** exceed 20. Endpoints that receive a version 1 long header with a value larger than 20 **MUST** drop the packet. Servers **SHOULD** be able to read longer connection IDs from other QUIC versions in order to properly form a version negotiation packet.

Destination Connection ID: The Destination Connection ID field follows the Destination Connection ID Length field and is between 0 and 20 bytes in length. [Section 7.2](#) describes the use of this field in more detail.

Source Connection ID Length: The byte following the Destination Connection ID contains the length in bytes of the Source Connection ID field that follows it. This length is encoded as a 8-bit unsigned integer. In QUIC version 1, this value **MUST NOT** exceed 20 bytes. Endpoints that receive a version 1 long header with a value larger than 20 **MUST** drop the packet. Servers **SHOULD** be able to read longer connection IDs from other QUIC versions in order to properly form a version negotiation packet.

Source Connection ID: The Source Connection ID field follows the Source Connection ID Length field and is between 0 and 20 bytes in length. [Section 7.2](#) describes the use of this field in more detail.

In this version of QUIC, the following packet types with the long header are defined:

Type	Name	Section
0x0	Initial	Section 17.2.2
0x1	0-RTT	Section 17.2.3
0x2	Handshake	Section 17.2.4
0x3	Retry	Section 17.2.5

Table 5: Long Header Packet Types

The header form bit, connection ID lengths byte, Destination and Source Connection ID fields, and Version fields of a long header packet are version-independent. The other fields in the first byte are version-specific. See [\[QUIC-INVARIANTS\]](#) for details on how packets from different versions of QUIC are interpreted.

The interpretation of the fields and the payload are specific to a version and packet type. While type-specific semantics for this version are described in the following sections, several long-header packets in this version of QUIC contain these additional fields:

Reserved Bits: Two bits (those with a mask of 0x0c) of byte 0 are

reserved across multiple packet types. These bits are protected using header protection; see Section 5.4 of [QUIC-TLS]. The value included prior to protection MUST be set to 0. An endpoint MUST treat receipt of a packet that has a non-zero value for these bits, after removing both packet and header protection, as a connection error of type `PROTOCOL_VIOLATION`. Discarding such a packet after only removing header protection can expose the endpoint to attacks; see Section 9.3 of [QUIC-TLS].

Packet Number Length: In packet types which contain a Packet Number field, the least significant two bits (those with a mask of `0x03`) of byte 0 contain the length of the packet number, encoded as an unsigned, two-bit integer that is one less than the length of the packet number field in bytes. That is, the length of the packet number field is the value of this field, plus one. These bits are protected using header protection; see Section 5.4 of [QUIC-TLS].

Length: The length of the remainder of the packet (that is, the Packet Number and Payload fields) in bytes, encoded as a variable-length integer (Section 16).

Packet Number: The packet number field is 1 to 4 bytes long. The packet number has confidentiality protection separate from packet protection, as described in Section 5.4 of [QUIC-TLS]. The length of the packet number field is encoded in the Packet Number Length bits of byte 0; see above.

17.2.1. Version Negotiation Packet

A Version Negotiation packet is inherently not version-specific. Upon receipt by a client, it will be identified as a Version Negotiation packet based on the Version field having a value of 0.

The Version Negotiation packet is a response to a client packet that contains a version that is not supported by the server, and is only sent by servers.

The layout of a Version Negotiation packet is:

```
Version Negotiation Packet {
  Header Form (1) = 1,
  Unused (7),
  Version (32) = 0,
  Destination Connection ID Length (8),
  Destination Connection ID (0..2040),
  Source Connection ID Length (8),
  Source Connection ID (0..2040),
  Supported Version (32) ...,
}
```

Figure 13: Version Negotiation Packet

The value in the Unused field is selected randomly by the server. Clients MUST ignore the value of this field. Servers SHOULD set the most significant bit of this field (0x40) to 1 so that Version Negotiation packets appear to have the Fixed Bit field.

The Version field of a Version Negotiation packet MUST be set to 0x00000000.

The server MUST include the value from the Source Connection ID field of the packet it receives in the Destination Connection ID field. The value for Source Connection ID MUST be copied from the Destination Connection ID of the received packet, which is initially randomly selected by a client. Echoing both connection IDs gives clients some assurance that the server received the packet and that the Version Negotiation packet was not generated by an off-path attacker.

As future versions of QUIC may support Connection IDs larger than the version 1 limit, Version Negotiation packets could carry Connection IDs that are longer than 20 bytes.

The remainder of the Version Negotiation packet is a list of 32-bit versions which the server supports.

A Version Negotiation packet cannot be explicitly acknowledged in an ACK frame by a client. Receiving another Initial packet implicitly acknowledges a Version Negotiation packet.

The Version Negotiation packet does not include the Packet Number and Length fields present in other packets that use the long header form. Consequently, a Version Negotiation packet consumes an entire UDP datagram.

A server MUST NOT send more than one Version Negotiation packet in response to a single UDP datagram.

See [Section 6](#) for a description of the version negotiation process.

17.2.2. Initial Packet

An Initial packet uses long headers with a type value of 0x0. It carries the first CRYPTO frames sent by the client and server to perform key exchange, and carries ACKs in either direction.

```
Initial Packet {
  Header Form (1) = 1,
  Fixed Bit (1) = 1,
  Long Packet Type (2) = 0,
  Reserved Bits (2),
  Packet Number Length (2),
  Version (32),
  Destination Connection ID Length (8),
  Destination Connection ID (0..160),
  Source Connection ID Length (8),
  Source Connection ID (0..160),
  Token Length (i),
  Token (...),
  Length (i),
  Packet Number (8..32),
  Packet Payload (...),
}
```

Figure 14: Initial Packet

The Initial packet contains a long header as well as the Length and Packet Number fields. The first byte contains the Reserved and Packet Number Length bits. Between the Source Connection ID and Length fields, there are two additional fields specific to the Initial packet.

Token Length: A variable-length integer specifying the length of the Token field, in bytes. This value is zero if no token is present. Initial packets sent by the server MUST set the Token Length field to zero; clients that receive an Initial packet with a non-zero Token Length field MUST either discard the packet or generate a connection error of type `PROTOCOL_VIOLATION`.

Token: The value of the token that was previously provided in a Retry packet or `NEW_TOKEN` frame.

Packet Payload: The payload of the packet.

In order to prevent tampering by version-unaware middleboxes, Initial packets are protected with connection- and version-specific keys (Initial keys) as described in [QUIC-TLS]. This protection does not provide confidentiality or integrity against on-path attackers, but provides some level of protection against off-path attackers.

The client and server use the Initial packet type for any packet that contains an initial cryptographic handshake message. This includes all cases where a new packet containing the initial cryptographic message needs to be created, such as the packets sent after receiving a Retry packet (Section 17.2.5).

A server sends its first Initial packet in response to a client Initial. A server may send multiple Initial packets. The cryptographic key exchange could require multiple round trips or retransmissions of this data.

The payload of an Initial packet includes a CRYPTO frame (or frames) containing a cryptographic handshake message, ACK frames, or both. PING, PADDING, and CONNECTION_CLOSE frames of type 0x1c are also permitted. An endpoint that receives an Initial packet containing other frames can either discard the packet as spurious or treat it as a connection error.

The first packet sent by a client always includes a CRYPTO frame that contains the start or all of the first cryptographic handshake message. The first CRYPTO frame sent always begins at an offset of 0; see Section 7.

Note that if the server sends a HelloRetryRequest, the client will send another series of Initial packets. These Initial packets will continue the cryptographic handshake and will contain CRYPTO frames starting at an offset matching the size of the CRYPTO frames sent in the first flight of Initial packets.

17.2.2.1. Abandoning Initial Packets

A client stops both sending and processing Initial packets when it sends its first Handshake packet. A server stops sending and processing Initial packets when it receives its first Handshake packet. Though packets might still be in flight or awaiting acknowledgment, no further Initial packets need to be exchanged beyond this point. Initial packet protection keys are discarded (see Section 4.11.1 of [QUIC-TLS]) along with any loss recovery and congestion control state; see Section 6.4 of [QUIC-RECOVERY].

Any data in CRYPTO frames is discarded - and no longer retransmitted - when Initial keys are discarded.

17.2.3. 0-RTT

A 0-RTT packet uses long headers with a type value of 0x1, followed by the Length and Packet Number fields. The first byte contains the Reserved and Packet Number Length bits. It is used to carry "early" data from the client to the server as part of the first flight, prior to handshake completion. As part of the TLS handshake, the server can accept or reject this early data.

See Section 2.3 of [TLS13] for a discussion of 0-RTT data and its limitations.

```
0-RTT Packet {
  Header Form (1) = 1,
  Fixed Bit (1) = 1,
  Long Packet Type (2) = 1,
  Reserved Bits (2),
  Packet Number Length (2),
  Version (32),
  Destination Connection ID Length (8),
  Destination Connection ID (0..160),
  Source Connection ID Length (8),
  Source Connection ID (0..160),
  Length (i),
  Packet Number (8..32),
  Packet Payload (...),
}
```

Figure 15: 0-RTT Packet

Packet numbers for 0-RTT protected packets use the same space as 1-RTT protected packets.

After a client receives a Retry packet, 0-RTT packets are likely to have been lost or discarded by the server. A client SHOULD attempt to resend data in 0-RTT packets after it sends a new Initial packet.

A client MUST NOT reset the packet number it uses for 0-RTT packets, since the keys used to protect 0-RTT packets will not change as a result of responding to a Retry packet. Sending packets with the same packet number in that case is likely to compromise the packet protection for all 0-RTT packets because the same key and nonce could be used to protect different content.

A client only receives acknowledgments for its 0-RTT packets once the handshake is complete. Consequently, a server might expect 0-RTT packets to start with a packet number of 0. Therefore, in determining the length of the packet number encoding for 0-RTT

packets, a client MUST assume that all packets up to the current packet number are in flight, starting from a packet number of 0. Thus, 0-RTT packets could need to use a longer packet number encoding.

A client MUST NOT send 0-RTT packets once it starts processing 1-RTT packets from the server. This means that 0-RTT packets cannot contain any response to frames from 1-RTT packets. For instance, a client cannot send an ACK frame in a 0-RTT packet, because that can only acknowledge a 1-RTT packet. An acknowledgment for a 1-RTT packet MUST be carried in a 1-RTT packet.

A server SHOULD treat a violation of remembered limits as a connection error of an appropriate type (for instance, a `FLOW_CONTROL_ERROR` for exceeding stream data limits).

17.2.4. Handshake Packet

A Handshake packet uses long headers with a type value of 0x2, followed by the Length and Packet Number fields. The first byte contains the Reserved and Packet Number Length bits. It is used to carry acknowledgments and cryptographic handshake messages from the server and client.

```
Handshake Packet {
  Header Form (1) = 1,
  Fixed Bit (1) = 1,
  Long Packet Type (2) = 2,
  Reserved Bits (2),
  Packet Number Length (2),
  Version (32),
  Destination Connection ID Length (8),
  Destination Connection ID (0..160),
  Source Connection ID Length (8),
  Source Connection ID (0..160),
  Length (i),
  Packet Number (8..32),
  Packet Payload (..),
}
```

Figure 16: Handshake Protected Packet

Once a client has received a Handshake packet from a server, it uses Handshake packets to send subsequent cryptographic handshake messages and acknowledgments to the server.

The Destination Connection ID field in a Handshake packet contains a connection ID that is chosen by the recipient of the packet; the Source Connection ID includes the connection ID that the sender of the packet wishes to use; see [Section 7.2](#).

Handshake packets are their own packet number space, and thus the first Handshake packet sent by a server contains a packet number of 0.

The payload of this packet contains CRYPTO frames and could contain PING, PADDING, or ACK frames. Handshake packets MAY contain CONNECTION_CLOSE frames of type 0x1c. Endpoints MUST treat receipt of Handshake packets with other frames as a connection error.

Like Initial packets (see [Section 17.2.2.1](#)), data in CRYPTO frames for Handshake packets is discarded - and no longer retransmitted - when Handshake protection keys are discarded.

17.2.5. Retry Packet

A Retry packet uses a long packet header with a type value of 0x3. It carries an address validation token created by the server. It is used by a server that wishes to perform a retry; see [Section 8.1](#).

```
Retry Packet {
  Header Form (1) = 1,
  Fixed Bit (1) = 1,
  Long Packet Type (2) = 3,
  Unused (4),
  Version (32),
  Destination Connection ID Length (8),
  Destination Connection ID (0..160),
  Source Connection ID Length (8),
  Source Connection ID (0..160),
  Retry Token (...),
  Retry Integrity Tag (128),
}
```

Figure 17: Retry Packet

A Retry packet (shown in Figure 17) does not contain any protected fields. The value in the Unused field is selected randomly by the server. In addition to the fields from the long header, it contains these additional fields:

Retry Token: An opaque token that the server can use to validate the client's address.

Retry Integrity Tag: See the Retry Packet Integrity section of [QUIC-TLS].

17.2.5.1. Sending a Retry Packet

The server populates the Destination Connection ID with the connection ID that the client included in the Source Connection ID of the Initial packet.

The server includes a connection ID of its choice in the Source Connection ID field. This value MUST NOT be equal to the Destination Connection ID field of the packet sent by the client. A client MUST discard a Retry packet that contains a Source Connection ID field that is identical to the Destination Connection ID field of its Initial packet. The client MUST use the value from the Source Connection ID field of the Retry packet in the Destination Connection ID field of subsequent packets that it sends.

A server MAY send Retry packets in response to Initial and 0-RTT packets. A server can either discard or buffer 0-RTT packets that it receives. A server can send multiple Retry packets as it receives Initial or 0-RTT packets. A server MUST NOT send more than one Retry packet in response to a single UDP datagram.

17.2.5.2. Handling a Retry Packet

A client MUST accept and process at most one Retry packet for each connection attempt. After the client has received and processed an Initial or Retry packet from the server, it MUST discard any subsequent Retry packets that it receives.

Clients MUST discard Retry packets that have a Retry Integrity Tag that cannot be validated, see the Retry Packet Integrity section of [QUIC-TLS]. This diminishes an off-path attacker's ability to inject a Retry packet and protects against accidental corruption of Retry packets. A client MUST discard a Retry packet with a zero-length Retry Token field.

The client responds to a Retry packet with an Initial packet that includes the provided Retry Token to continue connection establishment.

A client sets the Destination Connection ID field of this Initial packet to the value from the Source Connection ID in the Retry packet. Changing Destination Connection ID also results in a change to the keys used to protect the Initial packet. It also sets the Token field to the token provided in the Retry. The client **MUST NOT** change the Source Connection ID because the server could include the connection ID as part of its token validation logic; see [Section 8.1.4](#).

A Retry packet does not include a packet number and cannot be explicitly acknowledged by a client.

17.2.5.3. Continuing a Handshake After Retry

The next Initial packet from the client uses the connection ID and token values from the Retry packet; see [Section 7.2](#). Aside from this, the Initial packet sent by the client is subject to the same restrictions as the first Initial packet. A client **MUST** use the same cryptographic handshake message it includes in this packet. A server **MAY** treat a packet that contains a different cryptographic handshake message as a connection error or discard it.

A client **MAY** attempt 0-RTT after receiving a Retry packet by sending 0-RTT packets to the connection ID provided by the server. A client **MUST NOT** change the cryptographic handshake message it sends in response to receiving a Retry.

A client **MUST NOT** reset the packet number for any packet number space after processing a Retry packet; [Section 17.2.3](#) contains more information on this.

A server acknowledges the use of a Retry packet for a connection using the `retry_source_connection_id` transport parameter; see [Section 18.2](#). If the server sends a Retry packet, it also subsequently includes the value of the Source Connection ID field from the Retry packet in its `retry_source_connection_id` transport parameter.

If the client received and processed a Retry packet, it **MUST** validate that the `retry_source_connection_id` transport parameter is present and correct; otherwise, it **MUST** validate that the transport parameter is absent. A client **MUST** treat a failed validation as a connection error of type `PROTOCOL_VIOLATION`.

17.3. Short Header Packets

This version of QUIC defines a single packet type which uses the short packet header.

```
Short Header Packet {
  Header Form (1) = 0,
  Fixed Bit (1) = 1,
  Spin Bit (1),
  Reserved Bits (2),
  Key Phase (1),
  Packet Number Length (2),
  Destination Connection ID (0..160),
  Packet Number (8..32),
  Packet Payload (...),
}
```

Figure 18: Short Header Packet Format

The short header can be used after the version and 1-RTT keys are negotiated. Packets that use the short header contain the following fields:

Header Form: The most significant bit (0x80) of byte 0 is set to 0 for the short header.

Fixed Bit: The next bit (0x40) of byte 0 is set to 1. Packets containing a zero value for this bit are not valid packets in this version and MUST be discarded.

Spin Bit: The third most significant bit (0x20) of byte 0 is the latency spin bit, set as described in [Section 17.3.1](#).

Reserved Bits: The next two bits (those with a mask of 0x18) of byte 0 are reserved. These bits are protected using header protection; see Section 5.4 of [\[QUIC-TLS\]](#). The value included prior to protection MUST be set to 0. An endpoint MUST treat receipt of a packet that has a non-zero value for these bits, after removing both packet and header protection, as a connection error of type `PROTOCOL_VIOLATION`. Discarding such a packet after only removing header protection can expose the endpoint to attacks; see Section 9.3 of [\[QUIC-TLS\]](#).

Key Phase: The next bit (0x04) of byte 0 indicates the key phase, which allows a recipient of a packet to identify the packet protection keys that are used to protect the packet. See [\[QUIC-TLS\]](#) for details. This bit is protected using header protection; see Section 5.4 of [\[QUIC-TLS\]](#).

Packet Number Length: The least significant two bits (those with a mask of 0x03) of byte 0 contain the length of the packet number, encoded as an unsigned, two-bit integer that is one less than the length of the packet number field in bytes. That is, the length

of the packet number field is the value of this field, plus one. These bits are protected using header protection; see Section 5.4 of [QUIC-TLS].

Destination Connection ID: The Destination Connection ID is a connection ID that is chosen by the intended recipient of the packet. See Section 5.1 for more details.

Packet Number: The packet number field is 1 to 4 bytes long. The packet number has confidentiality protection separate from packet protection, as described in Section 5.4 of [QUIC-TLS]. The length of the packet number field is encoded in Packet Number Length field. See Section 17.1 for details.

Packet Payload: Packets with a short header always include a 1-RTT protected payload.

The header form bit and the connection ID field of a short header packet are version-independent. The remaining fields are specific to the selected QUIC version. See [QUIC-INVARIANTS] for details on how packets from different versions of QUIC are interpreted.

17.3.1. Latency Spin Bit

The latency spin bit enables passive latency monitoring from observation points on the network path throughout the duration of a connection. The spin bit is only present in the short packet header, since it is possible to measure the initial RTT of a connection by observing the handshake. Therefore, the spin bit is available after version negotiation and connection establishment are completed. On-path measurement and use of the latency spin bit is further discussed in [QUIC-MANAGEABILITY].

The spin bit is an OPTIONAL feature of QUIC. A QUIC stack that chooses to support the spin bit MUST implement it as specified in this section.

Each endpoint unilaterally decides if the spin bit is enabled or disabled for a connection. Implementations MUST allow administrators of clients and servers to disable the spin bit either globally or on a per-connection basis. Even when the spin bit is not disabled by the administrator, endpoints MUST disable their use of the spin bit for a random selection of at least one in every 16 network paths, or for one in every 16 connection IDs. As each endpoint disables the spin bit independently, this ensures that the spin bit signal is disabled on approximately one in eight network paths.

When the spin bit is disabled, endpoints MAY set the spin bit to any value, and MUST ignore any incoming value. It is RECOMMENDED that endpoints set the spin bit to a random value either chosen independently for each packet or chosen independently for each connection ID.

If the spin bit is enabled for the connection, the endpoint maintains a spin value and sets the spin bit in the short header to the currently stored value when a packet with a short header is sent out. The spin value is initialized to 0 in the endpoint at connection start. Each endpoint also remembers the highest packet number seen from its peer on the connection.

When a server receives a short header packet that increments the highest packet number seen by the server from the client, it sets the spin value to be equal to the spin bit in the received packet.

When a client receives a short header packet that increments the highest packet number seen by the client from the server, it sets the spin value to the inverse of the spin bit in the received packet.

An endpoint resets its spin value to zero when sending the first packet of a given connection with a new connection ID. This reduces the risk that transient spin bit state can be used to link flows across connection migration or ID change.

With this mechanism, the server reflects the spin value received, while the client 'spins' it after one RTT. On-path observers can measure the time between two spin bit toggle events to estimate the end-to-end RTT of a connection.

18. Transport Parameter Encoding

The `extension_data` field of the `quic_transport_parameters` extension defined in [QUIC-TLS] contains the QUIC transport parameters. They are encoded as a sequence of transport parameters, as shown in Figure 19:

```
Transport Parameters {  
  Transport Parameter (...) ...,  
}
```

Figure 19: Sequence of Transport Parameters

Each transport parameter is encoded as an (identifier, length, value) tuple, as shown in Figure 20:

```
Transport Parameter {  
  Transport Parameter ID (i),  
  Transport Parameter Length (i),  
  Transport Parameter Value (..),  
}
```

Figure 20: Transport Parameter Encoding

The Transport Parameter Length field contains the length of the Transport Parameter Value field.

QUIC encodes transport parameters into a sequence of bytes, which are then included in the cryptographic handshake.

18.1. Reserved Transport Parameters

Transport parameters with an identifier of the form "31 * N + 27" for integer values of N are reserved to exercise the requirement that unknown transport parameters be ignored. These transport parameters have no semantics, and may carry arbitrary values.

18.2. Transport Parameter Definitions

This section details the transport parameters defined in this document.

Many transport parameters listed here have integer values. Those transport parameters that are identified as integers use a variable-length integer encoding; see [Section 16](#). Transport parameters have a default value of 0 if the transport parameter is absent unless otherwise stated.

The following transport parameters are defined:

`original_destination_connection_id` (0x00): The value of the Destination Connection ID field from the first Initial packet sent by the client; see [Section 7.3](#). This transport parameter is only sent by a server.

`max_idle_timeout` (0x01): The max idle timeout is a value in milliseconds that is encoded as an integer; see ([Section 10.2](#)). Idle timeout is disabled when both endpoints omit this transport parameter or specify a value of 0.

`stateless_reset_token` (0x02): A stateless reset token is used in verifying a stateless reset; see [Section 10.4](#). This parameter is a sequence of 16 bytes. This transport parameter MUST NOT be sent by a client, but MAY be sent by a server. A server that does not

send this transport parameter cannot use stateless reset ([Section 10.4](#)) for the connection ID negotiated during the handshake.

`max_udp_payload_size (0x03)`: The maximum UDP payload size parameter is an integer value that limits the size of UDP payloads that the endpoint is willing to receive. UDP packets with payloads larger than this limit are not likely to be processed by the receiver.

The default for this parameter is the maximum permitted UDP payload of 65527. Values below 1200 are invalid.

This limit does act as an additional constraint on datagram size in the same way as the path MTU, but it is a property of the endpoint and not the path; see [Section 14](#). It is expected that this is the space an endpoint dedicates to holding incoming packets.

`initial_max_data (0x04)`: The initial maximum data parameter is an integer value that contains the initial value for the maximum amount of data that can be sent on the connection. This is equivalent to sending a `MAX_DATA` ([Section 19.9](#)) for the connection immediately after completing the handshake.

`initial_max_stream_data_bidi_local (0x05)`: This parameter is an integer value specifying the initial flow control limit for locally-initiated bidirectional streams. This limit applies to newly created bidirectional streams opened by the endpoint that sends the transport parameter. In client transport parameters, this applies to streams with an identifier with the least significant two bits set to 0x0; in server transport parameters, this applies to streams with the least significant two bits set to 0x1.

`initial_max_stream_data_bidi_remote (0x06)`: This parameter is an integer value specifying the initial flow control limit for peer-initiated bidirectional streams. This limit applies to newly created bidirectional streams opened by the endpoint that receives the transport parameter. In client transport parameters, this applies to streams with an identifier with the least significant two bits set to 0x1; in server transport parameters, this applies to streams with the least significant two bits set to 0x0.

`initial_max_stream_data_uni (0x07)`: This parameter is an integer value specifying the initial flow control limit for unidirectional streams. This limit applies to newly created unidirectional streams opened by the endpoint that receives the transport parameter. In client transport parameters, this applies to

streams with an identifier with the least significant two bits set to 0x3; in server transport parameters, this applies to streams with the least significant two bits set to 0x2.

`initial_max_streams_bidi (0x08)`: The initial maximum bidirectional streams parameter is an integer value that contains the initial maximum number of bidirectional streams the peer may initiate. If this parameter is absent or zero, the peer cannot open bidirectional streams until a `MAX_STREAMS` frame is sent. Setting this parameter is equivalent to sending a `MAX_STREAMS` ([Section 19.11](#)) of the corresponding type with the same value.

`initial_max_streams_uni (0x09)`: The initial maximum unidirectional streams parameter is an integer value that contains the initial maximum number of unidirectional streams the peer may initiate. If this parameter is absent or zero, the peer cannot open unidirectional streams until a `MAX_STREAMS` frame is sent. Setting this parameter is equivalent to sending a `MAX_STREAMS` ([Section 19.11](#)) of the corresponding type with the same value.

`ack_delay_exponent (0x0a)`: The ACK delay exponent is an integer value indicating an exponent used to decode the ACK Delay field in the ACK frame ([Section 19.3](#)). If this value is absent, a default value of 3 is assumed (indicating a multiplier of 8). Values above 20 are invalid.

`max_ack_delay (0x0b)`: The maximum ACK delay is an integer value indicating the maximum amount of time in milliseconds by which the endpoint will delay sending acknowledgments. This value SHOULD include the receiver's expected delays in alarms firing. For example, if a receiver sets a timer for 5ms and alarms commonly fire up to 1ms late, then it should send a `max_ack_delay` of 6ms. If this value is absent, a default of 25 milliseconds is assumed. Values of 2^{14} or greater are invalid.

`disable_active_migration (0x0c)`: The disable active migration transport parameter is included if the endpoint does not support active connection migration ([Section 9](#)) on the address being used during the handshake. When a peer sets this transport parameter, an endpoint MUST NOT use a new local address when sending to the address that the peer used during the handshake. This transport parameter does not prohibit connection migration after a client has acted on a `preferred_address` transport parameter. This parameter is a zero-length value.

`preferred_address (0x0d)`: The server's preferred address is used to

effect a change in server address at the end of the handshake, as described in [Section 9.6](#). The format of this transport parameter is shown in Figure 21. This transport parameter is only sent by a server. Servers MAY choose to only send a preferred address of one address family by sending an all-zero address and port (0.0.0.0:0 or ::.0) for the other family. IP addresses are encoded in network byte order.

The Connection ID field and the Stateless Reset Token field contain an alternative connection ID that has a sequence number of 1; see [Section 5.1.1](#). Having these values bundled with the preferred address ensures that there will be at least one unused active connection ID when the client initiates migration to the preferred address.

The Connection ID and Stateless Reset Token fields of a preferred address are identical in syntax and semantics to the corresponding fields of a NEW_CONNECTION_ID frame ([Section 19.15](#)). A server that chooses a zero-length connection ID MUST NOT provide a preferred address. Similarly, a server MUST NOT include a zero-length connection ID in this transport parameter. A client MUST treat violation of these requirements as a connection error of type TRANSPORT_PARAMETER_ERROR.

```
Preferred Address {
  IPv4 Address (32),
  IPv4 Port (16),
  IPv6 Address (128),
  IPv6 Port (16),
  CID Length (8),
  Connection ID (..),
  Stateless Reset Token (128),
}
```

Figure 21: Preferred Address format

active_connection_id_limit (0x0e): The active connection ID limit is

an integer value specifying the maximum number of connection IDs from the peer that an endpoint is willing to store. This value includes the connection ID received during the handshake, that received in the `preferred_address` transport parameter, and those received in `NEW_CONNECTION_ID` frames. The value of the `active_connection_id_limit` parameter MUST be at least 2. An endpoint that receives a value less than 2 MUST close the connection with an error of type `TRANSPORT_PARAMETER_ERROR`. If this transport parameter is absent, a default of 2 is assumed. If an endpoint issues a zero-length connection ID, it will never send a `NEW_CONNECTION_ID` frame and therefore ignores the `active_connection_id_limit` value received from its peer.

`initial_source_connection_id` (0x0f): The value that the endpoint included in the Source Connection ID field of the first Initial packet it sends for the connection; see [Section 7.3](#).

`retry_source_connection_id` (0x10): The value that the server included in the Source Connection ID field of a Retry packet; see [Section 7.3](#). This transport parameter is only sent by a server.

If present, transport parameters that set initial flow control limits (`initial_max_stream_data_bidi_local`, `initial_max_stream_data_bidi_remote`, and `initial_max_stream_data_uni`) are equivalent to sending a `MAX_STREAM_DATA` frame ([Section 19.10](#)) on every stream of the corresponding type immediately after opening. If the transport parameter is absent, streams of that type start with a flow control limit of 0.

A client MUST NOT include any server-only transport parameter: `original_destination_connection_id`, `preferred_address`, `retry_source_connection_id`, or `stateless_reset_token`. A server MUST treat receipt of any of these transport parameters as a connection error of type `TRANSPORT_PARAMETER_ERROR`.

19. Frame Types and Formats

As described in [Section 12.4](#), packets contain one or more frames. This section describes the format and semantics of the core QUIC frame types.

19.1. PADDING Frame

The PADDING frame (type=0x00) has no semantic value. PADDING frames can be used to increase the size of a packet. Padding can be used to increase an initial client packet to the minimum required size, or to provide protection against traffic analysis for protected packets.

As shown in Figure 22, a PADDING frame has no content. That is, a PADDING frame consists of the single byte that identifies the frame as a PADDING frame.

```
PADDING Frame {  
  Type (i) = 0x00,  
}
```

Figure 22: PADDING Frame Format

19.2. PING Frame

Endpoints can use PING frames (type=0x01) to verify that their peers are still alive or to check reachability to the peer. As shown in Figure 23 a PING frame contains no content.

```
PING Frame {  
  Type (i) = 0x01,  
}
```

Figure 23: PING Frame Format

The receiver of a PING frame simply needs to acknowledge the packet containing this frame.

The PING frame can be used to keep a connection alive when an application or application protocol wishes to prevent the connection from timing out; see [Section 10.2.2](#).

19.3. ACK Frames

Receivers send ACK frames (types 0x02 and 0x03) to inform senders of packets they have received and processed. The ACK frame contains one or more ACK Ranges. ACK Ranges identify acknowledged packets. If the frame type is 0x03, ACK frames also contain the sum of QUIC packets with associated ECN marks received on the connection up until this point. QUIC implementations MUST properly handle both types and, if they have enabled ECN for packets they send, they SHOULD use the information in the ECN section to manage their congestion state.

QUIC acknowledgements are irrevocable. Once acknowledged, a packet remains acknowledged, even if it does not appear in a future ACK frame. This is unlike TCP SACKs ([RFC2018]).

Packets from different packet number spaces can be identified using the same numeric value. An acknowledgment for a packet needs to indicate both a packet number and a packet number space. This is accomplished by having each ACK frame only acknowledge packet numbers in the same space as the packet in which the ACK frame is contained.

Version Negotiation and Retry packets cannot be acknowledged because they do not contain a packet number. Rather than relying on ACK frames, these packets are implicitly acknowledged by the next Initial packet sent by the client.

An ACK frame is shown in Figure 24.

```
ACK Frame {
  Type (i) = 0x02..0x03,
  Largest Acknowledged (i),
  ACK Delay (i),
  ACK Range Count (i),
  First ACK Range (i),
  ACK Range (...) ...,
  [ECN Counts (...)],
}
```

Figure 24: ACK Frame Format

ACK frames contain the following fields:

Largest Acknowledged: A variable-length integer representing the largest packet number the peer is acknowledging; this is usually the largest packet number that the peer has received prior to generating the ACK frame. Unlike the packet number in the QUIC long or short header, the value in an ACK frame is not truncated.

ACK Delay: A variable-length integer representing the time delta in microseconds between when this ACK was sent and when the largest acknowledged packet, as indicated in the Largest Acknowledged field, was received by this peer. The value of the ACK Delay field is scaled by multiplying the encoded value by 2 to the power of the value of the `ack_delay_exponent` transport parameter set by the sender of the ACK frame; see [Section 18.2](#). Scaling in this fashion allows for a larger range of values with a shorter encoding at the cost of lower resolution. Because the receiver doesn't use the ACK Delay for Initial and Handshake packets, a sender SHOULD send a value of 0.

ACK Range Count: A variable-length integer specifying the number of Gap and ACK Range fields in the frame.

First ACK Range: A variable-length integer indicating the number of contiguous packets preceding the Largest Acknowledged that are being acknowledged. The First ACK Range is encoded as an ACK Range; see [Section 19.3.1](#) starting from the Largest Acknowledged. That is, the smallest packet acknowledged in the range is determined by subtracting the First ACK Range value from the Largest Acknowledged.

ACK Ranges: Contains additional ranges of packets which are alternately not acknowledged (Gap) and acknowledged (ACK Range); see [Section 19.3.1](#).

ECN Counts: The three ECN Counts; see [Section 19.3.2](#).

19.3.1. ACK Ranges

Each ACK Range consists of alternating Gap and ACK Range values in descending packet number order. ACK Ranges can be repeated. The number of Gap and ACK Range values is determined by the ACK Range Count field; one of each value is present for each value in the ACK Range Count field.

ACK Ranges are structured as shown in Figure 25.

```
ACK Range {  
    Gap (i),  
    ACK Range Length (i),  
}
```

Figure 25: ACK Ranges

The fields that form each ACK Range are:

Gap: A variable-length integer indicating the number of contiguous unacknowledged packets preceding the packet number one lower than the smallest in the preceding ACK Range.

ACK Range Length: A variable-length integer indicating the number of contiguous acknowledged packets preceding the largest packet number, as determined by the preceding Gap.

Gap and ACK Range value use a relative integer encoding for efficiency. Though each encoded value is positive, the values are subtracted, so that each ACK Range describes progressively lower-numbered packets.

Each ACK Range acknowledges a contiguous range of packets by indicating the number of acknowledged packets that precede the largest packet number in that range. A value of zero indicates that only the largest packet number is acknowledged. Larger ACK Range values indicate a larger range, with corresponding lower values for the smallest packet number in the range. Thus, given a largest packet number for the range, the smallest value is determined by the formula:

$$\text{smallest} = \text{largest} - \text{ack_range}$$

An ACK Range acknowledges all packets between the smallest packet number and the largest, inclusive.

The largest value for an ACK Range is determined by cumulatively subtracting the size of all preceding ACK Ranges and Gaps.

Each Gap indicates a range of packets that are not being acknowledged. The number of packets in the gap is one higher than the encoded value of the Gap field.

The value of the Gap field establishes the largest packet number value for the subsequent ACK Range using the following formula:

$$\text{largest} = \text{previous_smallest} - \text{gap} - 2$$

If any computed packet number is negative, an endpoint MUST generate a connection error of type `FRAME_ENCODING_ERROR`.

19.3.2. ECN Counts

The ACK frame uses the least significant bit (that is, type 0x03) to indicate ECN feedback and report receipt of QUIC packets with associated ECN codepoints of ECT(0), ECT(1), or CE in the packet's IP header. ECN Counts are only present when the ACK frame type is 0x03.

ECN Counts are only parsed when the ACK frame type is 0x03. There are 3 ECN counts, as shown in Figure 26.

```
ECN Counts {
    ECT0 Count (i),
    ECT1 Count (i),
    ECN-CE Count (i),
}
```

Figure 26: ECN Count Format

The three ECN Counts are:

ECT0 Count: A variable-length integer representing the total number of packets received with the ECT(0) codepoint in the packet number space of the ACK frame.

ECT1 Count: A variable-length integer representing the total number of packets received with the ECT(1) codepoint in the packet number space of the ACK frame.

CE Count: A variable-length integer representing the total number of packets received with the CE codepoint in the packet number space of the ACK frame.

ECN counts are maintained separately for each packet number space.

19.4. RESET_STREAM Frame

An endpoint uses a RESET_STREAM frame (type=0x04) to abruptly terminate the sending part of a stream.

After sending a RESET_STREAM, an endpoint ceases transmission and retransmission of STREAM frames on the identified stream. A receiver of RESET_STREAM can discard any data that it already received on that stream.

An endpoint that receives a RESET_STREAM frame for a send-only stream MUST terminate the connection with error STREAM_STATE_ERROR.

The RESET_STREAM frame is shown in Figure 27.

```
RESET_STREAM Frame {
  Type (i) = 0x04,
  Stream ID (i),
  Application Protocol Error Code (i),
  Final Size (i),
}
```

Figure 27: RESET_STREAM Frame Format

RESET_STREAM frames contain the following fields:

Stream ID: A variable-length integer encoding of the Stream ID of the stream being terminated.

Application Protocol Error Code: A variable-length integer containing the application protocol error code (see [Section 20.1](#)) which indicates why the stream is being closed.

Final Size: A variable-length integer indicating the final size of

the stream by the RESET_STREAM sender, in unit of bytes.

19.5. STOP_SENDING Frame

An endpoint uses a STOP_SENDING frame (type=0x05) to communicate that incoming data is being discarded on receipt at application request. STOP_SENDING requests that a peer cease transmission on a stream.

A STOP_SENDING frame can be sent for streams in the Recv or Size Known states; see [Section 3.1](#). Receiving a STOP_SENDING frame for a locally-initiated stream that has not yet been created MUST be treated as a connection error of type STREAM_STATE_ERROR. An endpoint that receives a STOP_SENDING frame for a receive-only stream MUST terminate the connection with error STREAM_STATE_ERROR.

The STOP_SENDING frame is shown in Figure 28.

```
STOP_SENDING Frame {
  Type (i) = 0x05,
  Stream ID (i),
  Application Protocol Error Code (i),
}
```

Figure 28: STOP_SENDING Frame Format

STOP_SENDING frames contain the following fields:

Stream ID: A variable-length integer carrying the Stream ID of the stream being ignored.

Application Protocol Error Code: A variable-length integer containing the application-specified reason the sender is ignoring the stream; see [Section 20.1](#).

19.6. CRYPTO Frame

The CRYPTO frame (type=0x06) is used to transmit cryptographic handshake messages. It can be sent in all packet types except 0-RTT. The CRYPTO frame offers the cryptographic protocol an in-order stream of bytes. CRYPTO frames are functionally identical to STREAM frames, except that they do not bear a stream identifier; they are not flow controlled; and they do not carry markers for optional offset, optional length, and the end of the stream.

The CRYPTO frame is shown in Figure 29.

```
CRYPTO Frame {
  Type (i) = 0x06,
  Offset (i),
  Length (i),
  Crypto Data (...),
}
```

Figure 29: CRYPTO Frame Format

CRYPTO frames contain the following fields:

Offset: A variable-length integer specifying the byte offset in the stream for the data in this CRYPTO frame.

Length: A variable-length integer specifying the length of the Crypto Data field in this CRYPTO frame.

Crypto Data: The cryptographic message data.

There is a separate flow of cryptographic handshake data in each encryption level, each of which starts at an offset of 0. This implies that each encryption level is treated as a separate CRYPTO stream of data.

The largest offset delivered on a stream - the sum of the offset and data length - cannot exceed $2^{62}-1$. Receipt of a frame that exceeds this limit MUST be treated as a connection error of type `FRAME_ENCODING_ERROR` or `CRYPTO_BUFFER_EXCEEDED`.

Unlike STREAM frames, which include a Stream ID indicating to which stream the data belongs, the CRYPTO frame carries data for a single stream per encryption level. The stream does not have an explicit end, so CRYPTO frames do not have a FIN bit.

19.7. NEW_TOKEN Frame

A server sends a NEW_TOKEN frame (type=0x07) to provide the client with a token to send in the header of an Initial packet for a future connection.

The NEW_TOKEN frame is shown in Figure 30.

```
NEW_TOKEN Frame {
  Type (i) = 0x07,
  Token Length (i),
  Token (...),
}
```


Figure 30: NEW_TOKEN Frame Format

NEW_TOKEN frames contain the following fields:

Token Length: A variable-length integer specifying the length of the token in bytes.

Token: An opaque blob that the client may use with a future Initial packet. The token **MUST NOT** be empty. An endpoint **MUST** treat receipt of a NEW_TOKEN frame with an empty Token field as a connection error of type FRAME_ENCODING_ERROR.

An endpoint might receive multiple NEW_TOKEN frames that contain the same token value if packets containing the frame are incorrectly determined to be lost. Endpoints are responsible for discarding duplicate values, which might be used to link connection attempts; see [Section 8.1.3](#).

Clients **MUST NOT** send NEW_TOKEN frames. Servers **MUST** treat receipt of a NEW_TOKEN frame as a connection error of type PROTOCOL_VIOLATION.

19.8. STREAM Frames

STREAM frames implicitly create a stream and carry stream data. The STREAM frame takes the form 0b00001XXX (or the set of values from 0x08 to 0x0f). The value of the three low-order bits of the frame type determines the fields that are present in the frame.

- * The OFF bit (0x04) in the frame type is set to indicate that there is an Offset field present. When set to 1, the Offset field is present. When set to 0, the Offset field is absent and the Stream Data starts at an offset of 0 (that is, the frame contains the first bytes of the stream, or the end of a stream that includes no data).
- * The LEN bit (0x02) in the frame type is set to indicate that there is a Length field present. If this bit is set to 0, the Length field is absent and the Stream Data field extends to the end of the packet. If this bit is set to 1, the Length field is present.
- * The FIN bit (0x01) of the frame type is set only on frames that contain the final size of the stream. Setting this bit indicates that the frame marks the end of the stream.

An endpoint MUST terminate the connection with error `STREAM_STATE_ERROR` if it receives a `STREAM` frame for a locally-initiated stream that has not yet been created, or for a send-only stream.

The `STREAM` frames are shown in Figure 31.

```
STREAM Frame {
  Type (i) = 0x08..0x0f,
  Stream ID (i),
  [Offset (i)],
  [Length (i)],
  Stream Data (...),
}
```

Figure 31: `STREAM` Frame Format

`STREAM` frames contain the following fields:

Stream ID: A variable-length integer indicating the stream ID of the stream; see [Section 2.1](#).

Offset: A variable-length integer specifying the byte offset in the stream for the data in this `STREAM` frame. This field is present when the `OFF` bit is set to 1. When the `Offset` field is absent, the offset is 0.

Length: A variable-length integer specifying the length of the `Stream Data` field in this `STREAM` frame. This field is present when the `LEN` bit is set to 1. When the `LEN` bit is set to 0, the `Stream Data` field consumes all the remaining bytes in the packet.

Stream Data: The bytes from the designated stream to be delivered.

When a `Stream Data` field has a length of 0, the offset in the `STREAM` frame is the offset of the next byte that would be sent.

The first byte in the stream has an offset of 0. The largest offset delivered on a stream - the sum of the offset and data length - cannot exceed $2^{62}-1$, as it is not possible to provide flow control credit for that data. Receipt of a frame that exceeds this limit MUST be treated as a connection error of type `FRAME_ENCODING_ERROR` or `FLOW_CONTROL_ERROR`.

19.9. MAX_DATA Frame

The MAX_DATA frame (type=0x10) is used in flow control to inform the peer of the maximum amount of data that can be sent on the connection as a whole.

The MAX_DATA frame is shown in Figure 32.

```
MAX_DATA Frame {
  Type (i) = 0x10,
  Maximum Data (i),
}
```

Figure 32: MAX_DATA Frame Format

MAX_DATA frames contain the following fields:

Maximum Data: A variable-length integer indicating the maximum amount of data that can be sent on the entire connection, in units of bytes.

All data sent in STREAM frames counts toward this limit. The sum of the largest received offsets on all streams - including streams in terminal states - MUST NOT exceed the value advertised by a receiver. An endpoint MUST terminate a connection with a FLOW_CONTROL_ERROR error if it receives more data than the maximum data value that it has sent, unless this is a result of a change in the initial limits; see [Section 7.4.1](#).

19.10. MAX_STREAM_DATA Frame

The MAX_STREAM_DATA frame (type=0x11) is used in flow control to inform a peer of the maximum amount of data that can be sent on a stream.

A MAX_STREAM_DATA frame can be sent for streams in the Recv state; see [Section 3.1](#). Receiving a MAX_STREAM_DATA frame for a locally-initiated stream that has not yet been created MUST be treated as a connection error of type STREAM_STATE_ERROR. An endpoint that receives a MAX_STREAM_DATA frame for a receive-only stream MUST terminate the connection with error STREAM_STATE_ERROR.

The MAX_STREAM_DATA frame is shown in Figure 33.

```
MAX_STREAM_DATA Frame {
  Type (i) = 0x11,
  Stream ID (i),
  Maximum Stream Data (i),
}
```

Figure 33: MAX_STREAM_DATA Frame Format

MAX_STREAM_DATA frames contain the following fields:

Stream ID: The stream ID of the stream that is affected encoded as a variable-length integer.

Maximum Stream Data: A variable-length integer indicating the maximum amount of data that can be sent on the identified stream, in units of bytes.

When counting data toward this limit, an endpoint accounts for the largest received offset of data that is sent or received on the stream. Loss or reordering can mean that the largest received offset on a stream can be greater than the total size of data received on that stream. Receiving STREAM frames might not increase the largest received offset.

The data sent on a stream MUST NOT exceed the largest maximum stream data value advertised by the receiver. An endpoint MUST terminate a connection with a FLOW_CONTROL_ERROR error if it receives more data than the largest maximum stream data that it has sent for the affected stream, unless this is a result of a change in the initial limits; see [Section 7.4.1](#).

19.11. MAX_STREAMS Frames

The MAX_STREAMS frames (type=0x12 and 0x13) inform the peer of the cumulative number of streams of a given type it is permitted to open. A MAX_STREAMS frame with a type of 0x12 applies to bidirectional streams, and a MAX_STREAMS frame with a type of 0x13 applies to unidirectional streams.

The MAX_STREAMS frames are shown in Figure 34;

```
MAX_STREAMS Frame {
  Type (i) = 0x12..0x13,
  Maximum Streams (i),
}
```

Figure 34: MAX_STREAMS Frame Format

MAX_STREAMS frames contain the following fields:

Maximum Streams: A count of the cumulative number of streams of the corresponding type that can be opened over the lifetime of the connection. This value cannot exceed 2^{60} , as it is not possible to encode stream IDs larger than $2^{62}-1$. Receipt of a frame that permits opening of a stream larger than this limit MUST be treated as a FRAME_ENCODING_ERROR.

Loss or reordering can cause a MAX_STREAMS frame to be received which states a lower stream limit than an endpoint has previously received. MAX_STREAMS frames which do not increase the stream limit MUST be ignored.

An endpoint MUST NOT open more streams than permitted by the current stream limit set by its peer. For instance, a server that receives a unidirectional stream limit of 3 is permitted to open stream 3, 7, and 11, but not stream 15. An endpoint MUST terminate a connection with a STREAM_LIMIT_ERROR error if a peer opens more streams than was permitted.

Note that these frames (and the corresponding transport parameters) do not describe the number of streams that can be opened concurrently. The limit includes streams that have been closed as well as those that are open.

19.12. DATA_BLOCKED Frame

A sender SHOULD send a DATA_BLOCKED frame (type=0x14) when it wishes to send data, but is unable to due to connection-level flow control; see [Section 4](#). DATA_BLOCKED frames can be used as input to tuning of flow control algorithms; see [Section 4.2](#).

The DATA_BLOCKED frame is shown in Figure 35.

```
DATA_BLOCKED Frame {
  Type (i) = 0x14,
  Maximum Data (i),
}
```

Figure 35: DATA_BLOCKED Frame Format

DATA_BLOCKED frames contain the following fields:

Maximum Data: A variable-length integer indicating the connection-level limit at which blocking occurred.

19.13. STREAM_DATA_BLOCKED Frame

A sender SHOULD send a STREAM_DATA_BLOCKED frame (type=0x15) when it wishes to send data, but is unable to due to stream-level flow control. This frame is analogous to DATA_BLOCKED (Section 19.12).

An endpoint that receives a STREAM_DATA_BLOCKED frame for a send-only stream MUST terminate the connection with error STREAM_STATE_ERROR.

The STREAM_DATA_BLOCKED frame is shown in Figure 36.

```
STREAM_DATA_BLOCKED Frame {  
  Type (i) = 0x15,  
  Stream ID (i),  
  Maximum Stream Data (i),  
}
```

Figure 36: STREAM_DATA_BLOCKED Frame Format

STREAM_DATA_BLOCKED frames contain the following fields:

Stream ID: A variable-length integer indicating the stream which is flow control blocked.

Maximum Stream Data: A variable-length integer indicating the offset of the stream at which the blocking occurred.

19.14. STREAMS_BLOCKED Frames

A sender SHOULD send a STREAMS_BLOCKED frame (type=0x16 or 0x17) when it wishes to open a stream, but is unable to due to the maximum stream limit set by its peer; see Section 19.11. A STREAMS_BLOCKED frame of type 0x16 is used to indicate reaching the bidirectional stream limit, and a STREAMS_BLOCKED frame of type 0x17 indicates reaching the unidirectional stream limit.

A STREAMS_BLOCKED frame does not open the stream, but informs the peer that a new stream was needed and the stream limit prevented the creation of the stream.

The STREAMS_BLOCKED frames are shown in Figure 37.

```
STREAMS_BLOCKED Frame {  
  Type (i) = 0x16..0x17,  
  Maximum Streams (i),  
}
```

Figure 37: STREAMS_BLOCKED Frame Format

STREAMS_BLOCKED frames contain the following fields:

Maximum Streams: A variable-length integer indicating the maximum number of streams allowed at the time the frame was sent. This value cannot exceed 2^{60} , as it is not possible to encode stream IDs larger than $2^{62}-1$. Receipt of a frame that encodes a larger stream ID MUST be treated as a STREAM_LIMIT_ERROR or a FRAME_ENCODING_ERROR.

19.15. NEW_CONNECTION_ID Frame

An endpoint sends a NEW_CONNECTION_ID frame (type=0x18) to provide its peer with alternative connection IDs that can be used to break linkability when migrating connections; see [Section 9.5](#).

The NEW_CONNECTION_ID frame is shown in Figure 38.

```
NEW_CONNECTION_ID Frame {
  Type (i) = 0x18,
  Sequence Number (i),
  Retire Prior To (i),
  Length (8),
  Connection ID (8..160),
  Stateless Reset Token (128),
}
```

Figure 38: NEW_CONNECTION_ID Frame Format

NEW_CONNECTION_ID frames contain the following fields:

Sequence Number: The sequence number assigned to the connection ID by the sender. See [Section 5.1.1](#).

Retire Prior To: A variable-length integer indicating which connection IDs should be retired; see [Section 5.1.2](#).

Length: An 8-bit unsigned integer containing the length of the connection ID. Values less than 1 and greater than 20 are invalid and MUST be treated as a connection error of type FRAME_ENCODING_ERROR.

Connection ID: A connection ID of the specified length.

Stateless Reset Token: A 128-bit value that will be used for a stateless reset when the associated connection ID is used; see [Section 10.4](#).

An endpoint MUST NOT send this frame if it currently requires that its peer send packets with a zero-length Destination Connection ID. Changing the length of a connection ID to or from zero-length makes it difficult to identify when the value of the connection ID changed. An endpoint that is sending packets with a zero-length Destination Connection ID MUST treat receipt of a NEW_CONNECTION_ID frame as a connection error of type `PROTOCOL_VIOLATION`.

Transmission errors, timeouts and retransmissions might cause the same NEW_CONNECTION_ID frame to be received multiple times. Receipt of the same frame multiple times MUST NOT be treated as a connection error. A receiver can use the sequence number supplied in the NEW_CONNECTION_ID frame to identify new connection IDs from old ones.

If an endpoint receives a NEW_CONNECTION_ID frame that repeats a previously issued connection ID with a different Stateless Reset Token or a different sequence number, or if a sequence number is used for different connection IDs, the endpoint MAY treat that receipt as a connection error of type `PROTOCOL_VIOLATION`.

The `Retire Prior To` field counts connection IDs established during connection setup and the `preferred_address` transport parameter; see [Section 5.1.2](#). The `Retire Prior To` field MUST be less than or equal to the `Sequence Number` field. Receiving a value greater than the `Sequence Number` MUST be treated as a connection error of type `FRAME_ENCODING_ERROR`.

Once a sender indicates a `Retire Prior To` value, smaller values sent in subsequent NEW_CONNECTION_ID frames have no effect. A receiver MUST ignore any `Retire Prior To` fields that do not increase the largest received `Retire Prior To` value.

An endpoint that receives a NEW_CONNECTION_ID frame with a sequence number smaller than the `Retire Prior To` field of a previously received NEW_CONNECTION_ID frame MUST send a corresponding RETIRE_CONNECTION_ID frame that retires the newly received connection ID, unless it has already done so for that sequence number.

19.16. RETIRE_CONNECTION_ID Frame

An endpoint sends a RETIRE_CONNECTION_ID frame (type=0x19) to indicate that it will no longer use a connection ID that was issued by its peer. This may include the connection ID provided during the handshake. Sending a RETIRE_CONNECTION_ID frame also serves as a request to the peer to send additional connection IDs for future use; see [Section 5.1](#). New connection IDs can be delivered to a peer using the NEW_CONNECTION_ID frame ([Section 19.15](#)).

Retiring a connection ID invalidates the stateless reset token associated with that connection ID.

The RETIRE_CONNECTION_ID frame is shown in Figure 39.

```
RETIRE_CONNECTION_ID Frame {
  Type (i) = 0x19,
  Sequence Number (i),
}
```

Figure 39: RETIRE_CONNECTION_ID Frame Format

RETIRE_CONNECTION_ID frames contain the following fields:

Sequence Number: The sequence number of the connection ID being retired; see [Section 5.1.2](#).

Receipt of a RETIRE_CONNECTION_ID frame containing a sequence number greater than any previously sent to the peer MUST be treated as a connection error of type `PROTOCOL_VIOLATION`.

The sequence number specified in a RETIRE_CONNECTION_ID frame MUST NOT refer to the Destination Connection ID field of the packet in which the frame is contained. The peer MAY treat this as a connection error of type `FRAME_ENCODING_ERROR`.

An endpoint cannot send this frame if it was provided with a zero-length connection ID by its peer. An endpoint that provides a zero-length connection ID MUST treat receipt of a RETIRE_CONNECTION_ID frame as a connection error of type `PROTOCOL_VIOLATION`.

19.17. PATH_CHALLENGE Frame

Endpoints can use `PATH_CHALLENGE` frames (type=0x1a) to check reachability to the peer and for path validation during connection migration.

The `PATH_CHALLENGE` frame is shown in Figure 40.

```
PATH_CHALLENGE Frame {
  Type (i) = 0x1a,
  Data (64),
}
```

Figure 40: PATH_CHALLENGE Frame Format

`PATH_CHALLENGE` frames contain the following fields:

Data: This 8-byte field contains arbitrary data.

A `PATH_CHALLENGE` frame containing 8 bytes that are hard to guess is sufficient to ensure that it is easier to receive the packet than it is to guess the value correctly.

The recipient of this frame **MUST** generate a `PATH_RESPONSE` frame ([Section 19.18](#)) containing the same Data.

19.18. `PATH_RESPONSE` Frame

The `PATH_RESPONSE` frame (type=0x1b) is sent in response to a `PATH_CHALLENGE` frame. Its format, shown in Figure 41 is identical to the `PATH_CHALLENGE` frame ([Section 19.17](#)).

```
PATH_RESPONSE Frame {
  Type (i) = 0x1b,
  Data (64),
}
```

Figure 41: `PATH_RESPONSE` Frame Format

If the content of a `PATH_RESPONSE` frame does not match the content of a `PATH_CHALLENGE` frame previously sent by the endpoint, the endpoint **MAY** generate a connection error of type `PROTOCOL_VIOLATION`.

19.19. `CONNECTION_CLOSE` Frames

An endpoint sends a `CONNECTION_CLOSE` frame (type=0x1c or 0x1d) to notify its peer that the connection is being closed. The `CONNECTION_CLOSE` with a frame type of 0x1c is used to signal errors at only the QUIC layer, or the absence of errors (with the `NO_ERROR` code). The `CONNECTION_CLOSE` frame with a type of 0x1d is used to signal an error with the application that uses QUIC.

If there are open streams that haven't been explicitly closed, they are implicitly closed when the connection is closed.

The `CONNECTION_CLOSE` frames are shown in Figure 42.

```
CONNECTION_CLOSE Frame {
  Type (i) = 0x1c..0x1d,
  Error Code (i),
  [Frame Type (i)],
  Reason Phrase Length (i),
  Reason Phrase (...),
}
```

Figure 42: CONNECTION_CLOSE Frame Format

CONNECTION_CLOSE frames contain the following fields:

Error Code: A variable length integer error code which indicates the reason for closing this connection. A CONNECTION_CLOSE frame of type 0x1c uses codes from the space defined in [Section 20](#). A CONNECTION_CLOSE frame of type 0x1d uses codes from the application protocol error code space; see [Section 20.1](#).

Frame Type: A variable-length integer encoding the type of frame that triggered the error. A value of 0 (equivalent to the mention of the PADDING frame) is used when the frame type is unknown. The application-specific variant of CONNECTION_CLOSE (type 0x1d) does not include this field.

Reason Phrase Length: A variable-length integer specifying the length of the reason phrase in bytes. Because a CONNECTION_CLOSE frame cannot be split between packets, any limits on packet size will also limit the space available for a reason phrase.

Reason Phrase: A human-readable explanation for why the connection was closed. This can be zero length if the sender chooses to not give details beyond the Error Code. This SHOULD be a UTF-8 encoded string [[RFC3629](#)].

The application-specific variant of CONNECTION_CLOSE (type 0x1d) can only be sent using 0-RTT or 1-RTT packets; see [Section 4](#) of [[QUIC-TLS](#)]. When an application wishes to abandon a connection during the handshake, an endpoint can send a CONNECTION_CLOSE frame (type 0x1c) with an error code of APPLICATION_ERROR in an Initial or a Handshake packet.

19.20. HANDSHAKE_DONE frame

The server uses the HANDSHAKE_DONE frame (type=0x1e) to signal confirmation of the handshake to the client. As shown in [Figure 43](#), a HANDSHAKE_DONE frame has no content.

```
HANDSHAKE_DONE Frame {  
    Type (i) = 0x1e,  
}
```

Figure 43: HANDSHAKE_DONE Frame Format

A HANDSHAKE_DONE frame can only be sent by the server. Servers MUST NOT send a HANDSHAKE_DONE frame before completing the handshake. A server MUST treat receipt of a HANDSHAKE_DONE frame as a connection error of type PROTOCOL_VIOLATION.

19.21. Extension Frames

QUIC frames do not use a self-describing encoding. An endpoint therefore needs to understand the syntax of all frames before it can successfully process a packet. This allows for efficient encoding of frames, but it means that an endpoint cannot send a frame of a type that is unknown to its peer.

An extension to QUIC that wishes to use a new type of frame MUST first ensure that a peer is able to understand the frame. An endpoint can use a transport parameter to signal its willingness to receive one or more extension frame types with the one transport parameter.

Extensions that modify or replace core protocol functionality (including frame types) will be difficult to combine with other extensions which modify or replace the same functionality unless the behavior of the combination is explicitly defined. Such extensions SHOULD define their interaction with previously-defined extensions modifying the same protocol components.

Extension frames MUST be congestion controlled and MUST cause an ACK frame to be sent. The exception is extension frames that replace or supplement the ACK frame. Extension frames are not included in flow control unless specified in the extension.

An IANA registry is used to manage the assignment of frame types; see [Section 22.3](#).

20. Transport Error Codes

QUIC error codes are 62-bit unsigned integers.

This section lists the defined QUIC transport error codes that may be used in a CONNECTION_CLOSE frame. These errors apply to the entire connection.

NO_ERROR (0x0): An endpoint uses this with CONNECTION_CLOSE to signal that the connection is being closed abruptly in the absence of any error.

INTERNAL_ERROR (0x1): The endpoint encountered an internal error and cannot continue with the connection.

- CONNECTION_REFUSED (0x2): The server refused to accept a new connection.
- FLOW_CONTROL_ERROR (0x3): An endpoint received more data than it permitted in its advertised data limits; see [Section 4](#).
- STREAM_LIMIT_ERROR (0x4): An endpoint received a frame for a stream identifier that exceeded its advertised stream limit for the corresponding stream type.
- STREAM_STATE_ERROR (0x5): An endpoint received a frame for a stream that was not in a state that permitted that frame; see [Section 3](#).
- FINAL_SIZE_ERROR (0x6): An endpoint received a STREAM frame containing data that exceeded the previously established final size. Or an endpoint received a STREAM frame or a RESET_STREAM frame containing a final size that was lower than the size of stream data that was already received. Or an endpoint received a STREAM frame or a RESET_STREAM frame containing a different final size to the one already established.
- FRAME_ENCODING_ERROR (0x7): An endpoint received a frame that was badly formatted. For instance, a frame of an unknown type, or an ACK frame that has more acknowledgment ranges than the remainder of the packet could carry.
- TRANSPORT_PARAMETER_ERROR (0x8): An endpoint received transport parameters that were badly formatted, included an invalid value, was absent even though it is mandatory, was present though it is forbidden, or is otherwise in error.
- CONNECTION_ID_LIMIT_ERROR (0x9): The number of connection IDs provided by the peer exceeds the advertised `active_connection_id_limit`.
- PROTOCOL_VIOLATION (0xA): An endpoint detected an error with protocol compliance that was not covered by more specific error codes.
- INVALID_TOKEN (0xB): A server received a Retry Token in a client Initial that is invalid.
- APPLICATION_ERROR (0xC): The application or application protocol caused the connection to be closed.
- CRYPTO_BUFFER_EXCEEDED (0xD): An endpoint has received more data in CRYPTO frames than it can buffer.

CRYPTO_ERROR (0x1XX): The cryptographic handshake failed. A range of 256 values is reserved for carrying error codes specific to the cryptographic handshake that is used. Codes for errors occurring when TLS is used for the crypto handshake are described in Section 4.8 of [QUIC-TLS].

See [Section 22.4](#) for details of registering new error codes.

In defining these error codes, several principles are applied. Error conditions that might require specific action on the part of a recipient are given unique codes. Errors that represent common conditions are given specific codes. Absent either of these conditions, error codes are used to identify a general function of the stack, like flow control or transport parameter handling. Finally, generic errors are provided for conditions where implementations are unable or unwilling to use more specific codes.

20.1. Application Protocol Error Codes

Application protocol error codes are 62-bit unsigned integers, but the management of application error codes is left to application protocols. Application protocol error codes are used for the RESET_STREAM frame ([Section 19.4](#)), the STOP_SENDING frame ([Section 19.5](#)), and the CONNECTION_CLOSE frame with a type of 0x1d ([Section 19.19](#)).

21. Security Considerations

21.1. Handshake Denial of Service

As an encrypted and authenticated transport QUIC provides a range of protections against denial of service. Once the cryptographic handshake is complete, QUIC endpoints discard most packets that are not authenticated, greatly limiting the ability of an attacker to interfere with existing connections.

Once a connection is established QUIC endpoints might accept some unauthenticated ICMP packets (see [Section 14.2.1](#)), but the use of these packets is extremely limited. The only other type of packet that an endpoint might accept is a stateless reset ([Section 10.4](#)) which relies on the token being kept secret until it is used.

During the creation of a connection, QUIC only provides protection against attack from off the network path. All QUIC packets contain proof that the recipient saw a preceding packet from its peer.

Addresses cannot change during the handshake, so endpoints can discard packets that are received on a different network path.

The Source and Destination Connection ID fields are the primary means of protection against off-path attack during the handshake. These are required to match those set by a peer. Except for an Initial and stateless reset packets, an endpoint only accepts packets that include a Destination Connection ID field that matches a value the endpoint previously chose. This is the only protection offered for Version Negotiation packets.

The Destination Connection ID field in an Initial packet is selected by a client to be unpredictable, which serves an additional purpose. The packets that carry the cryptographic handshake are protected with a key that is derived from this connection ID and salt specific to the QUIC version. This allows endpoints to use the same process for authenticating packets that they receive as they use after the cryptographic handshake completes. Packets that cannot be authenticated are discarded. Protecting packets in this fashion provides a strong assurance that the sender of the packet saw the Initial packet and understood it.

These protections are not intended to be effective against an attacker that is able to receive QUIC packets prior to the connection being established. Such an attacker can potentially send packets that will be accepted by QUIC endpoints. This version of QUIC attempts to detect this sort of attack, but it expects that endpoints will fail to establish a connection rather than recovering. For the most part, the cryptographic handshake protocol [QUIC-TLS] is responsible for detecting tampering during the handshake.

Endpoints are permitted to use other methods to detect and attempt to recover from interference with the handshake. Invalid packets may be identified and discarded using other methods, but no specific method is mandated in this document.

21.2. Amplification Attack

An attacker might be able to receive an address validation token (Section 8) from a server and then release the IP address it used to acquire that token. At a later time, the attacker may initiate a 0-RTT connection with a server by spoofing this same address, which might now address a different (victim) endpoint. The attacker can thus potentially cause the server to send an initial congestion window's worth of data towards the victim.

Servers SHOULD provide mitigations for this attack by limiting the usage and lifetime of address validation tokens; see Section 8.1.3.

21.3. Optimistic ACK Attack

An endpoint that acknowledges packets it has not received might cause a congestion controller to permit sending at rates beyond what the network supports. An endpoint MAY skip packet numbers when sending packets to detect this behavior. An endpoint can then immediately close the connection with a connection error of type `PROTOCOL_VIOLATION`; see [Section 10.3](#).

21.4. Slowloris Attacks

The attacks commonly known as Slowloris [[SLOWLORIS](#)] try to keep many connections to the target endpoint open and hold them open as long as possible. These attacks can be executed against a QUIC endpoint by generating the minimum amount of activity necessary to avoid being closed for inactivity. This might involve sending small amounts of data, gradually opening flow control windows in order to control the sender rate, or manufacturing ACK frames that simulate a high loss rate.

QUIC deployments SHOULD provide mitigations for the Slowloris attacks, such as increasing the maximum number of clients the server will allow, limiting the number of connections a single IP address is allowed to make, imposing restrictions on the minimum transfer speed a connection is allowed to have, and restricting the length of time an endpoint is allowed to stay connected.

21.5. Stream Fragmentation and Reassembly Attacks

An adversarial sender might intentionally send fragments of stream data in order to cause disproportionate receive buffer memory commitment and/or creation of a large and inefficient data structure.

An adversarial receiver might intentionally not acknowledge packets containing stream data in order to force the sender to store the unacknowledged stream data for retransmission.

The attack on receivers is mitigated if flow control windows correspond to available memory. However, some receivers will over-commit memory and advertise flow control offsets in the aggregate that exceed actual available memory. The over-commitment strategy can lead to better performance when endpoints are well behaved, but renders endpoints vulnerable to the stream fragmentation attack.

QUIC deployments SHOULD provide mitigations against stream fragmentation attacks. Mitigations could consist of avoiding over-committing memory, limiting the size of tracking data structures, delaying reassembly of STREAM frames, implementing heuristics based on the age and duration of reassembly holes, or some combination.

21.6. Stream Commitment Attack

An adversarial endpoint can open lots of streams, exhausting state on an endpoint. The adversarial endpoint could repeat the process on a large number of connections, in a manner similar to SYN flooding attacks in TCP.

Normally, clients will open streams sequentially, as explained in [Section 2.1](#). However, when several streams are initiated at short intervals, loss or reordering may cause STREAM frames that open streams to be received out of sequence. On receiving a higher-numbered stream ID, a receiver is required to open all intervening streams of the same type; see [Section 3.2](#). Thus, on a new connection, opening stream 4000000 opens 1 million and 1 client-initiated bidirectional streams.

The number of active streams is limited by the `initial_max_streams_bidi` and `initial_max_streams_uni` transport parameters, as explained in [Section 4.5](#). If chosen judiciously, these limits mitigate the effect of the stream commitment attack. However, setting the limit too low could affect performance when applications expect to open large number of streams.

21.7. Peer Denial of Service

QUIC and TLS both contain frames or messages that have legitimate uses in some contexts, but that can be abused to cause a peer to expend processing resources without having any observable impact on the state of the connection.

Messages can also be used to change and revert state in small or inconsequential ways, such as by sending small increments to flow control limits.

If processing costs are disproportionately large in comparison to bandwidth consumption or effect on state, then this could allow a malicious peer to exhaust processing capacity.

While there are legitimate uses for all messages, implementations SHOULD track cost of processing relative to progress and treat excessive quantities of any non-productive packets as indicative of an attack. Endpoints MAY respond to this condition with a connection error, or by dropping packets.

21.8. Explicit Congestion Notification Attacks

An on-path attacker could manipulate the value of ECN codepoints in the IP header to influence the sender's rate. [RFC3168] discusses manipulations and their effects in more detail.

An on-the-side attacker can duplicate and send packets with modified ECN codepoints to affect the sender's rate. If duplicate packets are discarded by a receiver, an off-path attacker will need to race the duplicate packet against the original to be successful in this attack. Therefore, QUIC endpoints ignore the ECN codepoint field on an IP packet unless at least one QUIC packet in that IP packet is successfully processed; see [Section 13.4](#).

21.9. Stateless Reset Oracle

Stateless resets create a possible denial of service attack analogous to a TCP reset injection. This attack is possible if an attacker is able to cause a stateless reset token to be generated for a connection with a selected connection ID. An attacker that can cause this token to be generated can reset an active connection with the same connection ID.

If a packet can be routed to different instances that share a static key, for example by changing an IP address or port, then an attacker can cause the server to send a stateless reset. To defend against this style of denial service, endpoints that share a static key for stateless reset (see [Section 10.4.2](#)) MUST be arranged so that packets with a given connection ID always arrive at an instance that has connection state, unless that connection is no longer active.

In the case of a cluster that uses dynamic load balancing, it's possible that a change in load balancer configuration could happen while an active instance retains connection state; even if an instance retains connection state, the change in routing and resulting stateless reset will result in the connection being terminated. If there is no chance in the packet being routed to the correct instance, it is better to send a stateless reset than wait for connections to time out. However, this is acceptable only if the routing cannot be influenced by an attacker.

21.10. Version Downgrade

This document defines QUIC Version Negotiation packets in [Section 6](#), which can be used to negotiate the QUIC version used between two endpoints. However, this document does not specify how this negotiation will be performed between this version and subsequent future versions. In particular, Version Negotiation packets do not contain any mechanism to prevent version downgrade attacks. Future versions of QUIC that use Version Negotiation packets **MUST** define a mechanism that is robust against version downgrade attacks.

21.11. Targeted Attacks by Routing

Deployments should limit the ability of an attacker to target a new connection to a particular server instance. This means that client-controlled fields, such as the initial Destination Connection ID used on Initial and 0-RTT packets **SHOULD NOT** be used by themselves to make routing decisions. Ideally, routing decisions are made independently of client-selected values; a Source Connection ID can be selected to route later packets to the same server.

21.12. Overview of Security Properties

A complete security analysis of QUIC is outside the scope of this document. This section provides an informal description of the desired security properties as an aid to implementors and to help guide protocol analysis.

QUIC assumes the threat model described in [\[SEC-CONS\]](#) and provides protections against many of the attacks that arise from that model.

For this purpose, attacks are divided into passive and active attacks. Passive attackers have the capability to read packets from the network, while active attackers also have the capability to write packets into the network. However, a passive attack may involve an attacker with the ability to cause a routing change or other modification in the path taken by packets that comprise a connection.

Attackers are additionally categorized as either on-path attackers or off-path attackers; see Section 3.5 of [\[SEC-CONS\]](#). An on-path attacker can read, modify, or remove any packet it observes such that it no longer reaches its destination, while an off-path attacker observes the packets, but cannot prevent the original packet from reaching its intended destination. An off-path attacker can also transmit arbitrary packets.

Properties of the handshake, protected packets, and connection migration are considered separately.

21.12.1. Handshake

The QUIC handshake incorporates the TLS 1.3 handshake and inherits the cryptographic properties described in [Appendix E.1](#) of [TLS13]. Many of the security properties of QUIC depend on the TLS handshake providing these properties. Any attack on the TLS handshake could affect QUIC.

Any attack on the TLS handshake that compromises the secrecy or uniqueness of session keys affects other security guarantees provided by QUIC that depends on these keys. For instance, migration ([Section 9](#)) depends on the efficacy of confidentiality protections, both for the negotiation of keys using the TLS handshake and for QUIC packet protection, to avoid linkability across network paths.

An attack on the integrity of the TLS handshake might allow an attacker to affect the selection of application protocol or QUIC version.

In addition to the properties provided by TLS, the QUIC handshake provides some defense against DoS attacks on the handshake.

21.12.1.1. Anti-Amplification

Address validation ([Section 8](#)) is used to verify that an entity that claims a given address is able to receive packets at that address. Address validation limits amplification attack targets to addresses for which an attacker is either on-path or off-path.

Prior to validation, endpoints are limited in what they are able to send. During the handshake, a server cannot send more than three times the data it receives; clients that initiate new connections or migrate to a new network path are limited.

21.12.1.2. Server-Side DoS

Computing the server's first flight for a full handshake is potentially expensive, requiring both a signature and a key exchange computation. In order to prevent computational DoS attacks, the Retry packet provides a cheap token exchange mechanism which allows servers to validate a client's IP address prior to doing any expensive computations at the cost of a single round trip. After a successful handshake, servers can issue new tokens to a client which will allow new connection establishment without incurring this cost.

21.12.1.3. On-Path Handshake Termination

An on-path or off-path attacker can force a handshake to fail by replacing or racing Initial packets. Once valid Initial packets have been exchanged, subsequent Handshake packets are protected with the handshake keys and an on-path attacker cannot force handshake failure other than by dropping packets to cause endpoints to abandon the attempt.

An on-path attacker can also replace the addresses of packets on either side and therefore cause the client or server to have an incorrect view of the remote addresses. Such an attack is indistinguishable from the functions performed by a NAT.

21.12.1.4. Parameter Negotiation

The entire handshake is cryptographically protected, with the Initial packets being encrypted with per-version keys and the Handshake and later packets being encrypted with keys derived from the TLS key exchange. Further, parameter negotiation is folded into the TLS transcript and thus provides the same integrity guarantees as ordinary TLS negotiation. An attacker can observe the client's transport parameters (as long as it knows the version-specific keys) but cannot observe the server's transport parameters and cannot influence parameter negotiation.

Connection IDs are unencrypted but integrity protected in all packets.

This version of QUIC does not incorporate a version negotiation mechanism; implementations of incompatible versions will simply fail to establish a connection.

21.12.2. Protected Packets

Packet protection ([Section 12.1](#)) provides authentication and encryption of all packets except Version Negotiation packets, though Initial and Retry packets have limited encryption and authentication based on version-specific keys; see [[QUIC-TLS](#)] for more details. This section considers passive and active attacks against protected packets.

Both on-path and off-path attackers can mount a passive attack in which they save observed packets for an offline attack against packet protection at a future time; this is true for any observer of any packet on any network.

A blind attacker, one who injects packets without being able to observe valid packets for a connection, is unlikely to be successful, since packet protection ensures that valid packets are only generated by endpoints which possess the key material established during the handshake; see [Section 7](#) and [Section 21.12.1](#). Similarly, any active attacker that observes packets and attempts to insert new data or modify existing data in those packets should not be able to generate packets deemed valid by the receiving endpoint.

A spoofing attack, in which an active attacker rewrites unprotected parts of a packet that it forwards or injects, such as the source or destination address, is only effective if the attacker can forward packets to the original endpoint. Packet protection ensures that the packet payloads can only be processed by the endpoints that completed the handshake, and invalid packets are ignored by those endpoints.

An attacker can also modify the boundaries between packets and UDP datagrams, causing multiple packets to be coalesced into a single datagram, or splitting coalesced packets into multiple datagrams. Aside from datagrams containing Initial packets, which require padding, modification of how packets are arranged in datagrams has no functional effect on a connection, although it might change some performance characteristics.

21.12.3. Connection Migration

Connection Migration ([Section 9](#)) provides endpoints with the ability to transition between IP addresses and ports on multiple paths, using one path at a time for transmission and receipt of non-probing frames. Path validation ([Section 8.2](#)) establishes that a peer is both willing and able to receive packets sent on a particular path. This helps reduce the effects of address spoofing by limiting the number of packets sent to a spoofed address.

This section describes the intended security properties of connection migration when under various types of DoS attacks.

21.12.3.1. On-Path Active Attacks

An attacker that can cause a packet it observes to no longer reach its intended destination is considered an on-path attacker. When an attacker is present between a client and server, endpoints are required to send packets through the attacker to establish connectivity on a given path.

An on-path attacker can:

- * Inspect packets

- * Modify IP and UDP packet headers
- * Inject new packets
- * Delay packets
- * Reorder packets
- * Drop packets
- * Split and merge datagrams along packet boundaries

An on-path attacker cannot:

- * Modify an authenticated portion of a packet and cause the recipient to accept that packet

An on-path attacker has the opportunity to modify the packets that it observes, however any modifications to an authenticated portion of a packet will cause it to be dropped by the receiving endpoint as invalid, as packet payloads are both authenticated and encrypted.

In the presence of an on-path attacker, QUIC aims to provide the following properties:

1. An on-path attacker can prevent use of a path for a connection, causing it to fail if it cannot use a different path that does not contain the attacker. This can be achieved by dropping all packets, modifying them so that they fail to decrypt, or other methods.
2. An on-path attacker can prevent migration to a new path for which the attacker is also on-path by causing path validation to fail on the new path.
3. An on-path attacker cannot prevent a client from migrating to a path for which the attacker is not on-path.
4. An on-path attacker can reduce the throughput of a connection by delaying packets or dropping them.
5. An on-path attacker cannot cause an endpoint to accept a packet for which it has modified an authenticated portion of that packet.

21.12.3.2. Off-Path Active Attacks

An off-path attacker is not directly on the path between a client and server, but could be able to obtain copies of some or all packets sent between the client and the server. It is also able to send copies of those packets to either endpoint.

An off-path attacker can:

- * Inspect packets
- * Inject new packets
- * Reorder injected packets

An off-path attacker cannot:

- * Modify any part of a packet
- * Delay packets
- * Drop packets
- * Reorder original packets

An off-path attacker can modify packets that it has observed and inject them back into the network, potentially with spoofed source and destination addresses.

For the purposes of this discussion, it is assumed that an off-path attacker has the ability to observe, modify, and re-inject a packet into the network that will reach the destination endpoint prior to the arrival of the original packet observed by the attacker. In other words, an attacker has the ability to consistently "win" a race with the legitimate packets between the endpoints, potentially causing the original packet to be ignored by the recipient.

It is also assumed that an attacker has the resources necessary to affect NAT state, potentially both causing an endpoint to lose its NAT binding, and an attacker to obtain the same port for use with its traffic.

In the presence of an off-path attacker, QUIC aims to provide the following properties:

1. An off-path attacker can race packets and attempt to become a "limited" on-path attacker.

2. An off-path attacker can cause path validation to succeed for forwarded packets with the source address listed as the off-path attacker as long as it can provide improved connectivity between the client and the server.
3. An off-path attacker cannot cause a connection to close once the handshake has completed.
4. An off-path attacker cannot cause migration to a new path to fail if it cannot observe the new path.
5. An off-path attacker can become a limited on-path attacker during migration to a new path for which it is also an off-path attacker.
6. An off-path attacker can become a limited on-path attacker by affecting shared NAT state such that it sends packets to the server from the same IP address and port that the client originally used.

21.12.3.3. Limited On-Path Active Attacks

A limited on-path attacker is an off-path attacker that has offered improved routing of packets by duplicating and forwarding original packets between the server and the client, causing those packets to arrive before the original copies such that the original packets are dropped by the destination endpoint.

A limited on-path attacker differs from an on-path attacker in that it is not on the original path between endpoints, and therefore the original packets sent by an endpoint are still reaching their destination. This means that a future failure to route copied packets to the destination faster than their original path will not prevent the original packets from reaching the destination.

A limited on-path attacker can:

- * Inspect packets
- * Inject new packets
- * Modify unencrypted packet headers
- * Reorder packets

A limited on-path attacker cannot:

- * Delay packets so that they arrive later than packets sent on the original path
- * Drop packets
- * Modify the authenticated and encrypted portion of a packet and cause the recipient to accept that packet

A limited on-path attacker can only delay packets up to the point that the original packets arrive before the duplicate packets, meaning that it cannot offer routing with worse latency than the original path. If a limited on-path attacker drops packets, the original copy will still arrive at the destination endpoint.

In the presence of a limited on-path attacker, QUIC aims to provide the following properties:

1. A limited on-path attacker cannot cause a connection to close once the handshake has completed.
2. A limited on-path attacker cannot cause an idle connection to close if the client is first to resume activity.
3. A limited on-path attacker can cause an idle connection to be deemed lost if the server is the first to resume activity.

Note that these guarantees are the same guarantees provided for any NAT, for the same reasons.

22. IANA Considerations

This document establishes several registries for the management of codepoints in QUIC. These registries operate on a common set of policies as defined in [Section 22.1](#).

22.1. Registration Policies for QUIC Registries

All QUIC registries allow for both provisional and permanent registration of codepoints. This section documents policies that are common to these registries.

22.1.1. Provisional Registrations

Provisional registration of codepoints are intended to allow for private use and experimentation with extensions to QUIC. Provisional registrations only require the inclusion of the codepoint value and contact information. However, provisional registrations could be reclaimed and reassigned for another purpose.

Provisional registrations require Expert Review, as defined in [Section 4.5 of \[RFC8126\]](#). Designated expert(s) are advised that only registrations for an excessive proportion of remaining codepoint space or the very first unassigned value (see [Section 22.1.2](#)) can be rejected.

Provisional registrations will include a date field that indicates when the registration was last updated. A request to update the date on any provisional registration can be made without review from the designated expert(s).

All QUIC registries include the following fields to support provisional registration:

Value: The assigned codepoint.

Status: "Permanent" or "Provisional".

Specification: A reference to a publicly available specification for the value.

Date: The date of last update to the registration.

Contact: Contact details for the registrant.

Notes: Supplementary notes about the registration.

Provisional registrations MAY omit the Specification and Notes fields, plus any additional fields that might be required for a permanent registration. The Date field is not required as part of requesting a registration as it is set to the date the registration is created or updated.

[22.1.2. Selecting Codepoints](#)

New uses of codepoints from QUIC registries SHOULD use a randomly selected codepoint that excludes both existing allocations and the first unallocated codepoint in the selected space. Requests for multiple codepoints MAY use a contiguous range. This minimizes the risk that differing semantics are attributed to the same codepoint by different implementations. Use of the first codepoint in a range is intended for use by specifications that are developed through the standards process [[STD](#)] and its allocation MUST be negotiated with IANA before use.

For codepoints that are encoded in variable-length integers (Section 16), such as frame types, codepoints that encode to four or eight bytes (that is, values 2^{14} and above) SHOULD be used unless the usage is especially sensitive to having a longer encoding.

Applications to register codepoints in QUIC registries MAY include a codepoint as part of the registration. IANA MUST allocate the selected codepoint unless that codepoint is already assigned or the codepoint is the first unallocated codepoint in the registry.

22.1.3. Reclaiming Provisional Codepoints

A request might be made to remove an unused provisional registration from the registry to reclaim space in a registry, or portion of the registry (such as the 64-16383 range for codepoints that use variable-length encodings). This SHOULD be done only for the codepoints with the earliest recorded date and entries that have been updated less than a year prior SHOULD NOT be reclaimed.

A request to remove a codepoint MUST be reviewed by the designated expert(s). The expert(s) MUST attempt to determine whether the codepoint is still in use. Experts are advised to contact the listed contacts for the registration, plus as wide a set of protocol implementers as possible in order to determine whether any use of the codepoint is known. The expert(s) are advised to allow at least four weeks for responses.

If any use of the codepoints is identified by this search or a request to update the registration is made, the codepoint MUST NOT be reclaimed. Instead, the date on the registration is updated. A note might be added for the registration recording relevant information that was learned.

If no use of the codepoint was identified and no request was made to update the registration, the codepoint MAY be removed from the registry.

This process also applies to requests to change a provisional registration into a permanent registration, except that the goal is not to determine whether there is no use of the codepoint, but to determine that the registration is an accurate representation of any deployed usage.

22.1.4. Permanent Registrations

Permanent registrations in QUIC registries use the Specification Required policy [RFC8126], unless otherwise specified. The designated expert(s) verify that a specification exists and is readily accessible. Expert(s) are encouraged to be biased towards approving registrations unless they are abusive, frivolous, or actively harmful (not merely aesthetically displeasing, or architecturally dubious). The creation of a registry MAY specify additional constraints on permanent registrations.

The creation of a registries MAY identify a range of codepoints where registrations are governed by a different registration policy. For instance, the registries for 62-bit codepoints in this document have stricter policies for codepoints in the range from 0 to 63.

Any stricter requirements for permanent registrations do not prevent provisional registrations for affected codepoints. For instance, a provisional registration for a frame type Section 22.3 of 61 could be requested.

All registrations made by Standards Track publications MUST be permanent.

All registrations in this document are assigned a permanent status and list as contact both the IESG (ietf@ietf.org) and the QUIC working group (quic@ietf.org (mailto:quic@ietf.org)).

22.2. QUIC Transport Parameter Registry

IANA [SHALL add/has added] a registry for "QUIC Transport Parameters" under a "QUIC" heading.

The "QUIC Transport Parameters" registry governs a 62-bit space. This registry follows the registration policy from Section 22.1. Permanent registrations in this registry are assigned using the Specification Required policy [RFC8126].

In addition to the fields in Section 22.1.1, permanent registrations in this registry MUST include the following fields:

Parameter Name: A short mnemonic for the parameter.

The initial contents of this registry are shown in Table 6.

Value	Parameter Name	Specification
0x00	original_destination_connection_id	Section 18.2
0x01	max_idle_timeout	Section 18.2
0x02	stateless_reset_token	Section 18.2
0x03	max_udp_payload_size	Section 18.2
0x04	initial_max_data	Section 18.2
0x05	initial_max_stream_data_bidi_local	Section 18.2
0x06	initial_max_stream_data_bidi_remote	Section 18.2
0x07	initial_max_stream_data_uni	Section 18.2
0x08	initial_max_streams_bidi	Section 18.2
0x09	initial_max_streams_uni	Section 18.2
0x0a	ack_delay_exponent	Section 18.2
0x0b	max_ack_delay	Section 18.2
0x0c	disable_active_migration	Section 18.2
0x0d	preferred_address	Section 18.2
0x0e	active_connection_id_limit	Section 18.2
0x0f	initial_source_connection_id	Section 18.2
0x10	retry_source_connection_id	Section 18.2

Table 6: Initial QUIC Transport Parameters Entries

Additionally, each value of the format "31 * N + 27" for integer values of N (that is, 27, 58, 89, ...) are reserved and MUST NOT be assigned by IANA.

22.3. QUIC Frame Type Registry

IANA [SHALL add/has added] a registry for "QUIC Frame Types" under a "QUIC" heading.

The "QUIC Frame Types" registry governs a 62-bit space. This registry follows the registration policy from [Section 22.1](#). Permanent registrations in this registry are assigned using the Specification Required policy [[RFC8126](#)], except for values between 0x00 and 0x3f (in hexadecimal; inclusive), which are assigned using Standards Action or IESG Approval as defined in [Section 4.9](#) and 4.10 of [[RFC8126](#)].

In addition to the fields in [Section 22.1.1](#), permanent registrations in this registry MUST include the following fields:

Frame Name: A short mnemonic for the frame type.

In addition to the advice in [Section 22.1](#), specifications for new permanent registrations SHOULD describe the means by which an endpoint might determine that it can send the identified type of frame. An accompanying transport parameter registration is expected for most registrations; see [Section 22.2](#). Specifications for permanent registrations also needs to describe the format and assigned semantics of any fields in the frame.

The initial contents of this registry are tabulated in Table 3.

22.4. QUIC Transport Error Codes Registry

IANA [SHALL add/has added] a registry for "QUIC Transport Error Codes" under a "QUIC" heading.

The "QUIC Transport Error Codes" registry governs a 62-bit space. This space is split into three spaces that are governed by different policies. Permanent registrations in this registry are assigned using the Specification Required policy [[RFC8126](#)], except for values between 0x00 and 0x3f (in hexadecimal; inclusive), which are assigned using Standards Action or IESG Approval as defined in [Section 4.9](#) and 4.10 of [[RFC8126](#)].

In addition to the fields in [Section 22.1.1](#), permanent registrations in this registry MUST include the following fields:

Code: A short mnemonic for the parameter.

Description: A brief description of the error code semantics, which MAY be a summary if a specification reference is provided.

The initial contents of this registry are shown in Table 7.

Value	Error	Description	Specification
0x0	NO_ERROR	No error	Section 20
0x1	INTERNAL_ERROR	Implementation error	Section 20
0x2	CONNECTION_REFUSED_ERROR	Server refuses a connection	Section 20
0x3	FLOW_CONTROL_ERROR	Flow control error	Section 20
0x4	STREAM_LIMIT_ERROR	Too many streams opened	Section 20
0x5	STREAM_STATE_ERROR	Frame received in invalid stream state	Section 20
0x6	FINAL_SIZE_ERROR	Change to final size	Section 20
0x7	FRAME_ENCODING_ERROR	Frame encoding error	Section 20
0x8	TRANSPORT_PARAMETER_ERROR	Error in transport parameters	Section 20
0x9	CONNECTION_ID_LIMIT_ERROR	Too many connection IDs received	Section 20
0xA	PROTOCOL_VIOLATION	Generic protocol violation	Section 20
0xB	INVALID_TOKEN	Invalid Token Received	Section 20
0xC	APPLICATION_ERROR	Application error	Section 20
0xD	CRYPTO_BUFFER_EXCEEDED	CRYPTO data buffer overflowed	Section 20

Table 7: Initial QUIC Transport Error Codes Entries

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[Appendix A](#). Sample Packet Number Decoding Algorithm

The pseudo-code in Figure 44 shows how an implementation can decode packet numbers after header protection has been removed.

```
DecodePacketNumber(largest_pn, truncated_pn, pn_nbits):
    expected_pn = largest_pn + 1
    pn_win      = 1 << pn_nbits
    pn_hwin     = pn_win / 2
    pn_mask     = pn_win - 1
    // The incoming packet number should be greater than
    // expected_pn - pn_hwin and less than or equal to
    // expected_pn + pn_hwin
    //
    // This means we can't just strip the trailing bits from
    // expected_pn and add the truncated_pn because that might
    // yield a value outside the window.
    //
    // The following code calculates a candidate value and
    // makes sure it's within the packet number window.
    // Note the extra checks to prevent overflow and underflow.
    candidate_pn = (expected_pn & ~pn_mask) | truncated_pn
    if candidate_pn <= expected_pn - pn_hwin and
        candidate_pn < (1 << 62) - pn_win:
        return candidate_pn + pn_win
    if candidate_pn > expected_pn + pn_hwin and
        candidate_pn >= pn_win:
        return candidate_pn - pn_win
    return candidate_pn
```

Figure 44: Sample Packet Number Decoding Algorithm

Appendix B. Sample ECN Validation Algorithm

Each time an endpoint commences sending on a new network path, it determines whether the path supports ECN; see [Section 13.4](#). If the path supports ECN, the goal is to use ECN. Endpoints might also periodically reassess a path that was determined to not support ECN.

This section describes one method for testing new paths. This algorithm is intended to show how a path might be tested for ECN support. Endpoints can implement different methods.

The path is assigned an ECN state that is one of "testing", "unknown", "failed", or "capable". On paths with a "testing" or "capable" state the endpoint sends packets with an ECT marking, by default ECT(0); otherwise, the endpoint sends unmarked packets.

To start testing a path, the ECN state is set to "testing" and existing ECN counts are remembered as a baseline.

The testing period runs for a number of packets or round-trip times, as determined by the endpoint. The goal is not to limit the duration of the testing period, but to ensure that enough marked packets are sent for received ECN counts to provide a clear indication of how the path treats marked packets. [Section 13.4.2.2](#) suggests limiting this to 10 packets or 3 round-trip times.

After the testing period ends, the ECN state for the path becomes "unknown". From the "unknown" state, successful validation of the ECN counts an ACK frame (see [Section 13.4.2.2](#)) causes the ECN state for the path to become "capable", unless no marked packet has been acknowledged.

If validation of ECN counts fails at any time, the ECN state for the affected path becomes "failed". An endpoint can also mark the ECN state for a path as "failed" if marked packets are all declared lost or if they are all CE marked.

Following this algorithm ensures that ECN is rarely disabled for paths that properly support ECN. Any path that incorrectly modifies markings will cause ECN to be disabled. For those rare cases where marked packets are discarded by the path, the short duration of the testing period limits the number of losses incurred.

[Appendix C](#). Change Log

RFC Editor's Note: Please remove this section prior to publication of a final version of this document.

Issue and pull request numbers are listed with a leading octothorp.

[C.1](#). Since [draft-ietf-quick-transport-28](#)

- * Made SERVER_BUSY error (0x2) more generic, now CONNECTION_REFUSED (#3709, #3690, #3694)
- * Allow TRANSPORT_PARAMETER_ERROR when validating connection IDs (#3703, #3691)
- * Integrate QUIC-specific language from [draft-ietf-tsvwg-datagram-plpmtud](#) (#3695, #3702)
- * disable_active_migration does not apply to the addresses offered in server_preferred_address (#3608, #3670)

[C.2](#). Since [draft-ietf-quick-transport-27](#)

- * Allowed CONNECTION_CLOSE in any packet number space, with a requirement to use a new transport-level error for application-specific errors in Initial and Handshake packets (#3430, #3435, #3440)
- * Clearer requirements for address validation (#2125, #3327)
- * Security analysis of handshake and migration (#2143, #2387, #2925)
- * The entire payload of a datagram is used when counting bytes for mitigating amplification attacks (#3333, #3470)
- * Connection IDs can be used at any time, including in the handshake (#3348, #3560, #3438, #3565)
- * Only one ACK should be sent for each instance of reordering (#3357, #3361)
- * Remove text allowing a server to proceed with a bad Retry token (#3396, #3398)
- * Ignore active_connection_id_limit with a zero-length connection ID (#3427, #3426)
- * Require active_connection_id_limit be remembered for 0-RTT (#3423, #3425)
- * Require ack_delay not be remembered for 0-RTT (#3433, #3545)
- * Redefined max_packet_size to max_udp_datagram_size (#3471, #3473)
- * Guidance on limiting outstanding attempts to retire connection IDs (#3489, #3509, #3557, #3547)
- * Restored text on dropping bogus Version Negotiation packets (#3532, #3533)
- * Clarified that largest acknowledged needs to be saved, but not necessarily signaled in all cases (#3541, #3581)
- * Addressed linkability risk with the use of preferred_address (#3559, #3563)
- * Added authentication of handshake connection IDs (#3439, #3499)
- * Opening a stream in the wrong direction is an error (#3527)

C.3. Since [draft-ietf-quic-transport-26](#)

- * Change format of transport parameters to use varints (#3294, #3169)

C.4. Since [draft-ietf-quic-transport-25](#)

- * Define the use of CONNECTION_CLOSE prior to establishing connection state (#3269, #3297, #3292)
- * Allow use of address validation tokens after client address changes (#3307, #3308)
- * Define the timer for address validation (#2910, #3339)

C.5. Since [draft-ietf-quic-transport-24](#)

- * Added HANDSHAKE_DONE to signal handshake confirmation (#2863, #3142, #3145)
- * Add integrity check to Retry packets (#3014, #3274, #3120)
- * Specify handling of reordered NEW_CONNECTION_ID frames (#3194, #3202)
- * Require checking of sequence numbers in RETIRE_CONNECTION_ID (#3037, #3036)
- * active_connection_id_limit is enforced (#3193, #3197, #3200, #3201)
- * Correct overflow in packet number decode algorithm (#3187, #3188)
- * Allow use of CRYPTO_BUFFER_EXCEEDED for CRYPTO frame errors (#3258, #3186)
- * Define applicability and scope of NEW_TOKEN (#3150, #3152, #3155, #3156)
- * Tokens from Retry and NEW_TOKEN must be differentiated (#3127, #3128)
- * Allow CONNECTION_CLOSE in response to invalid token (#3168, #3107)
- * Treat an invalid CONNECTION_CLOSE as an invalid frame (#2475, #3230, #3231)

- * Throttle when sending CONNECTION_CLOSE after discarding state (#3095, #3157)
- * Application-variant of CONNECTION_CLOSE can only be sent in 0-RTT or 1-RTT packets (#3158, #3164)
- * Advise sending while blocked to avoid idle timeout (#2744, #3266)
- * Define error codes for invalid frames (#3027, #3042)
- * Idle timeout is symmetric (#2602, #3099)
- * Prohibit IP fragmentation (#3243, #3280)
- * Define the use of provisional registration for all registries (#3109, #3020, #3102, #3170)
- * Packets on one path must not adjust values for a different path (#2909, #3139)

C.6. Since [draft-ietf-quic-transport-23](#)

- * Allow ClientHello to span multiple packets (#2928, #3045)
- * Client Initial size constraints apply to UDP datagram payload (#3053, #3051)
- * Stateless reset changes (#2152, #2993)
 - tokens need to be compared in constant time
 - detection uses UDP datagrams, not packets
 - tokens cannot be reused (#2785, #2968)
- * Clearer rules for sharing of UDP ports and use of connection IDs when doing so (#2844, #2851)
- * A new connection ID is necessary when responding to migration (#2778, #2969)
- * Stronger requirements for connection ID retirement (#3046, #3096)
- * NEW_TOKEN cannot be empty (#2978, #2977)
- * PING can be sent at any encryption level (#3034, #3035)
- * CONNECTION_CLOSE is not ack-eliciting (#3097, #3098)

- * Frame encoding error conditions updated (#3027, #3042)
- * Non-ack-eliciting packets cannot be sent in response to non-ack-eliciting packets (#3100, #3104)
- * Servers have to change connection IDs in Retry (#2837, #3147)

C.7. Since [draft-ietf-quic-transport-22](#)

- * Rules for preventing correlation by connection ID tightened (#2084, #2929)
- * Clarified use of CONNECTION_CLOSE in Handshake packets (#2151, #2541, #2688)
- * Discourage regressions of largest acknowledged in ACK (#2205, #2752)
- * Improved robustness of validation process for ECN counts (#2534, #2752)
- * Require endpoints to ignore spurious migration attempts (#2342, #2893)
- * Transport parameter for disabling migration clarified to allow NAT rebinding (#2389, #2893)
- * Document principles for defining new error codes (#2388, #2880)
- * Reserve transport parameters for greasing (#2550, #2873)
- * A maximum ACK delay of 0 is used for handshake packet number spaces (#2646, #2638)
- * Improved rules for use of congestion control state on new paths (#2685, #2918)
- * Removed recommendation to coordinate spin for multiple connections that share a path (#2763, #2882)
- * Allow smaller stateless resets and recommend a smaller minimum on packets that might trigger a stateless reset (#2770, #2869, #2927, #3007).
- * Provide guidance around the interface to QUIC as used by application protocols (#2805, #2857)

- * Frames other than STREAM can cause STREAM_LIMIT_ERROR (#2825, #2826)
- * Tighter rules about processing of rejected 0-RTT packets (#2829, #2840, #2841)
- * Explanation of the effect of Retry on 0-RTT packets (#2842, #2852)
- * Cryptographic handshake needs to provide server transport parameter encryption (#2920, #2921)
- * Moved ACK generation guidance from recovery draft to transport draft (#1860, #2916).

C.8. Since [draft-ietf-quic-transport-21](#)

- * Connection ID lengths are now one octet, but limited in version 1 to 20 octets of length (#2736, #2749)

C.9. Since [draft-ietf-quic-transport-20](#)

- * Error codes are encoded as variable-length integers (#2672, #2680)
- * NEW_CONNECTION_ID includes a request to retire old connection IDs (#2645, #2769)
- * Tighter rules for generating and explicitly eliciting ACK frames (#2546, #2794)
- * Recommend having only one packet per encryption level in a datagram (#2308, #2747)
- * More normative language about use of stateless reset (#2471, #2574)
- * Allow reuse of stateless reset tokens (#2732, #2733)
- * Allow, but not require, enforcing non-duplicate transport parameters (#2689, #2691)
- * Added an active_connection_id_limit transport parameter (#1994, #1998)
- * max_ack_delay transport parameter defaults to 0 (#2638, #2646)
- * When sending 0-RTT, only remembered transport parameters apply (#2458, #2360, #2466, #2461)

- * Define handshake completion and confirmation; define clearer rules when it encryption keys should be discarded (#2214, #2267, #2673)
- * Prohibit path migration prior to handshake confirmation (#2309, #2370)
- * PATH_RESPONSE no longer needs to be received on the validated path (#2582, #2580, #2579, #2637)
- * PATH_RESPONSE frames are not stored and retransmitted (#2724, #2729)
- * Document hack for enabling routing of ICMP when doing PMTU probing (#1243, #2402)

C.10. Since [draft-ietf-quit-transport-19](#)

- * Refine discussion of 0-RTT transport parameters (#2467, #2464)
- * Fewer transport parameters need to be remembered for 0-RTT (#2624, #2467)
- * Spin bit text incorporated (#2564)
- * Close the connection when maximum stream ID in MAX_STREAMS exceeds $2^{62} - 1$ (#2499, #2487)
- * New connection ID required for intentional migration (#2414, #2413)
- * Connection ID issuance can be rate-limited (#2436, #2428)
- * The "QUIC bit" is ignored in Version Negotiation (#2400, #2561)
- * Initial packets from clients need to be padded to 1200 unless a Handshake packet is sent as well (#2522, #2523)
- * CRYPTO frames can be discarded if too much data is buffered (#1834, #2524)
- * Stateless reset uses a short header packet (#2599, #2600)

C.11. Since [draft-ietf-quit-transport-18](#)

- * Removed version negotiation; version negotiation, including authentication of the result, will be addressed in the next version of QUIC (#1773, #2313)

- * Added discussion of the use of IPv6 flow labels (#2348, #2399)
- * A connection ID can't be retired in a packet that uses that connection ID (#2101, #2420)
- * Idle timeout transport parameter is in milliseconds (from seconds) (#2453, #2454)
- * Endpoints are required to use new connection IDs when they use new network paths (#2413, #2414)
- * Increased the set of permissible frames in 0-RTT (#2344, #2355)

C.12. Since [draft-ietf-quic-transport-17](#)

- * Stream-related errors now use `STREAM_STATE_ERROR` (#2305)
- * Endpoints discard initial keys as soon as handshake keys are available (#1951, #2045)
- * Expanded conditions for ignoring ICMP packet too big messages (#2108, #2161)
- * Remove rate control from `PATH_CHALLENGE/PATH_RESPONSE` (#2129, #2241)
- * Endpoints are permitted to discard malformed initial packets (#2141)
- * Clarified ECN implementation and usage requirements (#2156, #2201)
- * Disable ECN count verification for packets that arrive out of order (#2198, #2215)
- * Use Probe Timeout (PTO) instead of RTO (#2206, #2238)
- * Loosen constraints on retransmission of ACK ranges (#2199, #2245)
- * Limit Retry and Version Negotiation to once per datagram (#2259, #2303)
- * Set a maximum value for `max_ack_delay` transport parameter (#2282, #2301)
- * Allow server preferred address for both IPv4 and IPv6 (#2122, #2296)

- * Corrected requirements for migration to a preferred address (#2146, #2349)
- * ACK of non-existent packet is illegal (#2298, #2302)

C.13. Since [draft-ietf-quic-transport-16](#)

- * Stream limits are defined as counts, not maximums (#1850, #1906)
- * Require amplification attack defense after closing (#1905, #1911)
- * Remove reservation of application error code 0 for STOPPING (#1804, #1922)
- * Renumbered frames (#1945)
- * Renumbered transport parameters (#1946)
- * Numeric transport parameters are expressed as varints (#1608, #1947, #1955)
- * Reorder the NEW_CONNECTION_ID frame (#1952, #1963)
- * Rework the first byte (#2006)
 - Fix the 0x40 bit
 - Change type values for long header
 - Add spin bit to short header (#631, #1988)
 - Encrypt the remainder of the first byte (#1322)
 - Move packet number length to first byte
 - Move ODCIL to first byte of retry packets
 - Simplify packet number protection (#1575)
- * Allow STOP_SENDING to open a remote bidirectional stream (#1797, #2013)
- * Added mitigation for off-path migration attacks (#1278, #1749, #2033)
- * Don't let the PMTU to drop below 1280 (#2063, #2069)
- * Require peers to replace retired connection IDs (#2085)

- * Servers are required to ignore Version Negotiation packets (#2088)
- * Tokens are repeated in all Initial packets (#2089)
- * Clarified how PING frames are sent after loss (#2094)
- * Initial keys are discarded once Handshake are available (#1951, #2045)
- * ICMP PTB validation clarifications (#2161, #2109, #2108)

C.14. Since [draft-ietf-quic-transport-15](#)

Substantial editorial reorganization; no technical changes.

C.15. Since [draft-ietf-quic-transport-14](#)

- * Merge ACK and ACK_ECN (#1778, #1801)
- * Explicitly communicate max_ack_delay (#981, #1781)
- * Validate original connection ID after Retry packets (#1710, #1486, #1793)
- * Idle timeout is optional and has no specified maximum (#1765)
- * Update connection ID handling; add RETIRE_CONNECTION_ID type (#1464, #1468, #1483, #1484, #1486, #1495, #1729, #1742, #1799, #1821)
- * Include a Token in all Initial packets (#1649, #1794)
- * Prevent handshake deadlock (#1764, #1824)

C.16. Since [draft-ietf-quic-transport-13](#)

- * Streams open when higher-numbered streams of the same type open (#1342, #1549)
- * Split initial stream flow control limit into 3 transport parameters (#1016, #1542)
- * All flow control transport parameters are optional (#1610)
- * Removed UNSOLICITED_PATH_RESPONSE error code (#1265, #1539)
- * Permit stateless reset in response to any packet (#1348, #1553)

- * Recommended defense against stateless reset spoofing (#1386, #1554)
- * Prevent infinite stateless reset exchanges (#1443, #1627)
- * Forbid processing of the same packet number twice (#1405, #1624)
- * Added a packet number decoding example (#1493)
- * More precisely define idle timeout (#1429, #1614, #1652)
- * Corrected format of Retry packet and prevented looping (#1492, #1451, #1448, #1498)
- * Permit 0-RTT after receiving Version Negotiation or Retry (#1507, #1514, #1621)
- * Permit Retry in response to 0-RTT (#1547, #1552)
- * Looser verification of ECN counters to account for ACK loss (#1555, #1481, #1565)
- * Remove frame type field from APPLICATION_CLOSE (#1508, #1528)

C.17. Since [draft-ietf-quic-transport-12](#)

- * Changes to integration of the TLS handshake (#829, #1018, #1094, #1165, #1190, #1233, #1242, #1252, #1450, #1458)
 - The cryptographic handshake uses CRYPTO frames, not stream 0
 - QUIC packet protection is used in place of TLS record protection
 - Separate QUIC packet number spaces are used for the handshake
 - Changed Retry to be independent of the cryptographic handshake
 - Added NEW_TOKEN frame and Token fields to Initial packet
 - Limit the use of HelloRetryRequest to address TLS needs (like key shares)
- * Enable server to transition connections to a preferred address (#560, #1251, #1373)
- * Added ECN feedback mechanisms and handling; new ACK_ECN frame (#804, #805, #1372)

- * Changed rules and recommendations for use of new connection IDs (#1258, #1264, #1276, #1280, #1419, #1452, #1453, #1465)
- * Added a transport parameter to disable intentional connection migration (#1271, #1447)
- * Packets from different connection ID can't be coalesced (#1287, #1423)
- * Fixed sampling method for packet number encryption; the length field in long headers includes the packet number field in addition to the packet payload (#1387, #1389)
- * Stateless Reset is now symmetric and subject to size constraints (#466, #1346)
- * Added frame type extension mechanism (#58, #1473)

C.18. Since [draft-ietf-quic-transport-11](#)

- * Enable server to transition connections to a preferred address (#560, #1251)
- * Packet numbers are encrypted (#1174, #1043, #1048, #1034, #850, #990, #734, #1317, #1267, #1079)
- * Packet numbers use a variable-length encoding (#989, #1334)
- * STREAM frames can now be empty (#1350)

C.19. Since [draft-ietf-quic-transport-10](#)

- * Swap payload length and packed number fields in long header (#1294)
- * Clarified that CONNECTION_CLOSE is allowed in Handshake packet (#1274)
- * Spin bit reserved (#1283)
- * Coalescing multiple QUIC packets in a UDP datagram (#1262, #1285)
- * A more complete connection migration (#1249)
- * Refine opportunistic ACK defense text (#305, #1030, #1185)
- * A Stateless Reset Token isn't mandatory (#818, #1191)

- * Removed implicit stream opening (#896, #1193)
- * An empty STREAM frame can be used to open a stream without sending data (#901, #1194)
- * Define stream counts in transport parameters rather than a maximum stream ID (#1023, #1065)
- * STOP_SENDING is now prohibited before streams are used (#1050)
- * Recommend including ACK in Retry packets and allow PADDING (#1067, #882)
- * Endpoints now become closing after an idle timeout (#1178, #1179)
- * Remove implication that Version Negotiation is sent when a packet of the wrong version is received (#1197)

C.20. Since [draft-ietf-quic-transport-09](#)

- * Added PATH_CHALLENGE and PATH_RESPONSE frames to replace PING with Data and PONG frame. Changed ACK frame type from 0x0e to 0x0d. (#1091, #725, #1086)
- * A server can now only send 3 packets without validating the client address (#38, #1090)
- * Delivery order of stream data is no longer strongly specified (#252, #1070)
- * Rework of packet handling and version negotiation (#1038)
- * Stream 0 is now exempt from flow control until the handshake completes (#1074, #725, #825, #1082)
- * Improved retransmission rules for all frame types: information is retransmitted, not packets or frames (#463, #765, #1095, #1053)
- * Added an error code for server busy signals (#1137)
- * Endpoints now set the connection ID that their peer uses. Connection IDs are variable length. Removed the omit_connection_id transport parameter and the corresponding short header flag. (#1089, #1052, #1146, #821, #745, #821, #1166, #1151)

C.21. Since [draft-ietf-quic-transport-08](#)

- * Clarified requirements for BLOCKED usage (#65, #924)
- * BLOCKED frame now includes reason for blocking (#452, #924, #927, #928)
- * GAP limitation in ACK Frame (#613)
- * Improved PMTUD description (#614, #1036)
- * Clarified stream state machine (#634, #662, #743, #894)
- * Reserved versions don't need to be generated deterministically (#831, #931)
- * You don't always need the draining period (#871)
- * Stateless reset clarified as version-specific (#930, #986)
- * `initial_max_stream_id_x` transport parameters are optional (#970, #971)
- * Ack Delay assumes a default value during the handshake (#1007, #1009)
- * Removed transport parameters from `NewSessionTicket` (#1015)

C.22. Since [draft-ietf-quic-transport-07](#)

- * The long header now has version before packet number (#926, #939)
- * Rename and consolidate packet types (#846, #822, #847)
- * Packet types are assigned new codepoints and the Connection ID Flag is inverted (#426, #956)
- * Removed type for Version Negotiation and use Version 0 (#963, #968)
- * Streams are split into unidirectional and bidirectional (#643, #656, #720, #872, #175, #885)
 - Stream limits now have separate uni- and bi-directional transport parameters (#909, #958)
 - Stream limit transport parameters are now optional and default to 0 (#970, #971)

- * The stream state machine has been split into read and write (#634, #894)
- * Employ variable-length integer encodings throughout (#595)
- * Improvements to connection close
 - Added distinct closing and draining states (#899, #871)
 - Draining period can terminate early (#869, #870)
 - Clarifications about stateless reset (#889, #890)
- * Address validation for connection migration (#161, #732, #878)
- * Clearly defined retransmission rules for BLOCKED (#452, #65, #924)
- * negotiated_version is sent in server transport parameters (#710, #959)
- * Increased the range over which packet numbers are randomized (#864, #850, #964)

C.23. Since [draft-ietf-quic-transport-06](#)

- * Replaced FNV-1a with AES-GCM for all "Cleartext" packets (#554)
- * Split error code space between application and transport (#485)
- * Stateless reset token moved to end (#820)
- * 1-RTT-protected long header types removed (#848)
- * No acknowledgments during draining period (#852)
- * Remove "application close" as a separate close type (#854)
- * Remove timestamps from the ACK frame (#841)
- * Require transport parameters to only appear once (#792)

C.24. Since [draft-ietf-quic-transport-05](#)

- * Stateless token is server-only (#726)
- * Refactor section on connection termination (#733, #748, #328, #177)

- * Limit size of Version Negotiation packet (#585)
- * Clarify when and what to ack (#736)
- * Renamed STREAM_ID_NEEDED to STREAM_ID_BLOCKED
- * Clarify Keep-alive requirements (#729)

C.25. Since [draft-ietf-quick-transport-04](#)

- * Introduce STOP_SENDING frame, RESET_STREAM only resets in one direction (#165)
- * Removed GOAWAY; application protocols are responsible for graceful shutdown (#696)
- * Reduced the number of error codes (#96, #177, #184, #211)
- * Version validation fields can't move or change (#121)
- * Removed versions from the transport parameters in a NewSessionTicket message (#547)
- * Clarify the meaning of "bytes in flight" (#550)
- * Public reset is now stateless reset and not visible to the path (#215)
- * Reordered bits and fields in STREAM frame (#620)
- * Clarifications to the stream state machine (#572, #571)
- * Increased the maximum length of the Largest Acknowledged field in ACK frames to 64 bits (#629)
- * truncate_connection_id is renamed to omit_connection_id (#659)
- * CONNECTION_CLOSE terminates the connection like TCP RST (#330, #328)
- * Update labels used in HKDF-Expand-Label to match TLS 1.3 (#642)

C.26. Since [draft-ietf-quick-transport-03](#)

- * Change STREAM and RESET_STREAM layout
- * Add MAX_STREAM_ID settings

C.27. Since [draft-ietf-quic-transport-02](#)

- * The size of the initial packet payload has a fixed minimum (#267, #472)
- * Define when Version Negotiation packets are ignored (#284, #294, #241, #143, #474)
- * The 64-bit FNV-1a algorithm is used for integrity protection of unprotected packets (#167, #480, #481, #517)
- * Rework initial packet types to change how the connection ID is chosen (#482, #442, #493)
- * No timestamps are forbidden in unprotected packets (#542, #429)
- * Cryptographic handshake is now on stream 0 (#456)
- * Remove congestion control exemption for cryptographic handshake (#248, #476)
- * Version 1 of QUIC uses TLS; a new version is needed to use a different handshake protocol (#516)
- * STREAM frames have a reduced number of offset lengths (#543, #430)
- * Split some frames into separate connection- and stream- level frames (#443)
 - WINDOW_UPDATE split into MAX_DATA and MAX_STREAM_DATA (#450)
 - BLOCKED split to match WINDOW_UPDATE split (#454)
 - Define STREAM_ID_NEEDED frame (#455)
- * A NEW_CONNECTION_ID frame supports connection migration without linkability (#232, #491, #496)
- * Transport parameters for 0-RTT are retained from a previous connection (#405, #513, #512)
 - A client in 0-RTT no longer required to reset excess streams (#425, #479)
- * Expanded security considerations (#440, #444, #445, #448)

C.28. Since [draft-ietf-quic-transport-01](#)

- * Defined short and long packet headers (#40, #148, #361)
- * Defined a versioning scheme and stable fields (#51, #361)
- * Define reserved version values for "greasing" negotiation (#112, #278)
- * The initial packet number is randomized (#35, #283)
- * Narrow the packet number encoding range requirement (#67, #286, #299, #323, #356)
- * Defined client address validation (#52, #118, #120, #275)
- * Define transport parameters as a TLS extension (#49, #122)
- * SCUP and COPT parameters are no longer valid (#116, #117)
- * Transport parameters for 0-RTT are either remembered from before, or assume default values (#126)
- * The server chooses connection IDs in its final flight (#119, #349, #361)
- * The server echoes the Connection ID and packet number fields when sending a Version Negotiation packet (#133, #295, #244)
- * Defined a minimum packet size for the initial handshake packet from the client (#69, #136, #139, #164)
- * Path MTU Discovery (#64, #106)
- * The initial handshake packet from the client needs to fit in a single packet (#338)
- * Forbid acknowledgment of packets containing only ACK and PADDING (#291)
- * Require that frames are processed when packets are acknowledged (#381, #341)
- * Removed the STOP_WAITING frame (#66)
- * Don't require retransmission of old timestamps for lost ACK frames (#308)
- * Clarified that frames are not retransmitted, but the information in them can be (#157, #298)

- * Error handling definitions (#335)
- * Split error codes into four sections (#74)
- * Forbid the use of Public Reset where CONNECTION_CLOSE is possible (#289)
- * Define packet protection rules (#336)
- * Require that stream be entirely delivered or reset, including acknowledgment of all STREAM frames or the RESET_STREAM, before it closes (#381)
- * Remove stream reservation from state machine (#174, #280)
- * Only stream 1 does not contribute to connection-level flow control (#204)
- * Stream 1 counts towards the maximum concurrent stream limit (#201, #282)
- * Remove connection-level flow control exclusion for some streams (except 1) (#246)
- * RESET_STREAM affects connection-level flow control (#162, #163)
- * Flow control accounting uses the maximum data offset on each stream, rather than bytes received (#378)
- * Moved length-determining fields to the start of STREAM and ACK (#168, #277)
- * Added the ability to pad between frames (#158, #276)
- * Remove error code and reason phrase from GOAWAY (#352, #355)
- * GOAWAY includes a final stream number for both directions (#347)
- * Error codes for RESET_STREAM and CONNECTION_CLOSE are now at a consistent offset (#249)
- * Defined priority as the responsibility of the application protocol (#104, #303)

C.29. Since [draft-ietf-quic-transport-00](#)

- * Replaced DIVERSIFICATION_NONCE flag with KEY_PHASE flag

- * Defined versioning
- * Reworked description of packet and frame layout
- * Error code space is divided into regions for each component
- * Use big endian for all numeric values

C.30. Since [draft-hamilton-quick-transport-protocol-01](#)

- * Adopted as base for [draft-ietf-quick-tls](#)
- * Updated authors/editors list
- * Added IANA Considerations section
- * Moved Contributors and Acknowledgments to appendices

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